



26.10.20 12:00

# CSCI 150

## Introduction to Digital and Computer System Design

### Midterm Review I



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2020 Fall Semester (S3)

# Overview

- Focus: Review
- Architecture: Combinational Logic Circuit
- Textbook v4: Ch1-4; v5: Ch1-3
- Core Ideas:
  1. Digital Information Representation (Lecture 1)
  2. Combinational Logic Circuits (Lecture 2)
  3. Combinational Functional Blocks, Arithmetic Blocks (Lecture 3)

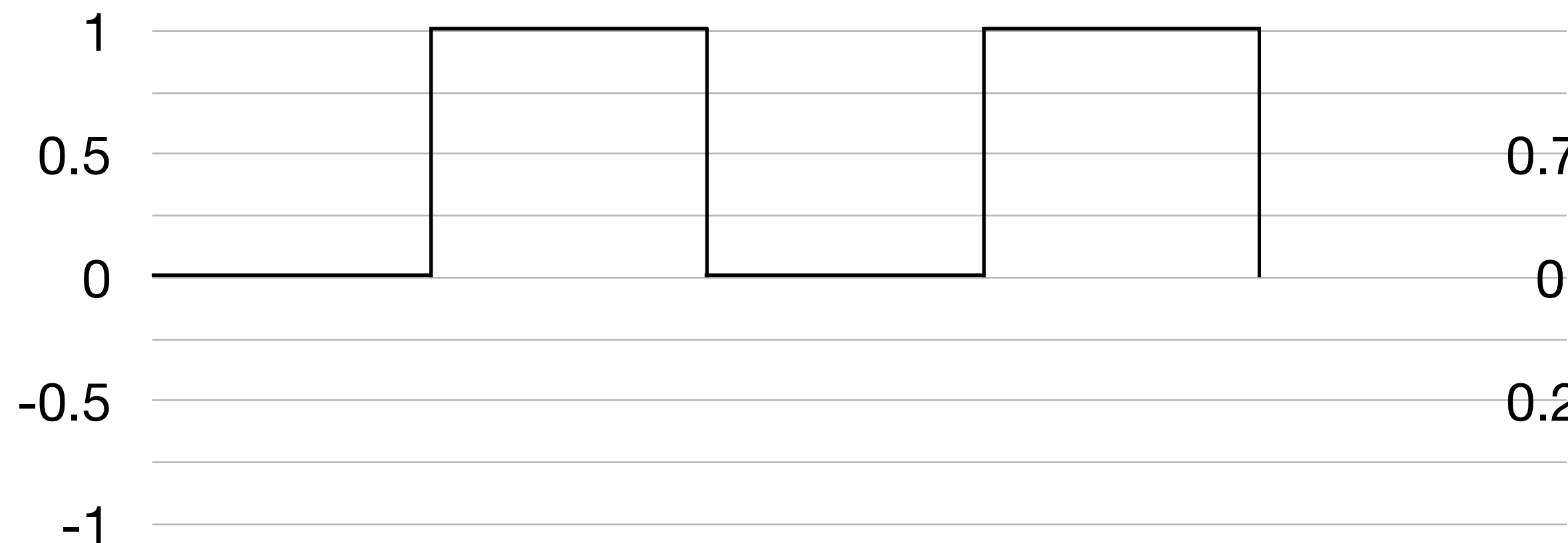
# Lecture 1: Digital Information Representation

Analog vs Digital circuits; Modern computer architectures; Embedded systems;  
Number Systems; Conversions;  
Arithmetic Operations; Alphanumeric Codes

# Analog vs Digital circuits

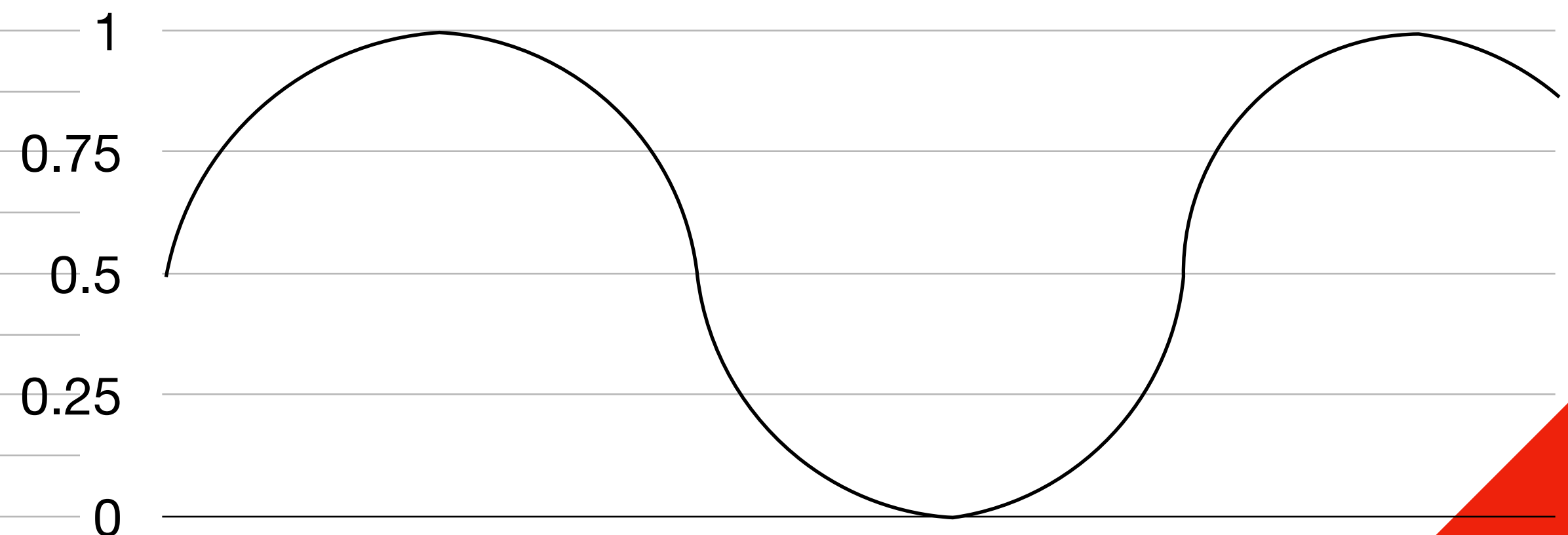
- Digital Circuits

- Process digital signals
- Current/Voltage represent discrete logical and numeric values



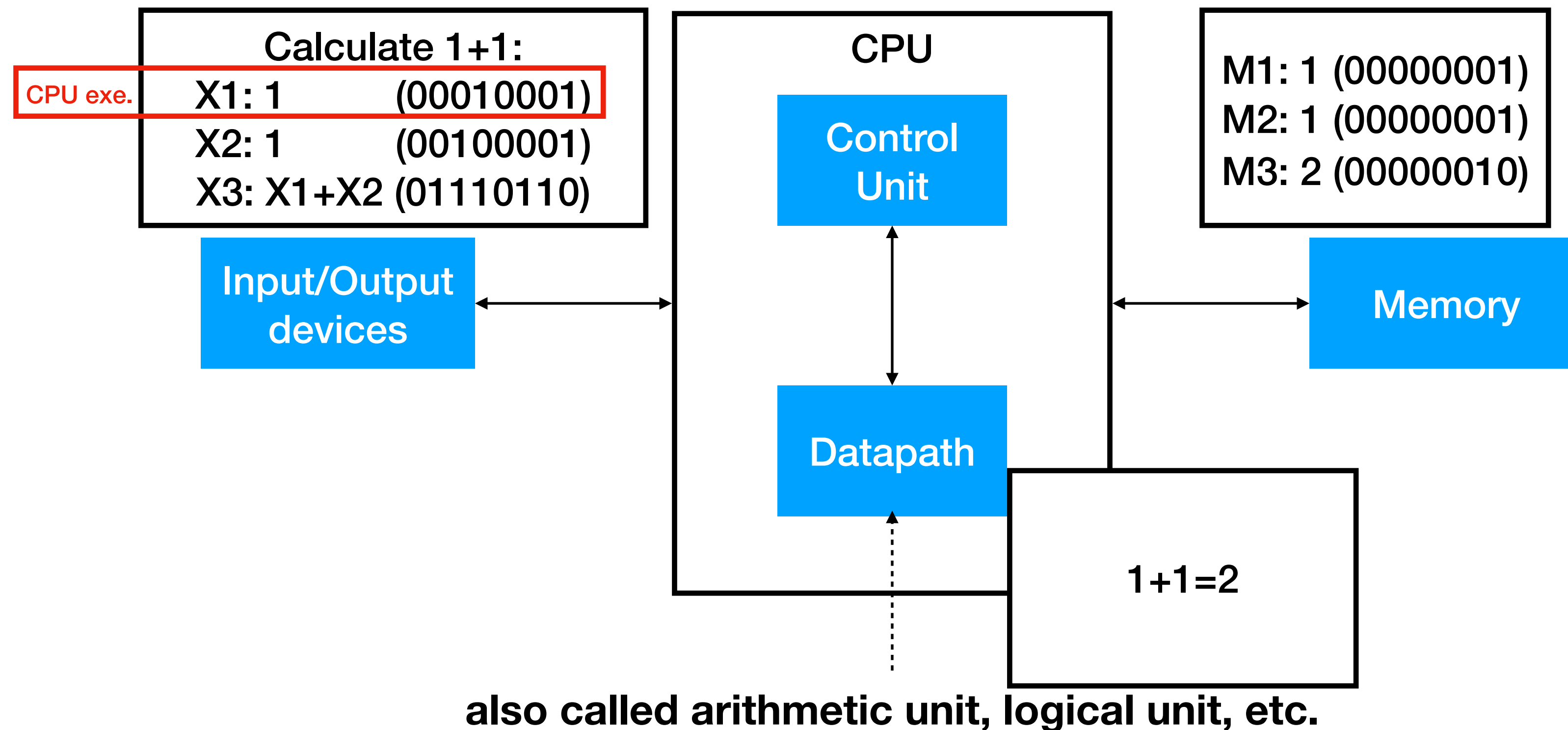
- Analog Circuits

- Process analog signals
- Current/Voltage vary continuously to represent information



# Von Neumann Architecture

A very rough example



# Computer

What's it like compared to a human?

