

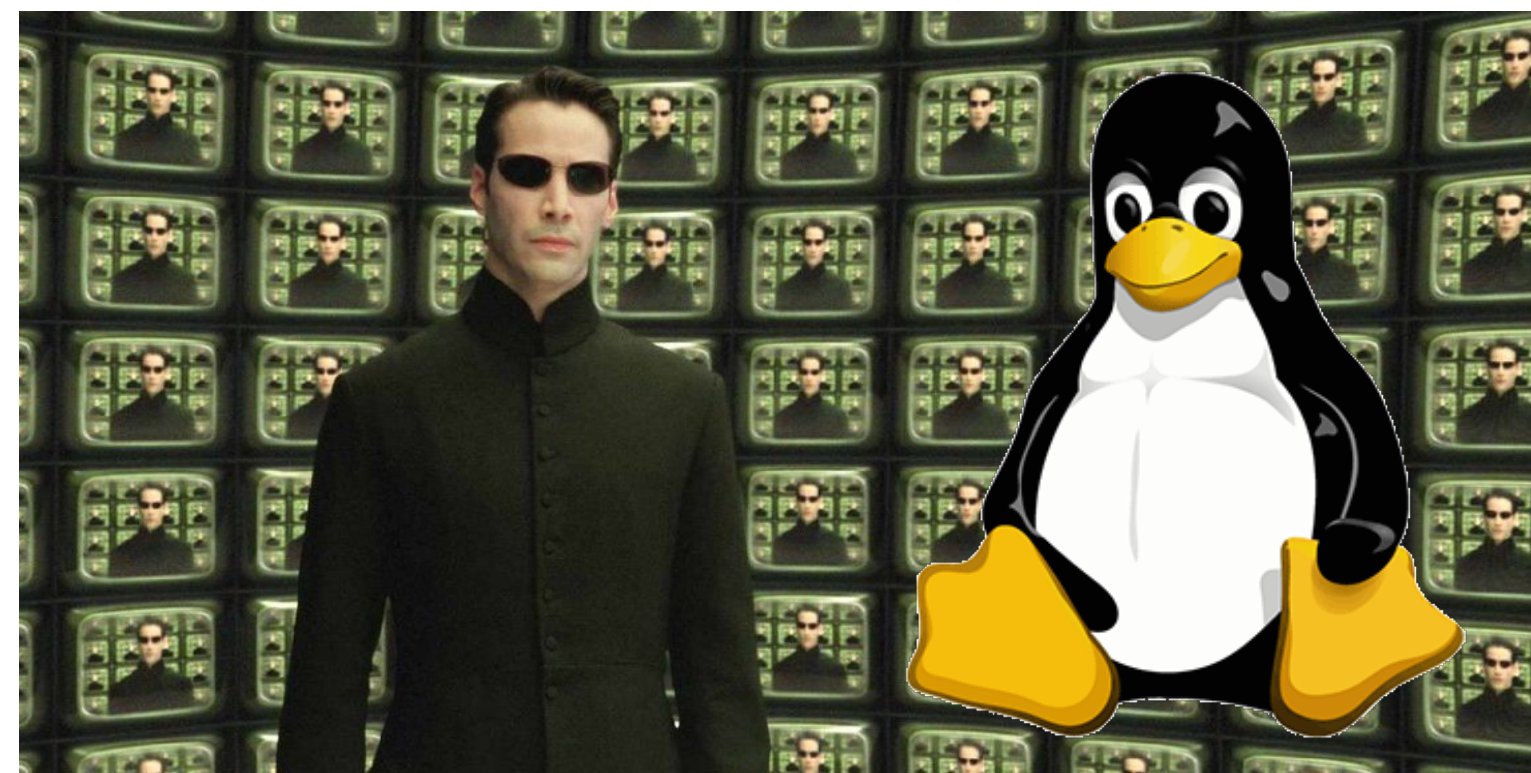


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CSCI 125

Introduction to Computer Science and Programming II

Lecture 7: Data Structure IV



Jetic Gū
2020 Summer Semester (S2)

Some changes

- Assignment 4 and Lab 4 due 9 Aug, covering Lecture 7
- No class next Monday (BC day)
- Last batch of OJ problem: 6 problems p025-p030