- Input/Output devices
  - Interaction (Mouth, hands and feet, eyes, etc.)
- CPU + Memory
  - Processing information, thinking (Brain, short-term memory)
- Storage?
  - Part of I/O devices (Books, long-term memory)

Concept

# Embedded Systems

- Similar to computers: processes information
- Difference
  - Function is usually simpler, and very very specific
  - Not programmable

# Decimal System

7 2 4 . 0 5  
2 1 0 -1 -2

- Numbers as strings of digits, each ranging from 0-9
- The decimal system is of base(radix) 10



# Decimal System

$$\begin{array}{cccccc} 7 & 2 & 4 & . & 0 & 5 \\ 2 & 1 & 0 & -1 & -2 & \end{array}$$
$$= 7 \times 10^2 + 2 \times 10^1 + 4 \times 10^0 + 0 \times 10^{-1} + 5 \times 10^{-2}$$

- Numbers as strings of digits, each ranging from 0-9
- The decimal system is of base(radix) 10

# Numbers of base N

- Default base: 10
- When there are numbers represented in different bases, attach base
  - Decimal:  $754.05 \rightarrow (754.05)_{10}$
  - e.g. Base 5:  $(432.1)_5 = ?$

$$= 4 \times 5^2 + 3 \times 5^1 + 2 \times 5^0 + 1 \times 5^{-1} = (117.2)_{10}$$

# Binary Systems in Computers

- Every 8bit is called a Byte
- $1,024 = 2^{10}$  is called K (Kilo)
- $1,024 \times 1,024 = 2^{20}$  is called M (Mega)
- $1,024 \times 1,024 \times 1,024 = 2^{40}$  is called G (Giga)
- Tera, Peta, Exa, Zetta, Yotta

# Binary Systems in Computers



- What is the difference between MBps and Mbps?
- MegaBytes per second vs MegaBits per second
- 8x difference!



Concept

# Octal and Hexadecimal Systems

- Octal: base 8
  - digits: 0-7
- Hexadecimal: base 16
  - digits: 0-9, A-F (10-15)

# Conversions

10	9	8	7	6	5	4	3	2	1
1024	512	256	128	64	32	16	8	4	2

- Binary-to-  
Octal: 3bits per octal digit  
Hexadecimal: 4bits per hexa digit  
Decimal: use the chart
- Decimal-to-  
Binary: use the chart  
Oct/Hex: do binary first

# Arithmetics

- The same as decimal (mostly)
- 

$$\begin{array}{r} 0010 \\ +0011 \\ \hline 0101 \end{array} \quad \begin{array}{r} 0101 \\ -0011 \\ \hline 0010 \end{array}$$

Example (binary)

# Arithmetics

## OCTAL Multiplication

Octal

$$\begin{array}{r} \text{Octal} \\ 762 \\ \times 54 \\ \hline 4672 \\ 3710 \\ \hline 43772 \end{array}$$

$$5 \times 2 = 12$$

$$5 \times 6 + 1 = 37$$

$$5 \times 7 + 3 = 46$$

...

Decimal

$$10 = (12)_8$$

$$31 = (37)_8$$

$$38 = (46)_8$$

...



# Signed & Unsigned Integers

- Unsigned 8bit:
  - $(11111111)_2 = 255$
- Signed 8bit (only in digital circuits):
  - $127 \rightarrow '01111111'$
  - $-127 \rightarrow '11111111'$

First digit:

- 0 for positive
- 1 for negative

**1**0001111

(binary, 8bit, signed)

# Signed & Unsigned Integers

- Unsigned 8bit integer: 0 - 255
  - Signed 8bit integer: -128 - 127
- Unsigned 32bit integer: 0 - 4,294,967,295
  - Signed 32bit integer: -2,147,483,648 - 2,147,483,647
- Unless otherwise specified, treat as unsigned

# Binary Coded Decimal

- Decimal numbers, each digit represented in 4bit binary, but separately
- $185 = (0001\ 1000\ 0101)_{\text{BCD}} = (10111001)_2$
- Used in places where using decimals directly is more convenient, such as digital watches etc.

# ASCII

- American Standard Code for Information Interchange
- Assign each character with a 8bit binary code (e.g. '0'-'9', 'A'-'Z', 'a'-'z')
- The first bit is always 0

# Parity Code

- For error detection in data communication
  - e.g. resulting from packet loss or other forms of interference
- One parity bit for n-bits
  - An extra even parity bit: whether the number of 1s is not even
  - An extra odd parity: whether the number of 1s is not odd
  - Can be placed in any fixed position
  - Does it always work?

# Parity Code

Original 7bits	with Even parity	with Odd parity
1000001	<u>0</u> 1000001	<u>1</u> 1000001
1010100	<u>1</u> 1010100	<u>0</u> 1010100

# Circuits

- Circuits
  - Digital and Analog
- Integrated systems
  - Von Neumann computers
  - Embedded systems

# Number Systems

- Number systems of base N
- Binary systems
- Octal and Hexadecimal systems
- Arithmetics



# Number Systems in DC

- Bit, Byte, Representation ranges
- Signed and Unsigned Binary Integers
- BCD, ASCII, UTF8
- Parity bit

# Digital to Analog Conversion

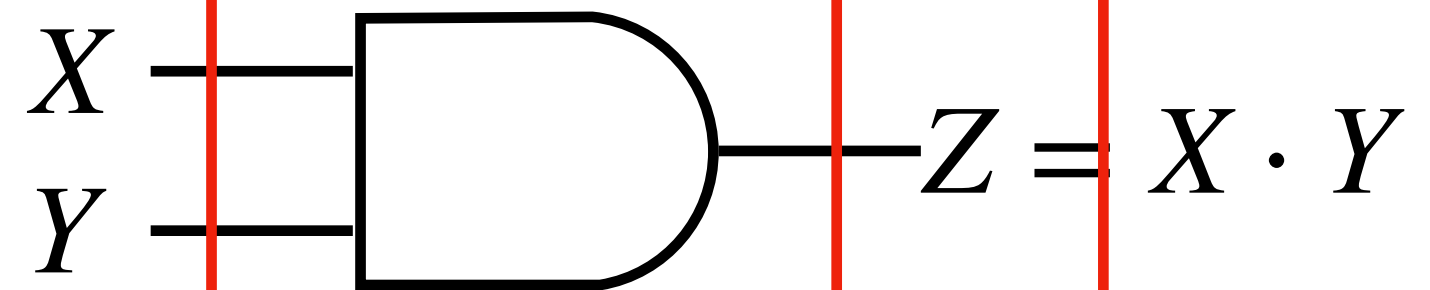
- Frequency: number of cycles per second
- Sample rate: number of samples per unit time
- Bitrate: number of bits per second

# Lecture 2: Combinational Logic Circuits

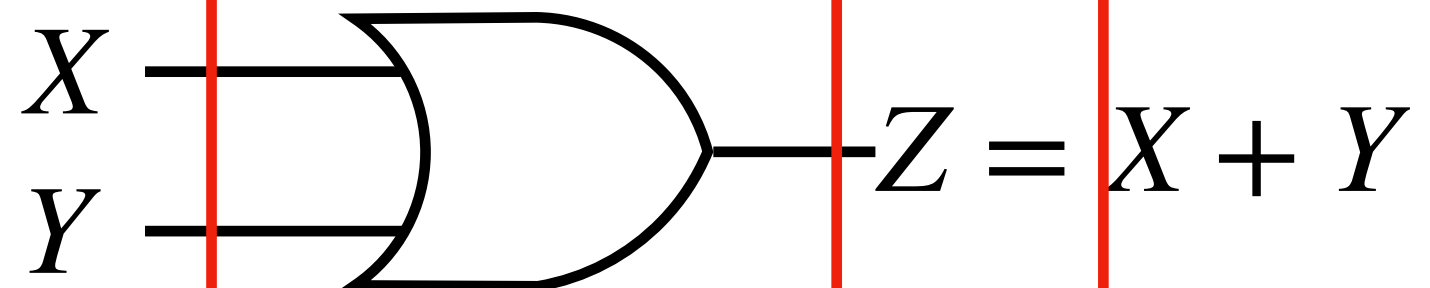
Logic Gates; Boolean Algebra; Minterm/Maxterm; K-Map; Some Other Gate Types

# First 3 Gates

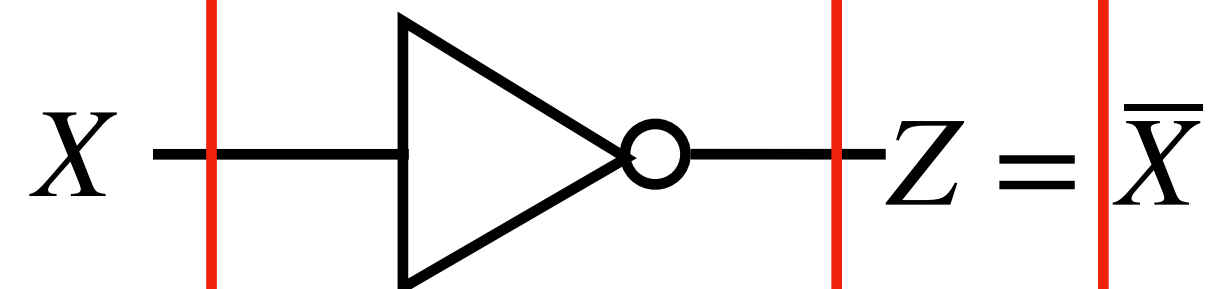
AND Gate



OR Gate



NOT Gate



Input

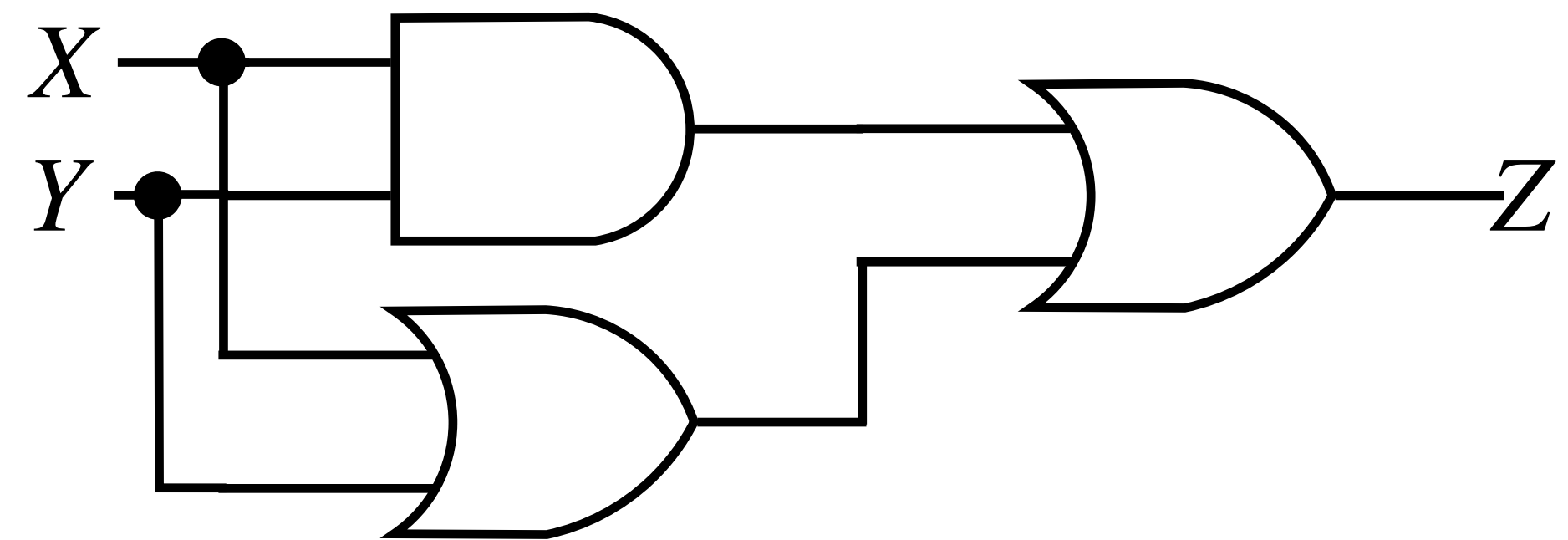
Output

Concept

# Truth Table

Truth Table

$X$	$Y$	$Z = (X \cdot Y) + (X + Y)$
0	0	0
0	1	1
1	0	1
1	1	1



# Basic Identities

- Boolean Algebra solving
  - **Identify** rules **applicable** to the expression
  - **Apply** rules that can help you **simplify** the expression
  - **Simplification**: reducing the number of variables and operators in an expression without changing its truth table values
  - **Atomic element**: an element that can't have the number of its variables and operators reduced any further

# Basic Identities

1.  $X + 0 = X$

2.  $X \cdot 1 = X$

3.  $X + 1 = 1$

4.  $X \cdot 0 = 0$

5.  $X + X = X$

6.  $X \cdot X = X$

7.  $X + \bar{X} = 1$

8.  $X \cdot \bar{X} = 0$

9.  $\bar{\bar{X}} = X$

# Basic Identities

- Commutative

$$10. X + Y = Y + X$$

$$11. XY = YX$$

- Associative

$$12. X + (Y + Z) = (X + Y) + Z$$

$$13. X(YZ) = (XY)Z$$

- Distributive

$$14. X(Y + Z) = XY + XZ$$

$$15. X + (YZ) = (X + Y)(X + Z)$$

- DeMorgan's

$$16. \overline{X + Y} = \bar{X} \cdot \bar{Y}$$

$$17. \overline{\bar{X} \cdot \bar{Y}} = \bar{X} + \bar{Y}$$



# Basic Identities

A.  $X + XY = X$

B.  $XY + X\bar{Y} = X$

C.  $X + \bar{X}Y = X + Y$

D.  $X(X + Y) = X$

E.  $(X + Y)(X + \bar{Y}) = X$

F.  $X(\bar{X} + Y) = XY$

# Complementation

- $\overline{F}$ : complement (invert) representation for a function  $F$ , obtained from an interchange of 1s to 0s and 0s to 1s for the values of  $F$  in the truth table
- Apply DeMorgan's Rule

$$16. \overline{X_1 + X_2 + \dots + X_n} = \overline{X_1} \cdot \overline{X_2} \cdot \dots \cdot \overline{X_n}$$

$$17. \overline{X_1 \cdot X_2 \cdot \dots \cdot X_n} = \overline{X_1} + \overline{X_2} + \dots + \overline{X_n}$$

# Algebraic Manipulation

Difficulty: Simple

Simplify the following expressions

- $\bar{X} \cdot \bar{Y} + XYZ + \bar{X}Y$
- $X + Y(Z + \overline{X + Z})$

# Algebraic Manipulation

Difficulty: Mid

Simplify the following expressions

- $\overline{W}X(\overline{Z} + \overline{Y}Z) + X(W + \overline{W}YZ)$
- $(AB + \overline{A}\overline{B})(\overline{C}\overline{D} + CD) + AC$

# Algebraic Manipulation

Difficulty: Mid

Simplify the following expressions

- $\bar{A} \cdot \bar{C} + \bar{A}BC + \bar{B}C$

- $\overline{A + B + C} \cdot \overline{ABC}$

# Algebraic Manipulation

Difficulty: Mid

Simplify the following expressions

- $AB\overline{C} + AC$
- $\overline{A} \cdot \overline{B}D + \overline{A} \cdot \overline{C}D + BD$

# Algebraic Manipulation

Difficulty: HARDCORE

Prove the identity of each of the following Boolean equations

- $AB\bar{C} + B\bar{C} \cdot \bar{D} + BC + \bar{C}D = B + \bar{C}D$
- $WY + \bar{W}Y\bar{Z} + WXZ + \bar{W}X\bar{Y} = WY + \bar{W}X\bar{Z} + \bar{X}Y\bar{Z} + X\bar{Y}Z$
- $A\bar{D} + \bar{A}B + \bar{C}D + \bar{B}C = (\bar{A} + \bar{B} + \bar{C} + \bar{D})(A + B + C + D)$

# Standard Forms

- Equivalent expressions can be written in a variety of ways  
**Standard forms:** typical such ways that incorporates some **unique characteristics** -> **simplify the implementation** of these designs
- **Product terms** (AND terms): e.g.  $\bar{X}YZ$   
Literals with inverts connected through only AND operators
- **Sum terms** (OR terms): e.g.  $X + \bar{Y} + Z$   
Literals with inverts connected through only OR operators



# Minterms and Maxterms

- Minterm**

**Product term;** Contains **all variables**; Has only **one Positive row** in the truth table

	X	Y	$m_0 = \bar{X}\bar{Y}$	$m_1 = \bar{X}Y$	$m_2 = X\bar{Y}$	$m_3 = XY$
(00) <sub>2</sub> =0	0	0	1	0	0	0
(01) <sub>2</sub> =1	0	1	0	1	0	0
(10) <sub>2</sub> =2	1	0	0	0	1	0
(11) <sub>2</sub> =3	1	1	0	0	0	1

# Minterms and Maxterms

- Maxterm**

**Sum term**; Contains **all variables**; Has only **one Negative row** in the truth table

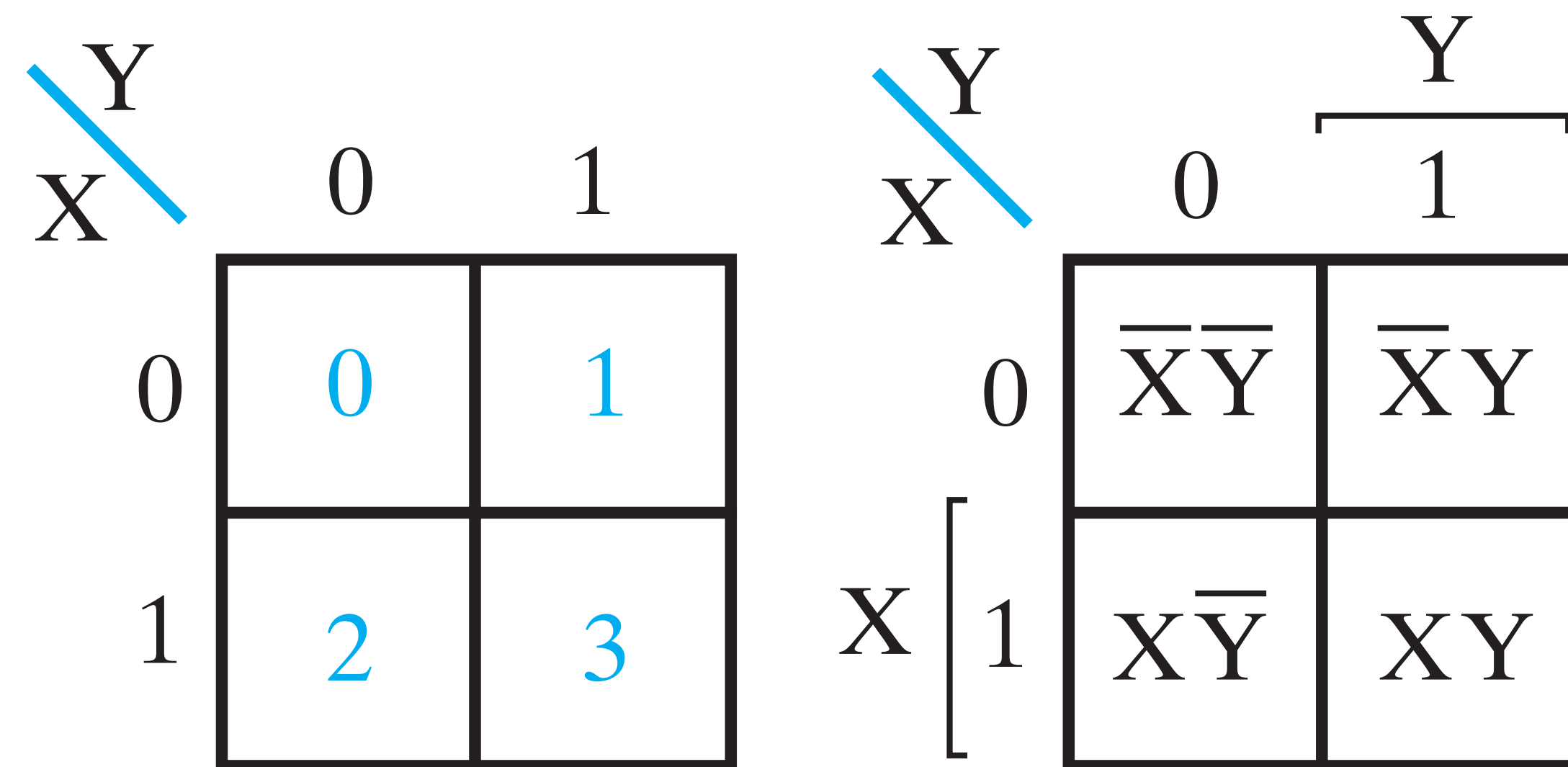
$$M_i = \overline{m_i}$$

X	Y	$M_0 = X + Y$	$M_1 = X + \bar{Y}$	$M_2 = \bar{X} + Y$	$M_3 = \bar{X} + \bar{Y}$
0	0	<b>0</b>	1	1	1
0	1	1	<b>0</b>	1	1
1	0	1	1	<b>0</b>	1
1	1	1	1	1	<b>0</b>