Overview

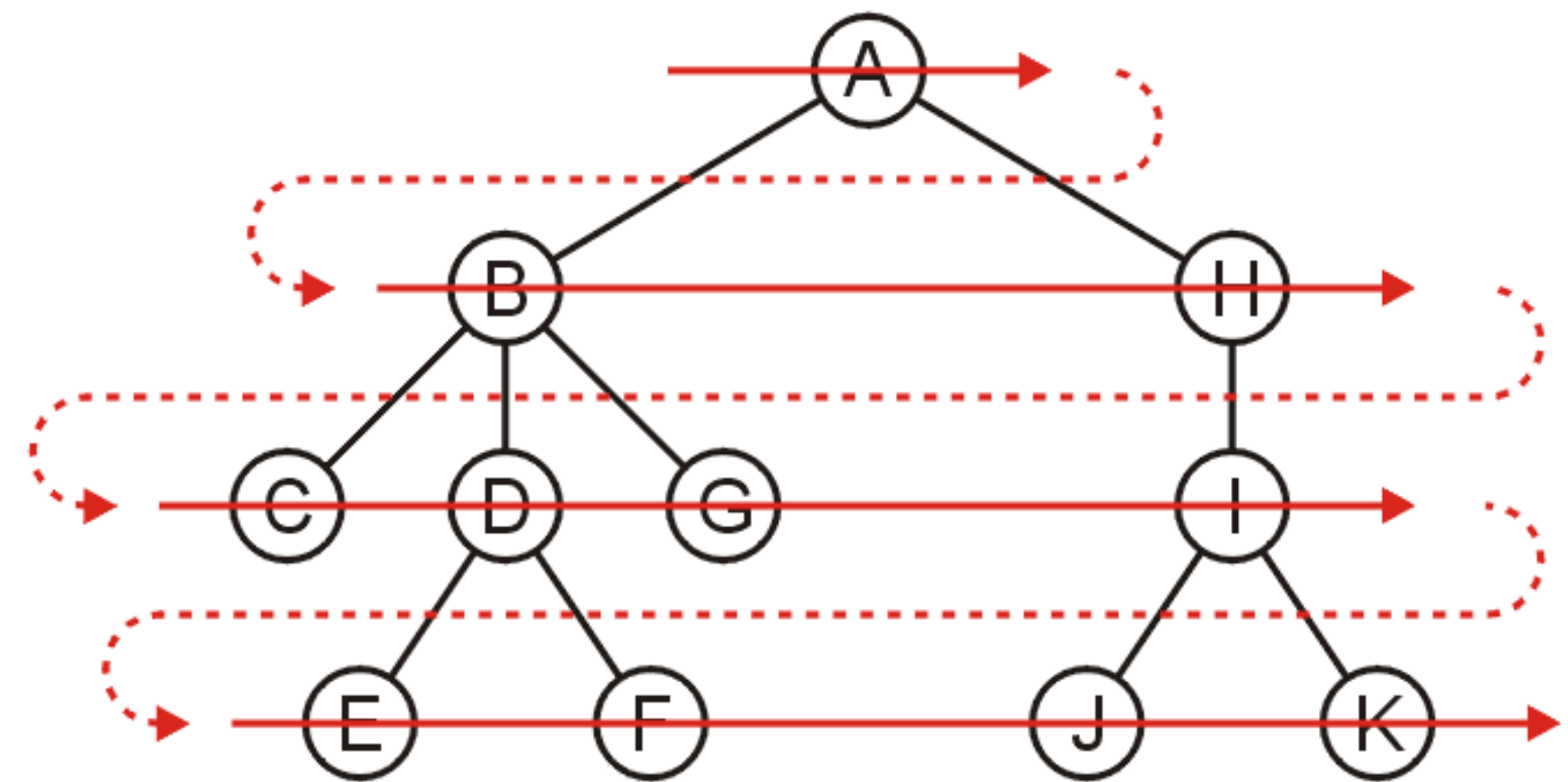
- Focus: Data Structures
- Architecture: Linux/Unix OS
- Core Ideas:
 1. Breath-First Search VS Depth-First Search