# Minterms and Maxterms

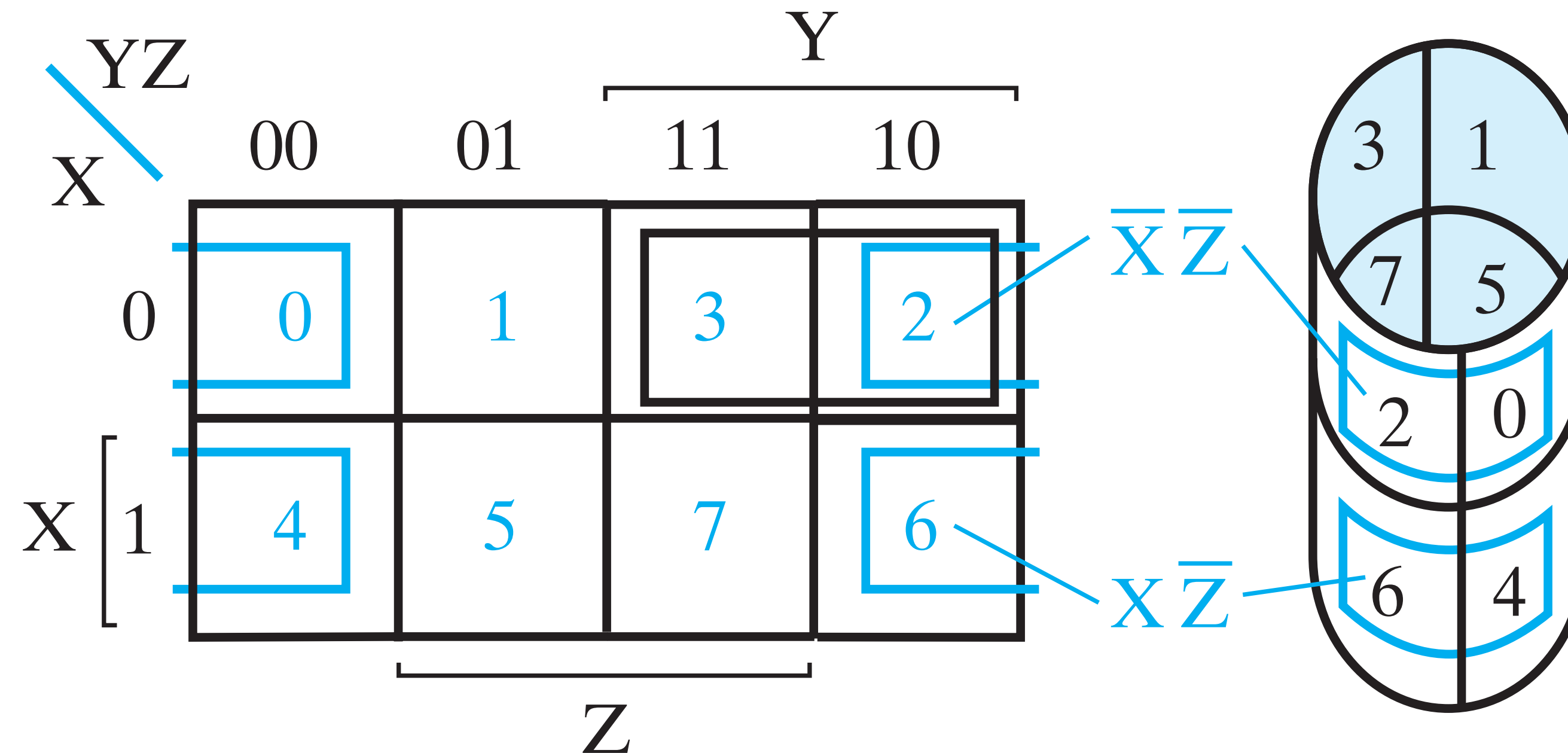
- e.g.  $M_3 = X + \bar{Y} + \bar{Z} = \overline{\bar{X}Y\bar{Z}} = \overline{m_3}$
- Sum of Minterms
  - e.g.  $F = \bar{X}\bar{Y}\bar{Z} + \bar{X}Y\bar{Z} + X\bar{Y}Z + XYZ = m_0 + m_2 + m_5 + m_7$   
 $= \Sigma m(0,2,5,7)$
- Product of Maxterm
  - e.g.  $F = (X + Y + Z)(X + \bar{Y} + Z)(\bar{X} + Y + \bar{Z})(\bar{X} + \bar{Y} + \bar{Z})$   
 $= M_0M_2M_5M_7$   
 $= \Pi M(0,2,5,7)$

# Two Variable Maps



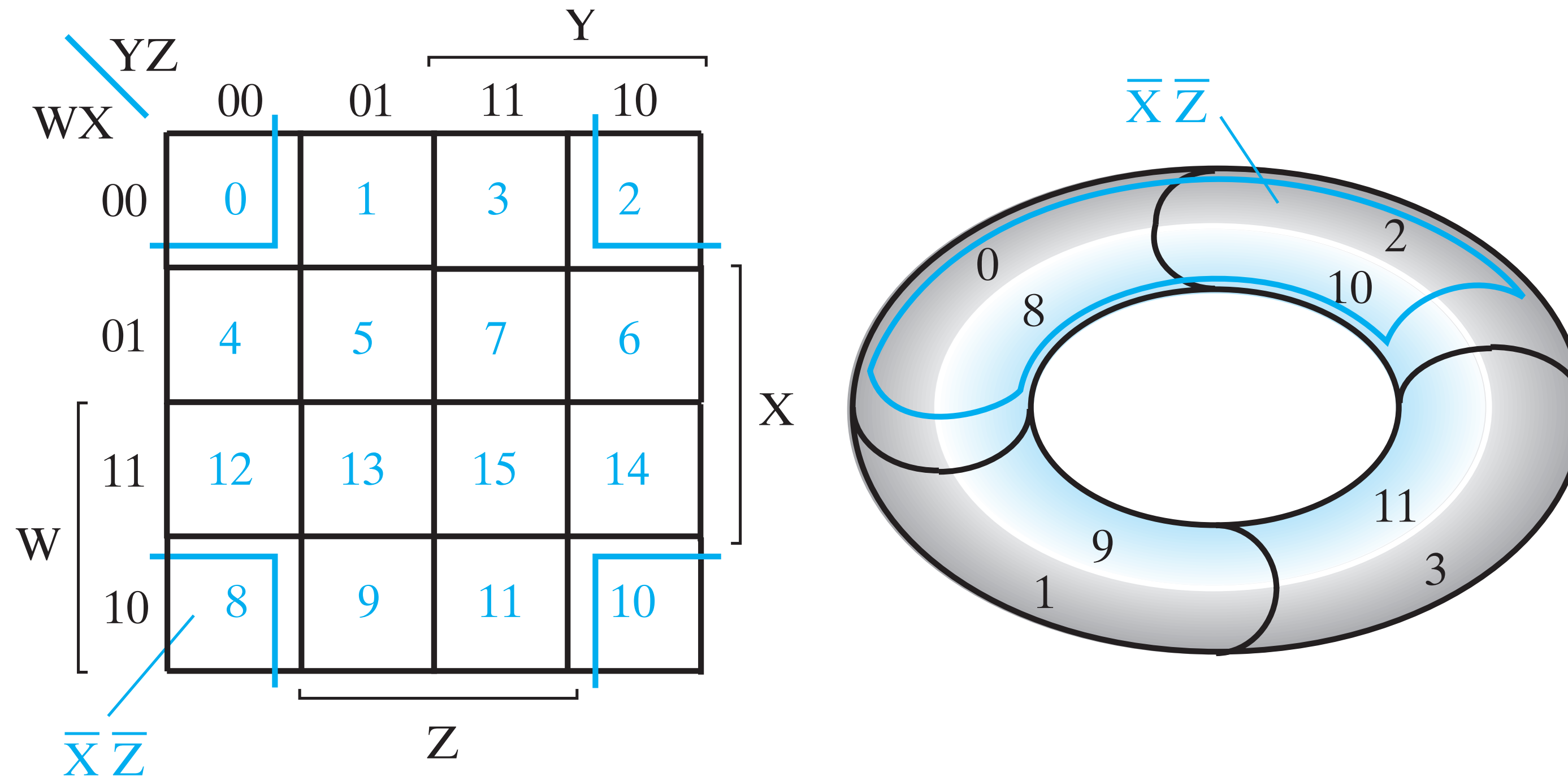
- Number of squares in each map is equal to the number of minterms for the same number of variables, light blue digit above is the index (of minterm)
- Two squares are adjacent if they only differ in one variable
- Binary value inside at each position indicates the truth table value for that term