Search Algorithms

Breath-First Search VS Depth-First Search

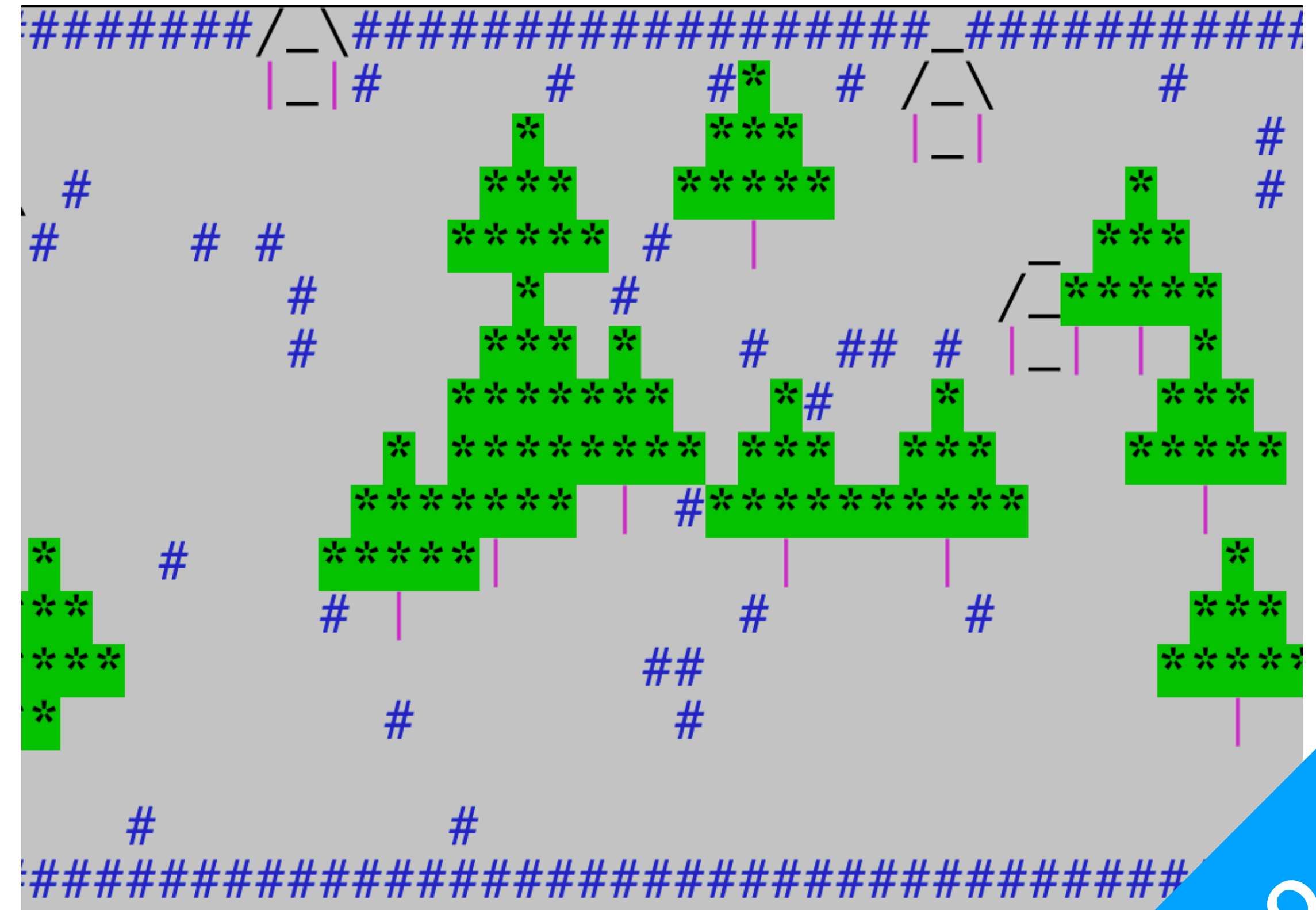
Breadth-First Search

- Algorithm for traversing or searching
- Visit all nodes on the same depth, before moving on
- Implementation using Queue
- Works on trees and graphs



BFS Example

- Village of the Sorcerer
- Certain areas of the map are not reachable (fences, houses)
- Certain areas might be entirely blocked off by fences



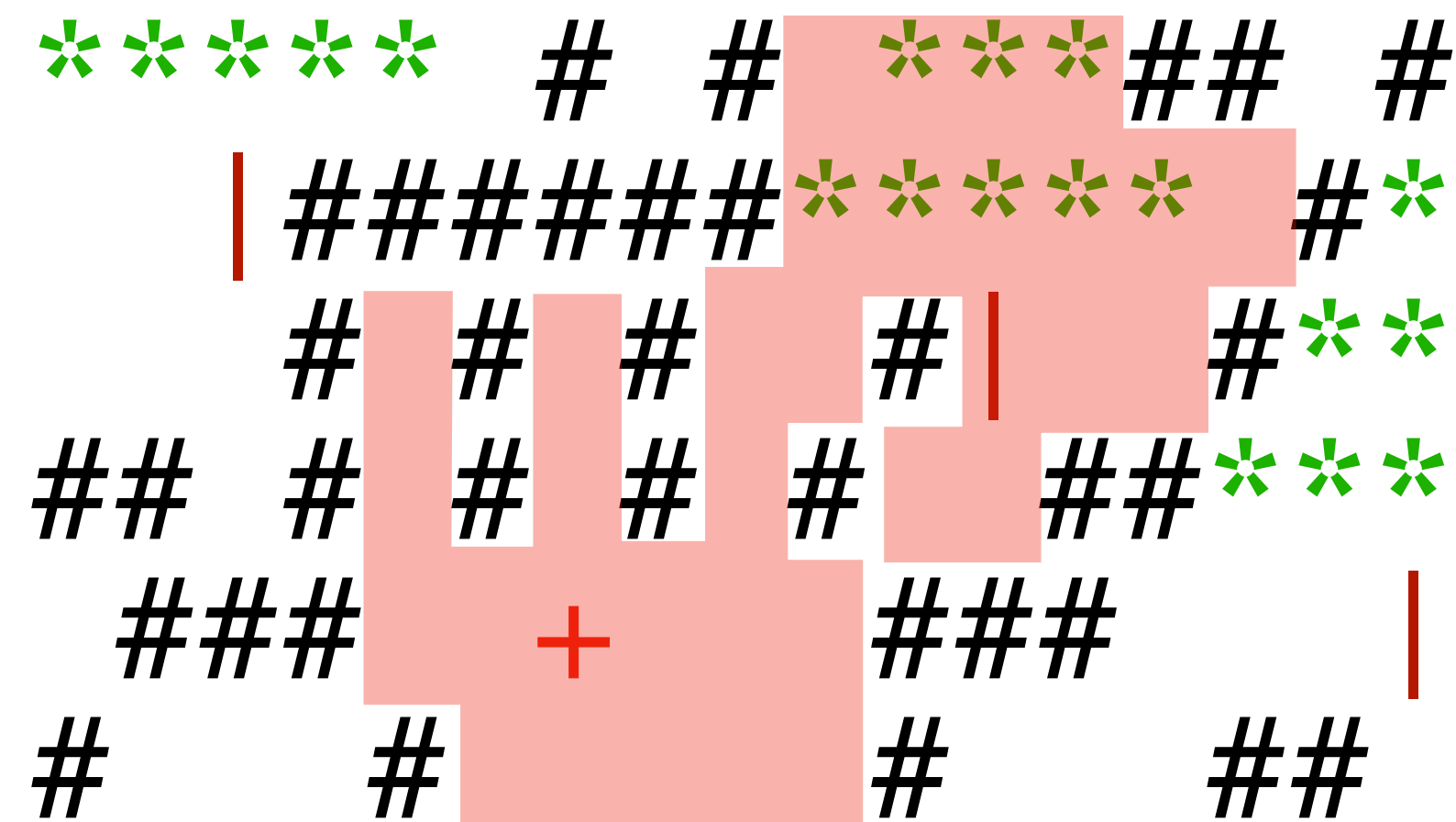
BFS Example

- Village of the Sorcerer
- Certain areas of the map are not reachable (fences, houses)
- Certain areas might be entirely blocked off by fences
- **+**: player

```
***** # # *****## #  
| #####***** #*  
# # # # | #**  
## # # # ##***  
### + ### |  
# # # ##
```


BFS Example

- Village of the Sorcerer
- Certain areas of the map are not reachable (fences, houses)
- Certain areas might be entirely blocked off by fences
- **+**: player
- **Light Red**: player accessible regions



BFS Example

- We start by initialising a Queue to store coordinates
- `push(playerCoordinate)`, which is 4,2
 - `push ((4, 2)) ;`
- In C++, you can simulate this using 2 `int`-based queues

(4, 2)
↑
Front

	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

(4, 2)
↑
Front

	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited

(4, 2)
↑
Front

	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (4, 2);`

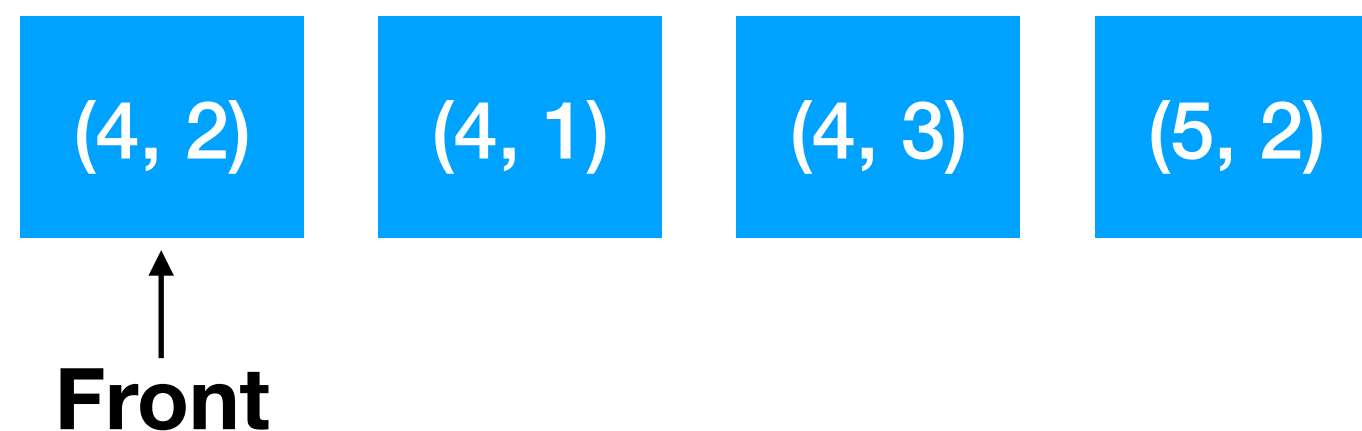
(4, 2)
↑
Front

	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (4, 2);`
- `push(4, 1); push(4, 3); push(5, 2)`

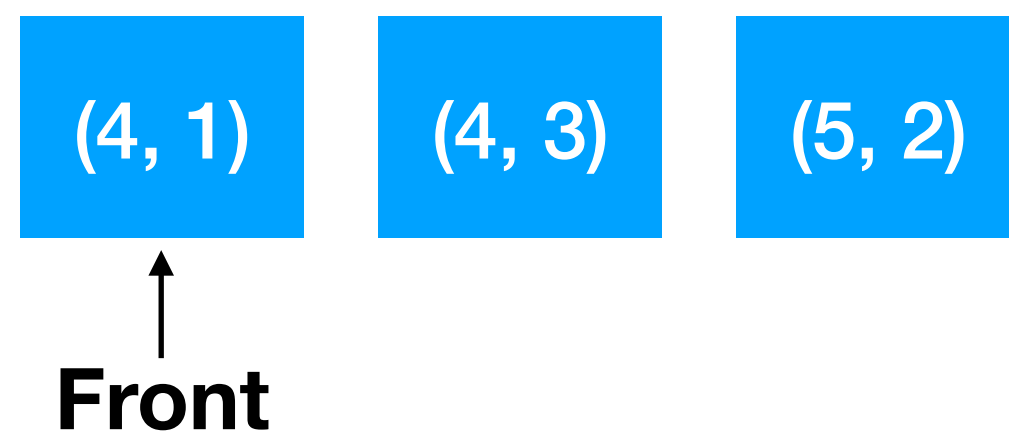


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (4, 2);`
- `push(4, 1); push(4, 3); push(5, 2)`
- `pop();`