# Three Variable Maps



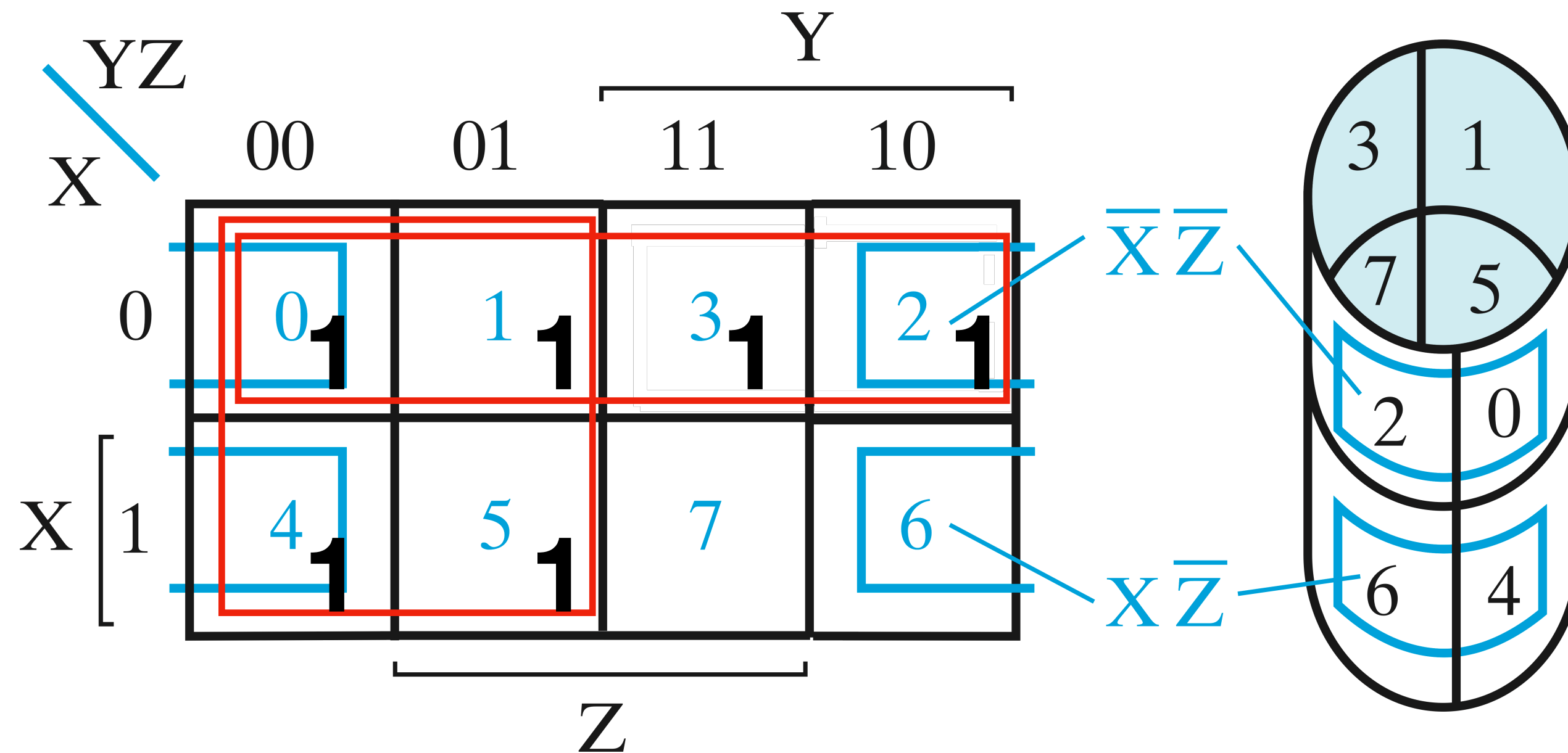
- Number of squares in each map is equal to the number of minterms for the same number of variables, light blue digit above is the index (of minterm)
- Two squares are adjacent if they only differ in one variable
- Binary value inside at each position indicates the truth table value for that term

# Four Variable Maps



- Number of squares in each map is equal to the number of minterms for the same number of variables, light blue digit above is the index (of minterm)
- Two squares are adjacent if they only differ in one variable
- Binary value inside at each position indicates the truth table value for that term

# K Map Optimisation



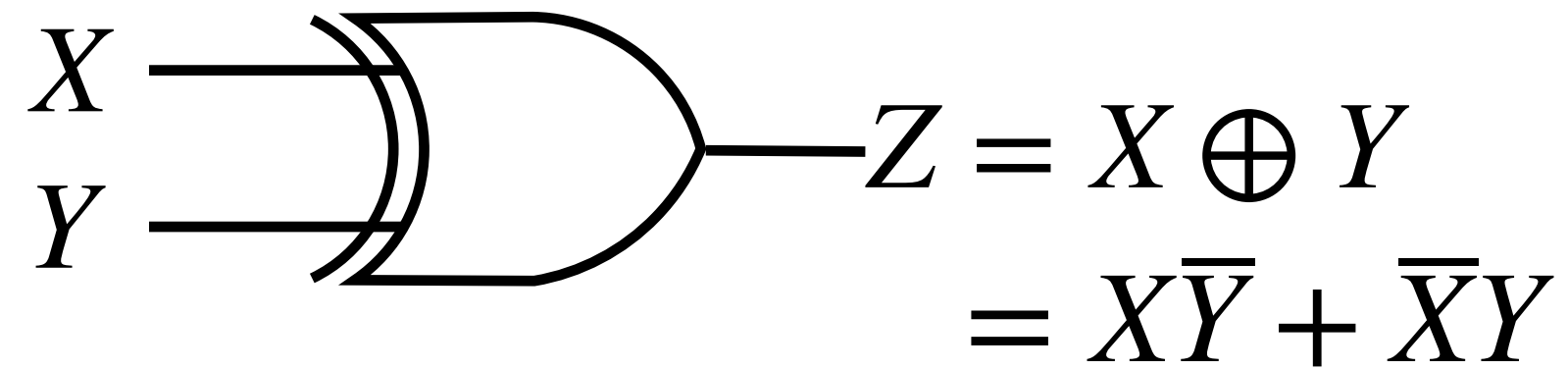
$$F(X, Y, Z) = \sum m(0, 1, 2, 3, 4, 5)$$

$$= \bar{X} + \bar{Y}$$

- Step 1: Enter the values
- Step 2: Identify the set of **largest** rectangles in which **all values are 1**, covering **all 1s**
- Step 3: **Read off** the selected rectangles. If rectangle has odd length edges (excluding 1), split

# XOR Gate

**XOR Gate**  
**Exclusive-OR**



- $X \oplus 0 = X$
- $X \oplus 1 = \bar{X}$
- $X \oplus X = 0$
- $X \oplus \bar{X} = 1$
- $X \oplus \bar{Y} = \overline{X \oplus Y}$
- $\bar{X} \oplus Y = \overline{X \oplus Y}$

XOR Truth Table

$X$	$Y$	$Z = X \oplus Y$
0	0	0
0	1	1
1	0	1
1	1	0



# XOR Gate

- $X \oplus 0 = X$

- $X \oplus X = 0$

- $X \oplus \bar{Y} = \overline{X \oplus Y}$

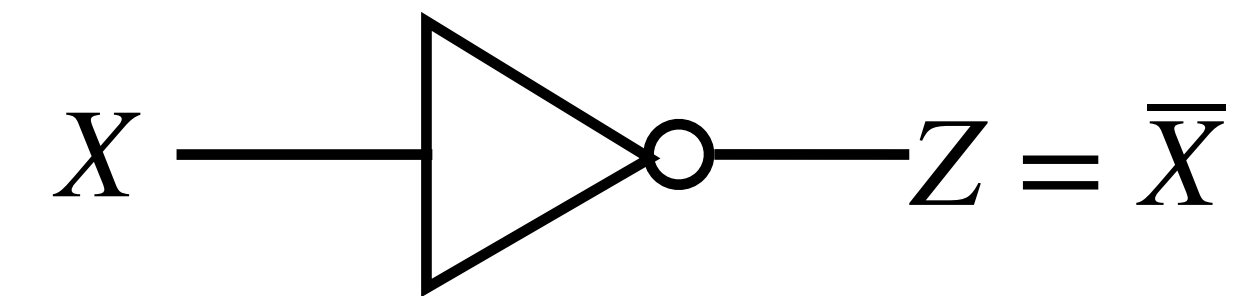
- $X \oplus 1 = \bar{X}$

- $X \oplus \bar{X} = 1$

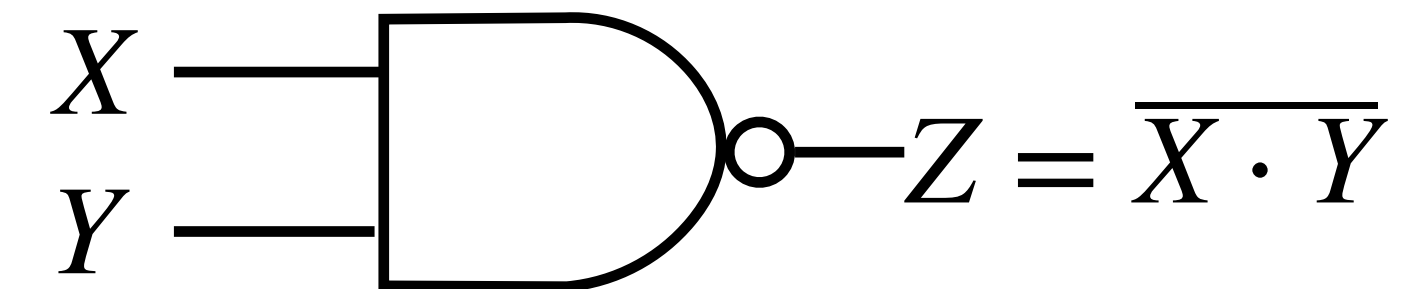
- $\bar{X} \oplus Y = \overline{X \oplus Y}$