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

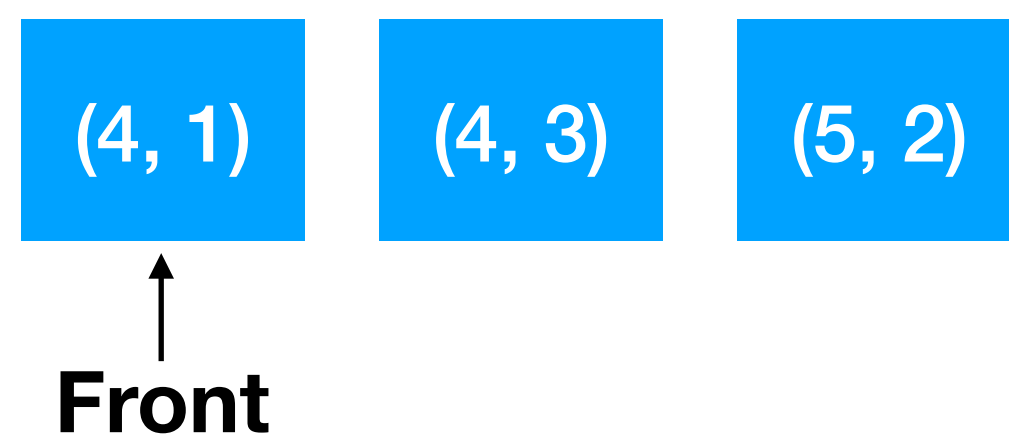


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited

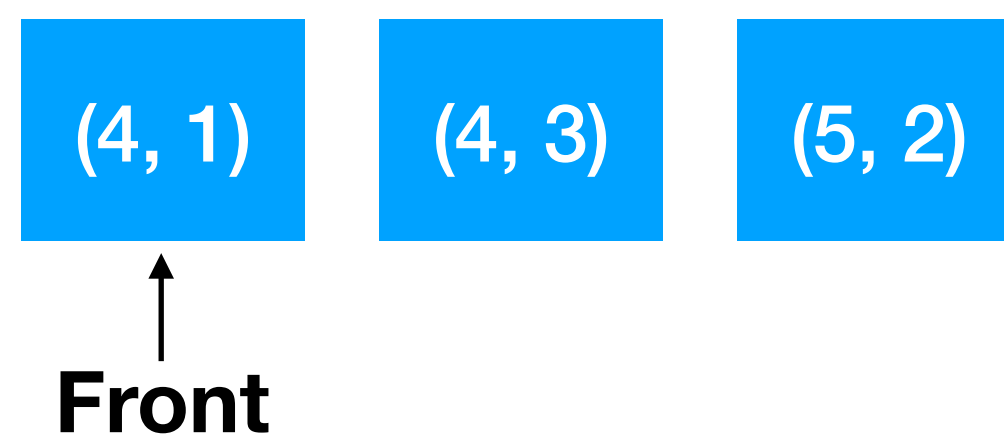


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (4, 1);`

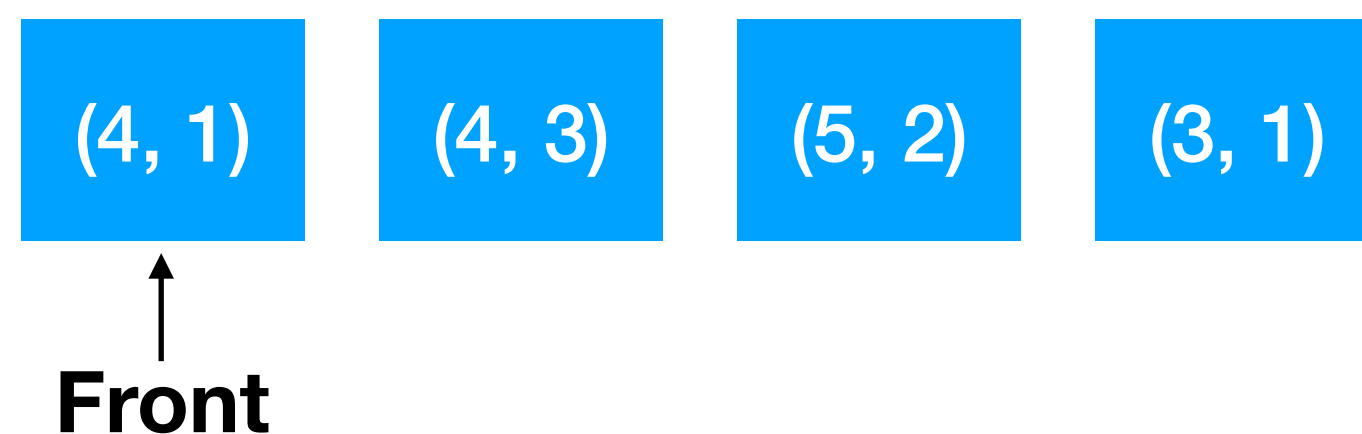


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (4, 1);`
- `push(3, 1);`

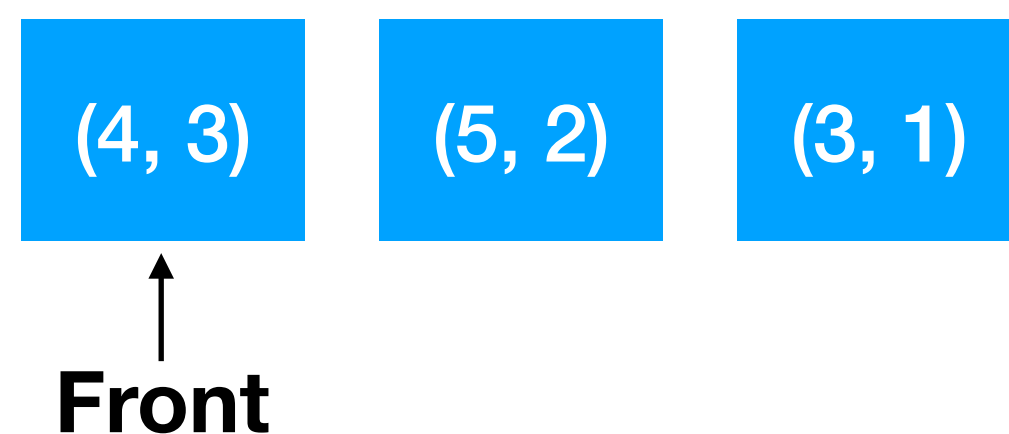


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (4, 1);`
- `push(3, 1);`
- `pop();`