# N-Gates

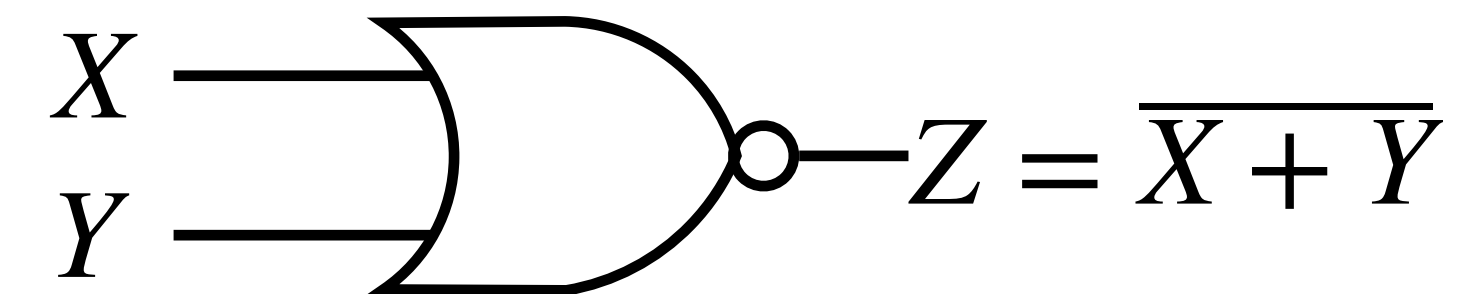
**NOT Gate**



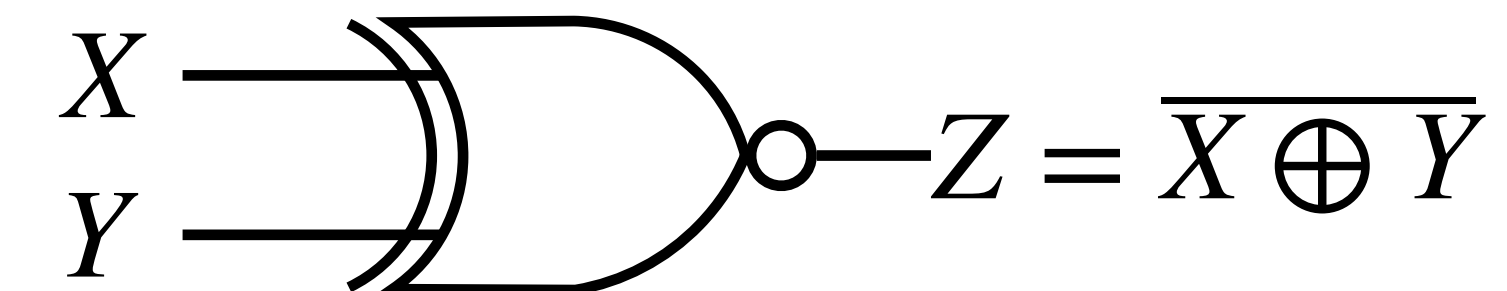
**NAND Gate**



**NOR Gate**



**XNOR Gate**



# Boolean Algebra

## I. AND, OR, NOT Operators and Gates

- Simple digital circuit implementation
- Algebraic manipulation using Binary Identities

## II. Standard Forms

- Minterm & Maxterm
- Sum of Products & Product of Sums

## III. Optimisation Using K-Map (For 2,3,4 Variables)

## IV. XOR, NAND, NOR, XNOR

# Lecture 3: Combinational Logic Design

5 Steps Systematic Design Procedures; Functional  
Blocks; Decoder, Enabler, Multiplexer; Arithmetic Blocks

# Systematic Design Procedures

1. **Specification:** Write a specification for the circuit
2. **Formulation:** Derive relationship between inputs and outputs of the system  
e.g. using truth table or Boolean expressions
3. **Optimisation:** Apply optimisation, minimise the number of logic gates and literals required
4. **Technology Mapping:** Transform design to new diagram using available implementation technology
5. **Verification:** Verify the correctness of the final design in meeting the specifications

# Hierarchical Design

- "divide-and-conquer"
- Circuit is broken up into individual functional pieces (blocks)
  - Each block has explicitly defined **Interface** (I/O) and **Behaviour**
  - A single block can be **reused** multiple times to simplify design process
  - If a single block is too complex, it can be **further divided into smaller blocks**, to allow for easier designs

# Value-Fixing, Transferring, and Inverting

- ① **Value-Fixing:** giving a constant value to a wire
  - $F = 0; F = 1;$
- ② **Transferring:** giving a variable (wire) value from another variable (wire)
  - $F = X;$
- ③ **Inverting:** inverting the value of a variable
  - $F = \bar{X}$

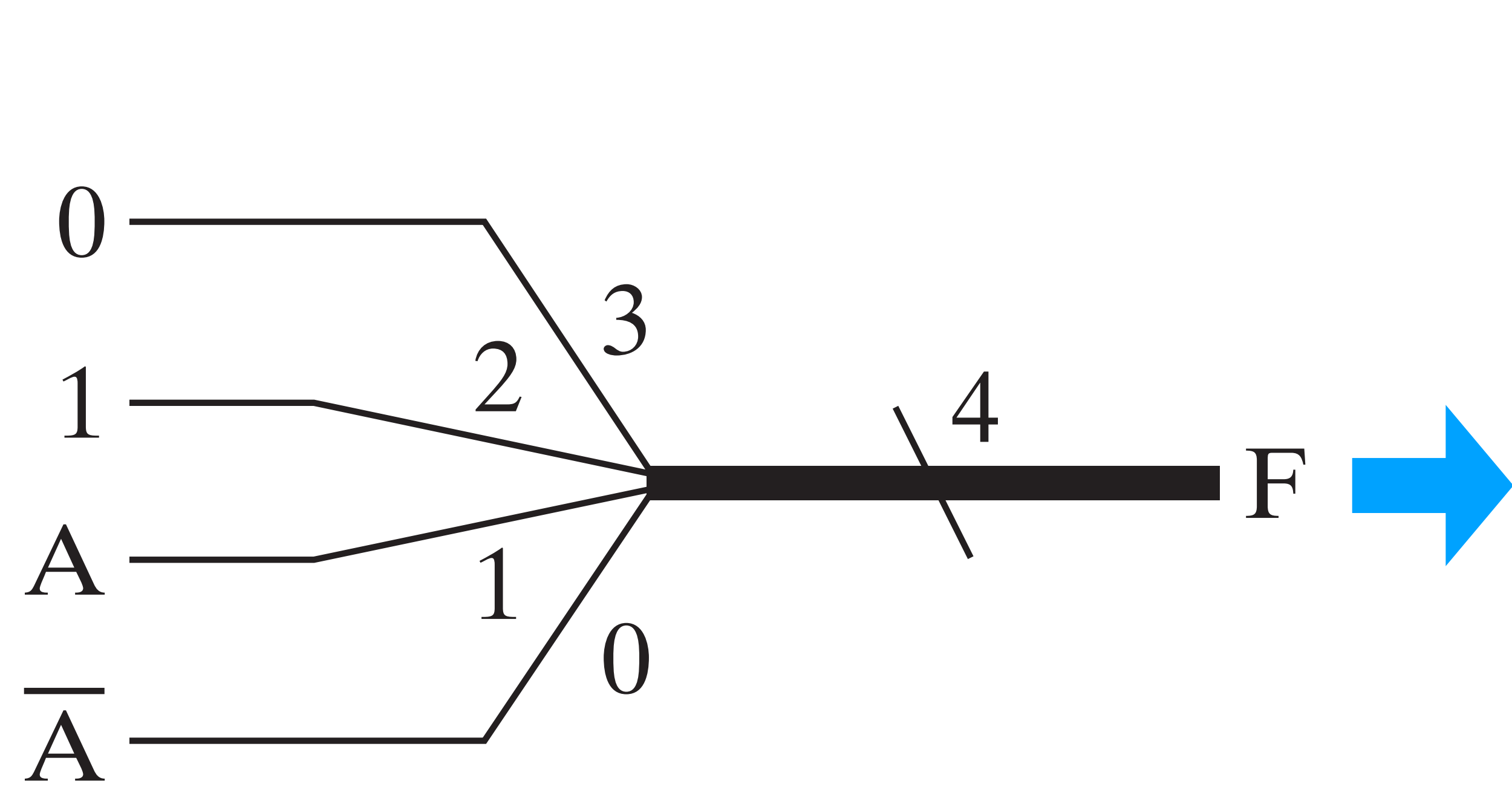
# Vector Denotation

## ④ Multiple-bit Function

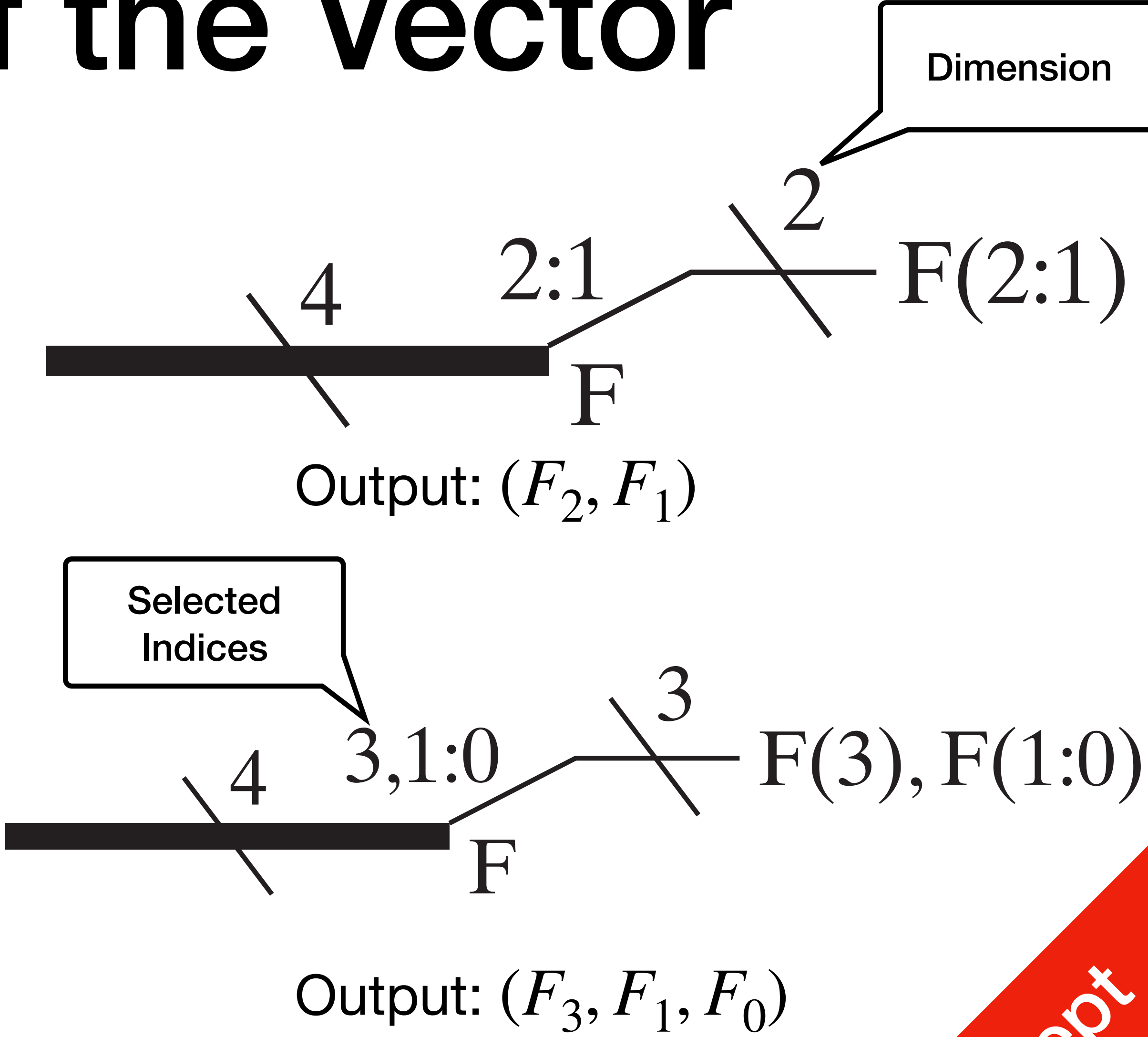
- Functions we've seen so far has only one-bit output: 0/1
- Certain functions may have  $n$ -bit output
- $F(n - 1 : 0) = (F_{n-1}, F_{n-2}, \dots, F_0)$ , each  $F_i$  is a one-bit function
- Curtain Motor Control Circuit:  $F = (F_{\text{Motor}_1}, F_{\text{Motor}_2}, F_{\text{Light}})$



# Taking part of the Vector



④ Multiple-bit Function



Concept

# Enabler

## ⑤ Enabler

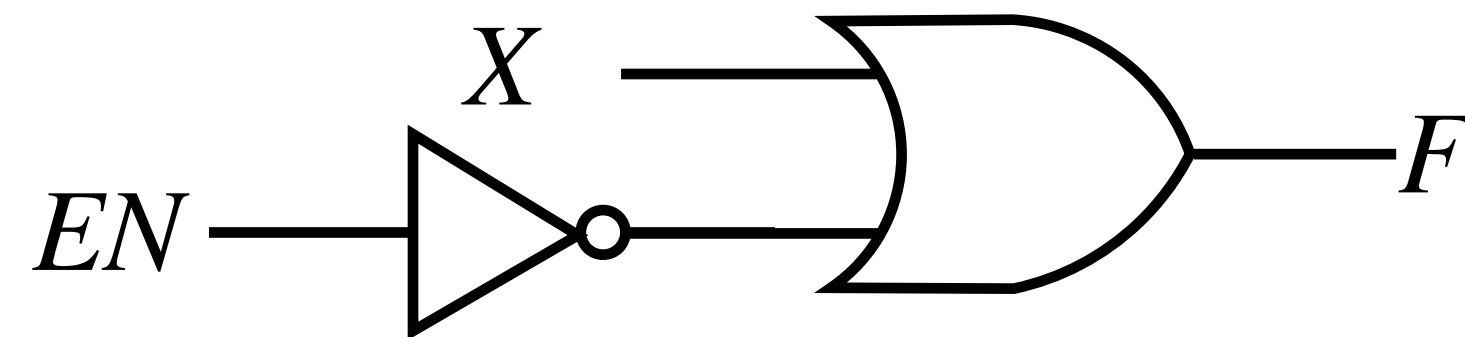
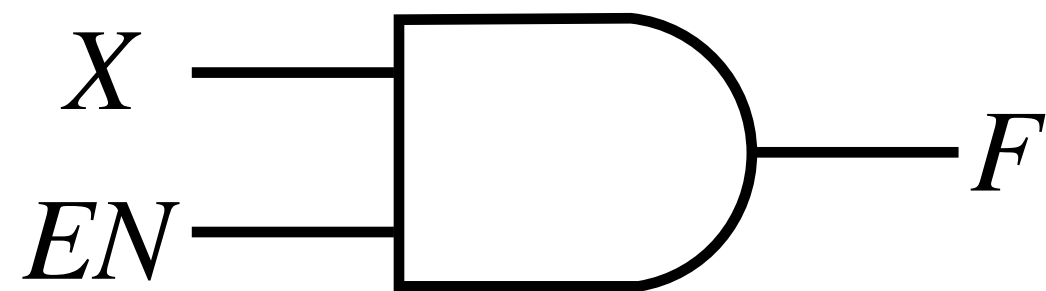
- Transferring function, but with an additional  $EN$  signal acting as switch

EN	X	F
0	X	0
1	0	0
1	1	1

# Enabler

## ⑤ Enabler

- Transferring function, but with an additional  $EN$  signal acting as switch



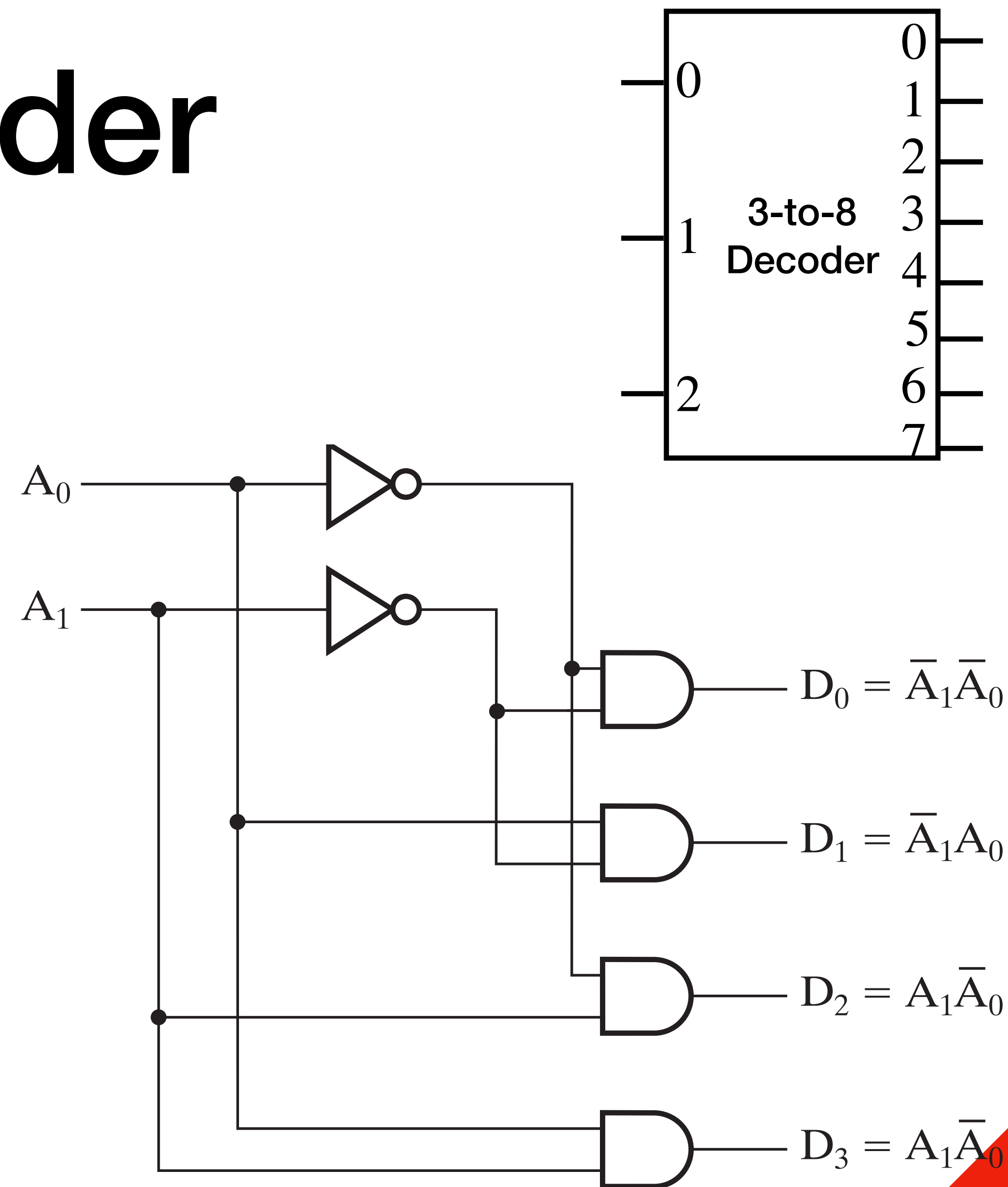
# Decoder

- $n$ -bit input,  $2^n$ bits output

- $D_i = m_i$

- Design: use hierarchical designs!

A <sub>1</sub>	A <sub>0</sub>	D <sub>0</sub>	D <sub>1</sub>	D <sub>2</sub>	D <sub>3</sub>
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1