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited

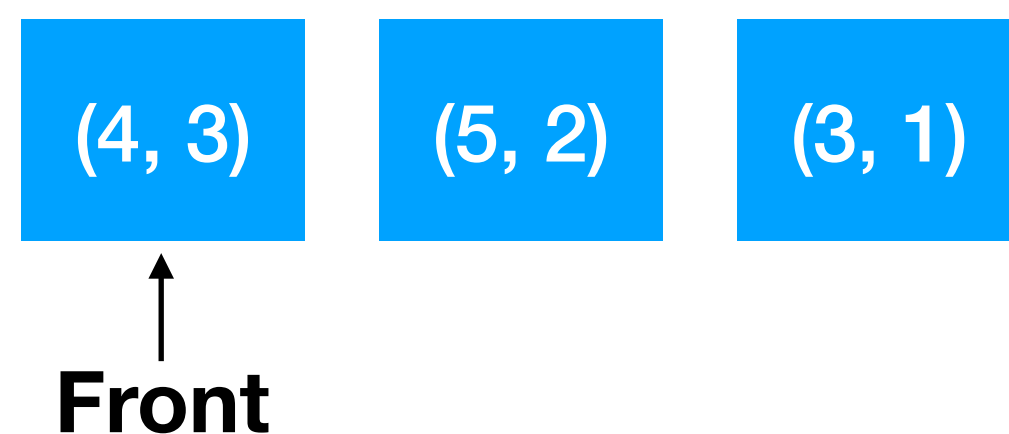


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (4, 3);`

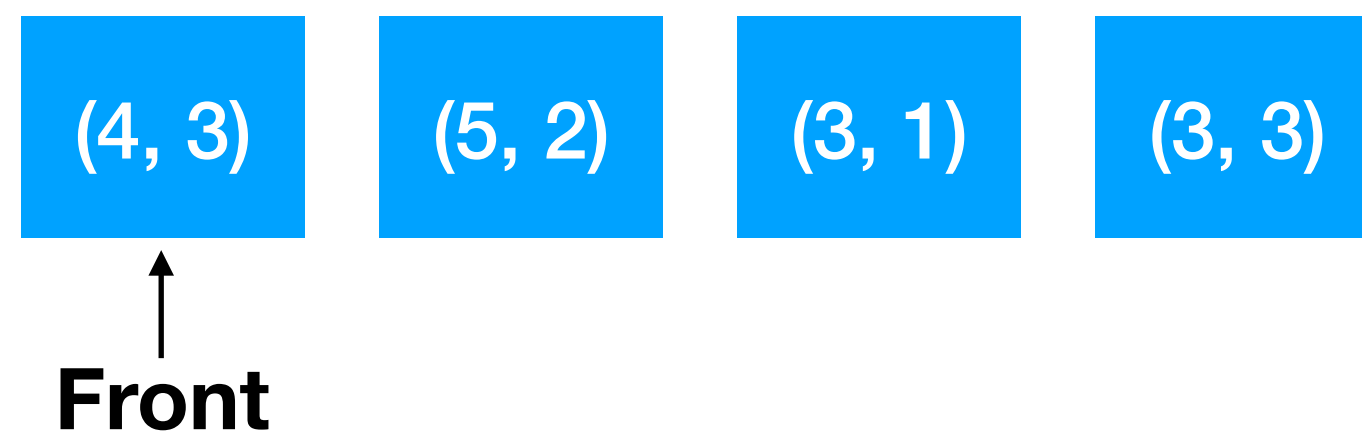


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (4, 3);`
- `push(3, 3);`



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

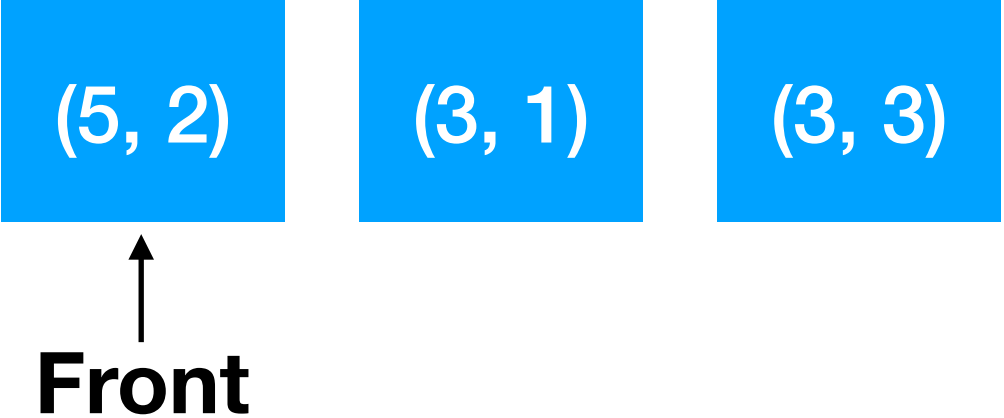
- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (4, 3);`
- `push(3, 3);`
- `pop();`