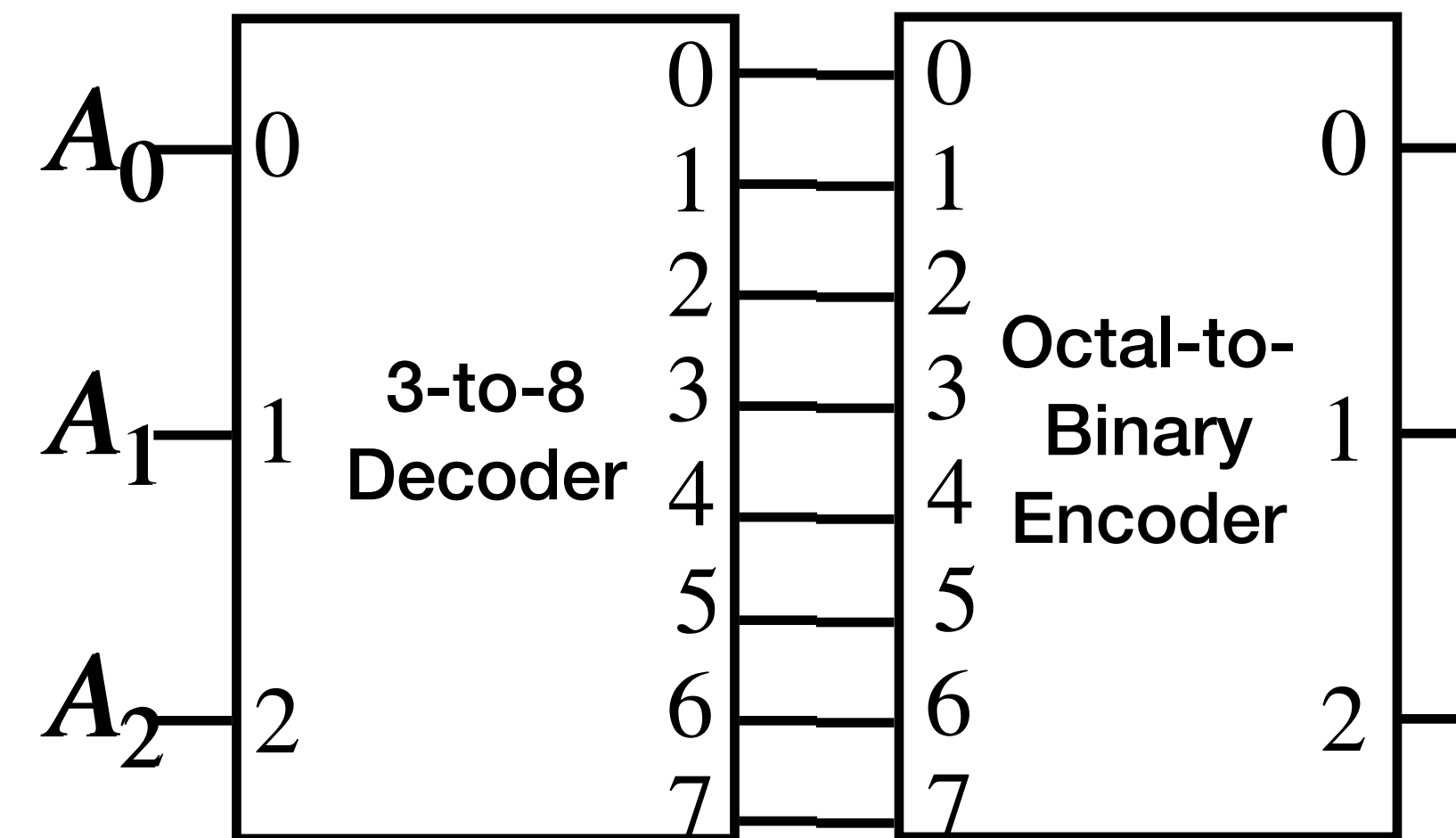


$A_1 \bar{A}_0$

Concept

# Encoder

- Inverse operation of a decoder
- $2^n$  inputs, only one is giving positive input<sup>1</sup>
- $n$  outputs



1. In reality, could be less

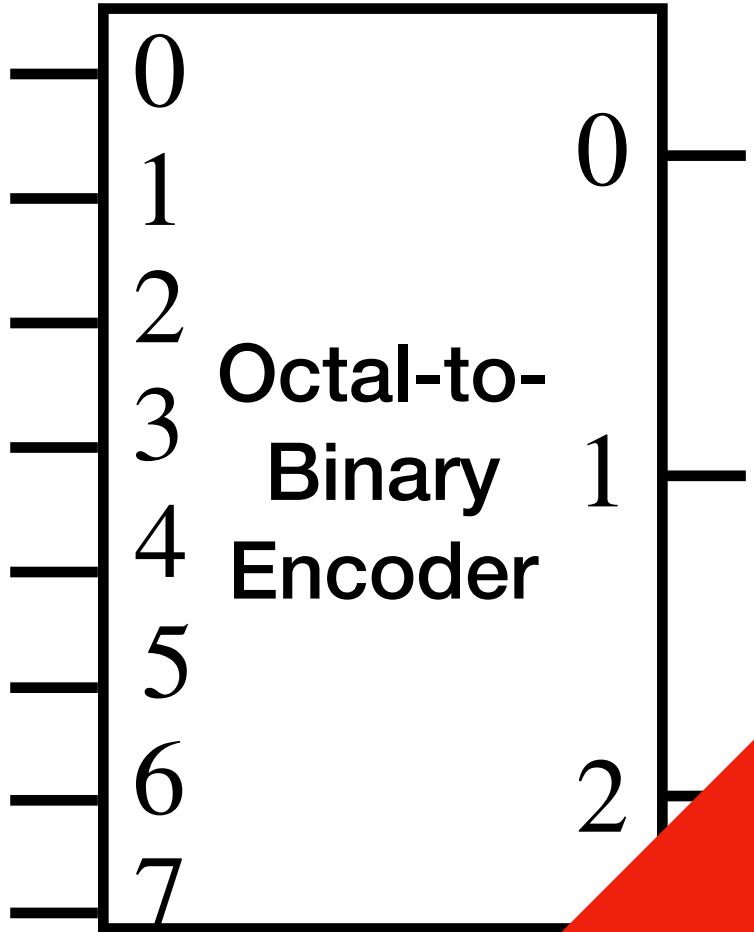
# Encoder

D <sub>7</sub>	D <sub>6</sub>	D <sub>5</sub>	D <sub>4</sub>	D <sub>3</sub>	D <sub>2</sub>	D <sub>1</sub>	D <sub>0</sub>	A <sub>2</sub>	A <sub>1</sub>	A <sub>0</sub>
							1	0	0	0
						1		0	0	1
					1			0	1	0
				1				0	1	1
			1					1	0	0
		1						1	0	1
	1							1	1	0
1								1	1	1

$A_0 = D_1 + D_3 + D_5 + D_7$

$A_1 = D_2 + D_3 + D_6 + D_7$

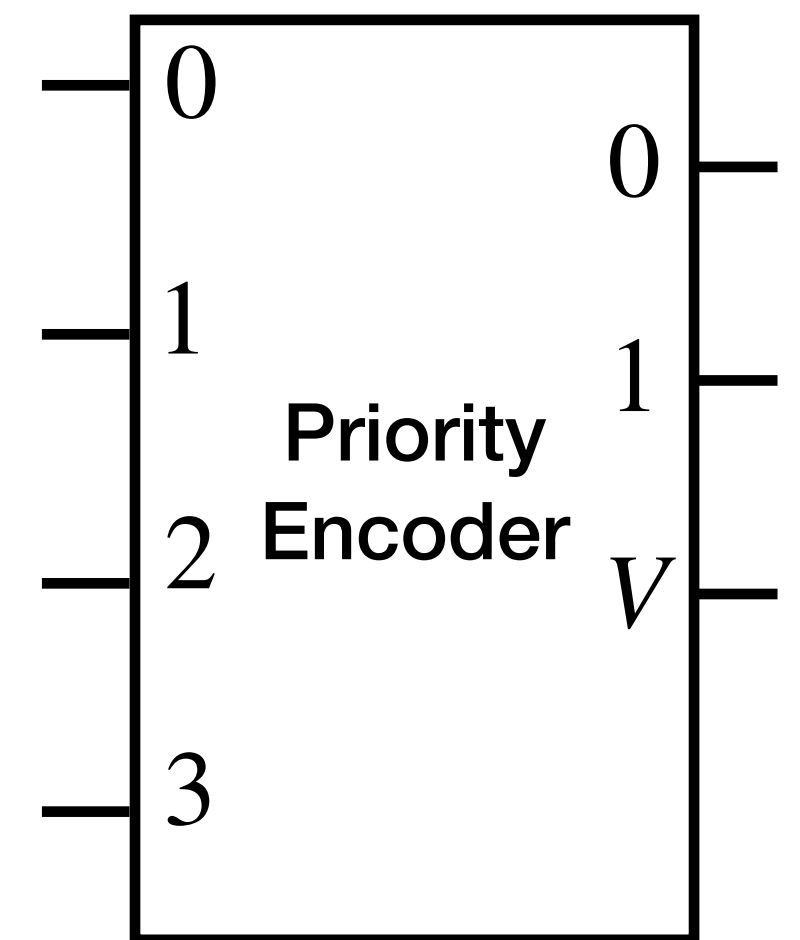
$A_2 = D_4 + D_5 + D_6 + D_7$



Concept

# Priority Encoder

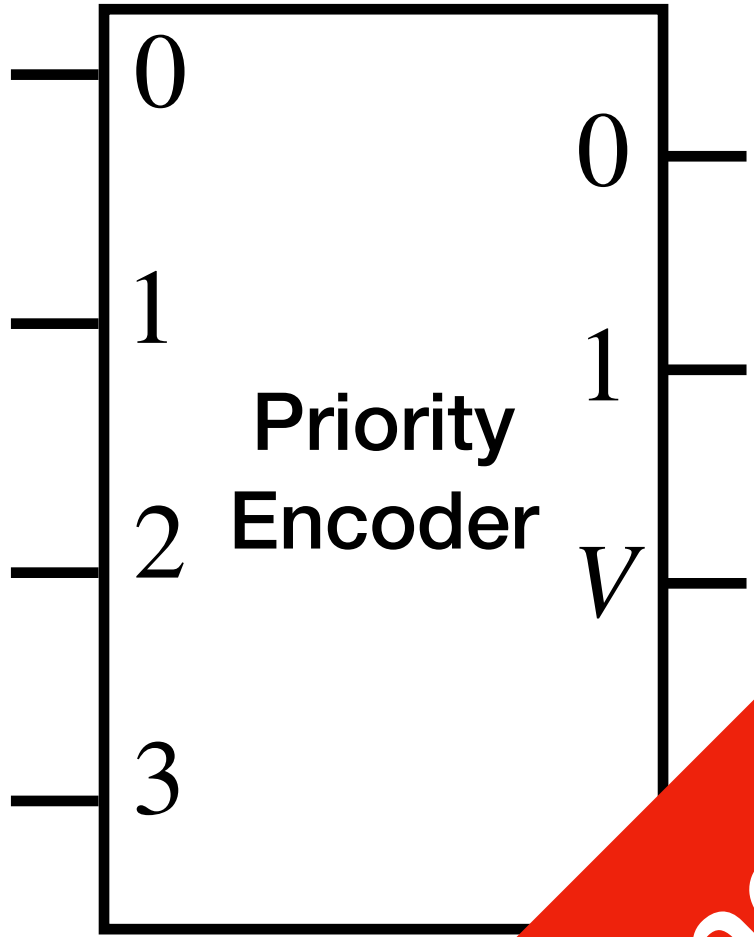
- Additional Validity Output  $V$ 
  - Indicating whether the input is valid (contains 1)
- Priority
  - Ignores  $D_{<i}$  if  $D_i = 1$



# Priority Encoder

D <sub>3</sub>	D <sub>2</sub>	D <sub>1</sub>	D <sub>0</sub>	A <sub>1</sub>	A <sub>0</sub>	V
0	0	0	0	0	0	0
0	0	0	1	0	0	1
0	0	1	X	0	1	1
0	1	X	X	1	0	1
1	X	X	X	1	1	1

$$V = D_3 + D_2 + D_1 + D_0$$
$$A_1 = D_3 + \overline{D_3}D_2 = D_2 + D_3$$
$$A_0 = \overline{D_3}\overline{D_2}D_1 + D_3$$
$$= \overline{D_2}D_1 + D_3$$



Concept



# Multiplexer

- Multiple  $n$ -variable input vectors
- Single  $n$ -variable output vector
- Switches: which input vectors to output