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (5, 2);`



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (5, 2);`
- no push here

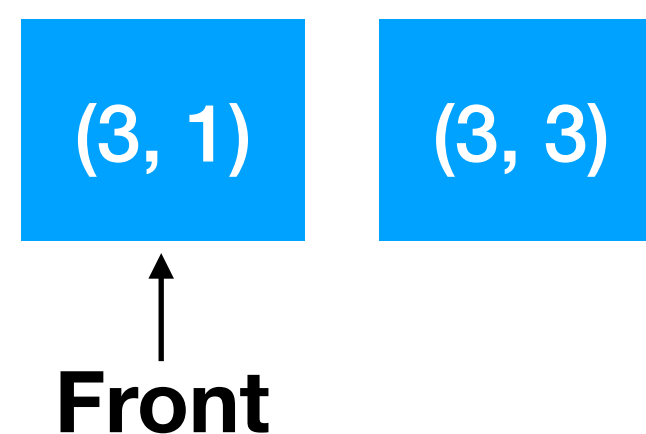


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

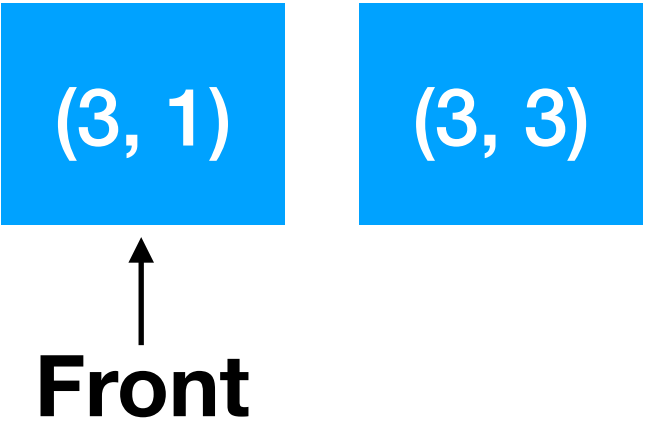
- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (5, 2);`
- no push here
- `pop();`



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

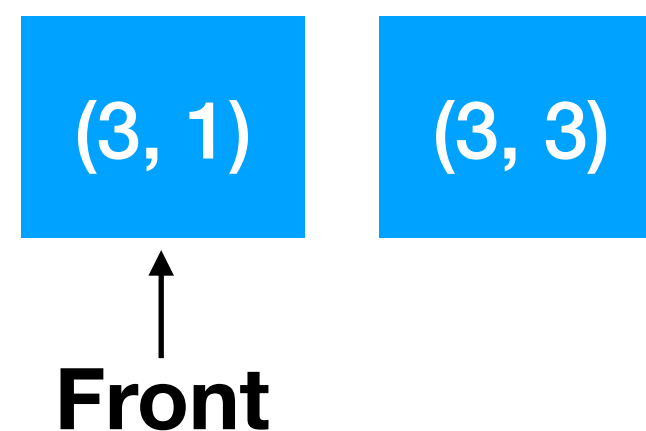


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited

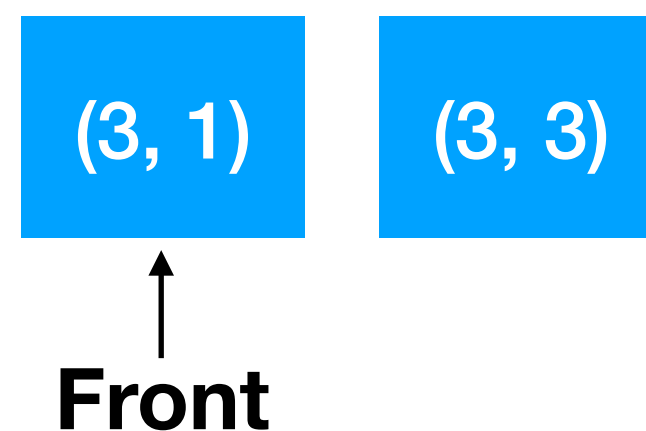


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (3, 1);`



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (3, 1);`
- `push(2, 1);`

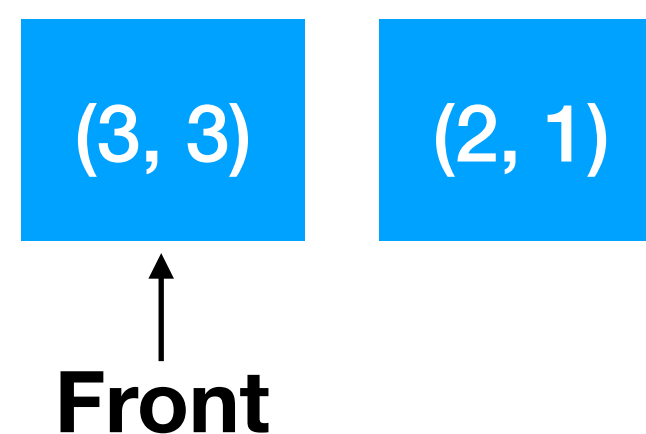


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

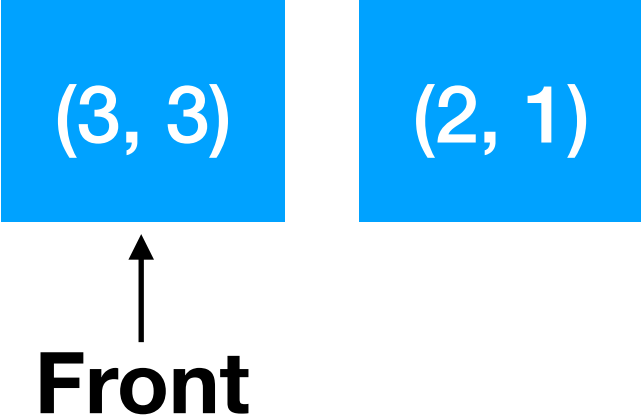
- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (3, 1);`
- `push(2, 1);`
- `pop();`



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

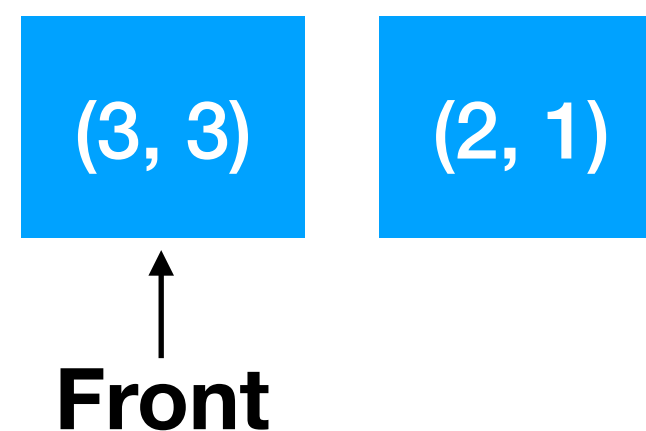


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited

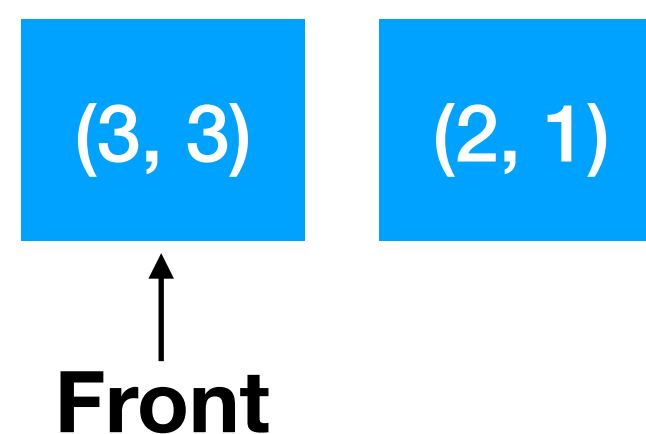


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (3, 3);`



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

BFS Example

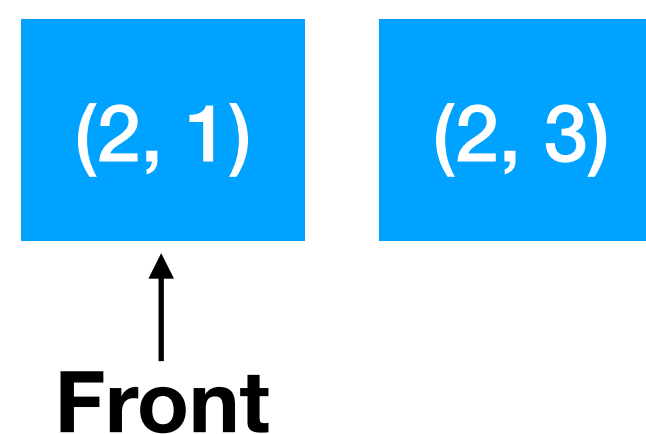
- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (3, 3);`
- `push(2, 3);`



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

BFS Example

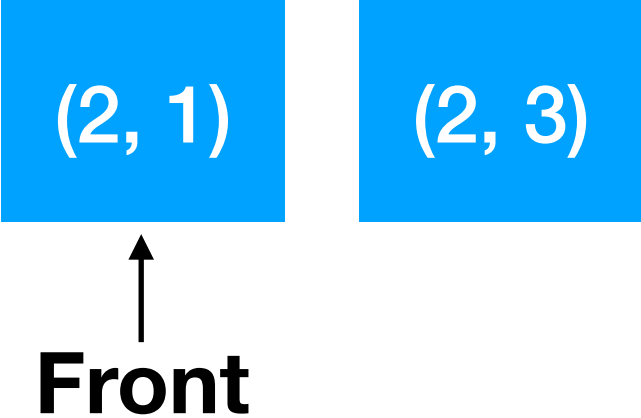
- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (3, 3);`
- `push(2, 3);`
- `pop();`



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

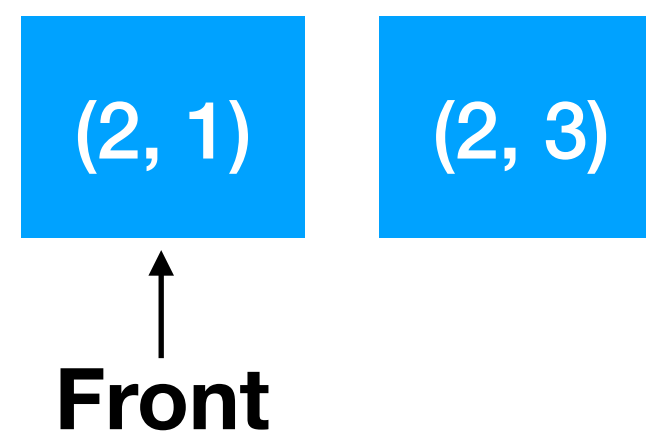


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (2, 1); pop();`

(2, 3)
↑
Front

	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

- For every frontal element, push it's neighbouring reachable coordinates, and mark itself as visited
- `front() = (2, 1); pop();`
- `front() = (2, 3); pop();`

↑
Front

	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#		#		#
4	#		+		#
5		#		#	

Demo

BFS Example

BFS Example

- Assuming map size $N \times M$, what is the time and space complexity?

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 - Is it $\Theta(N \times M)$ as well? Why?

BFS Example

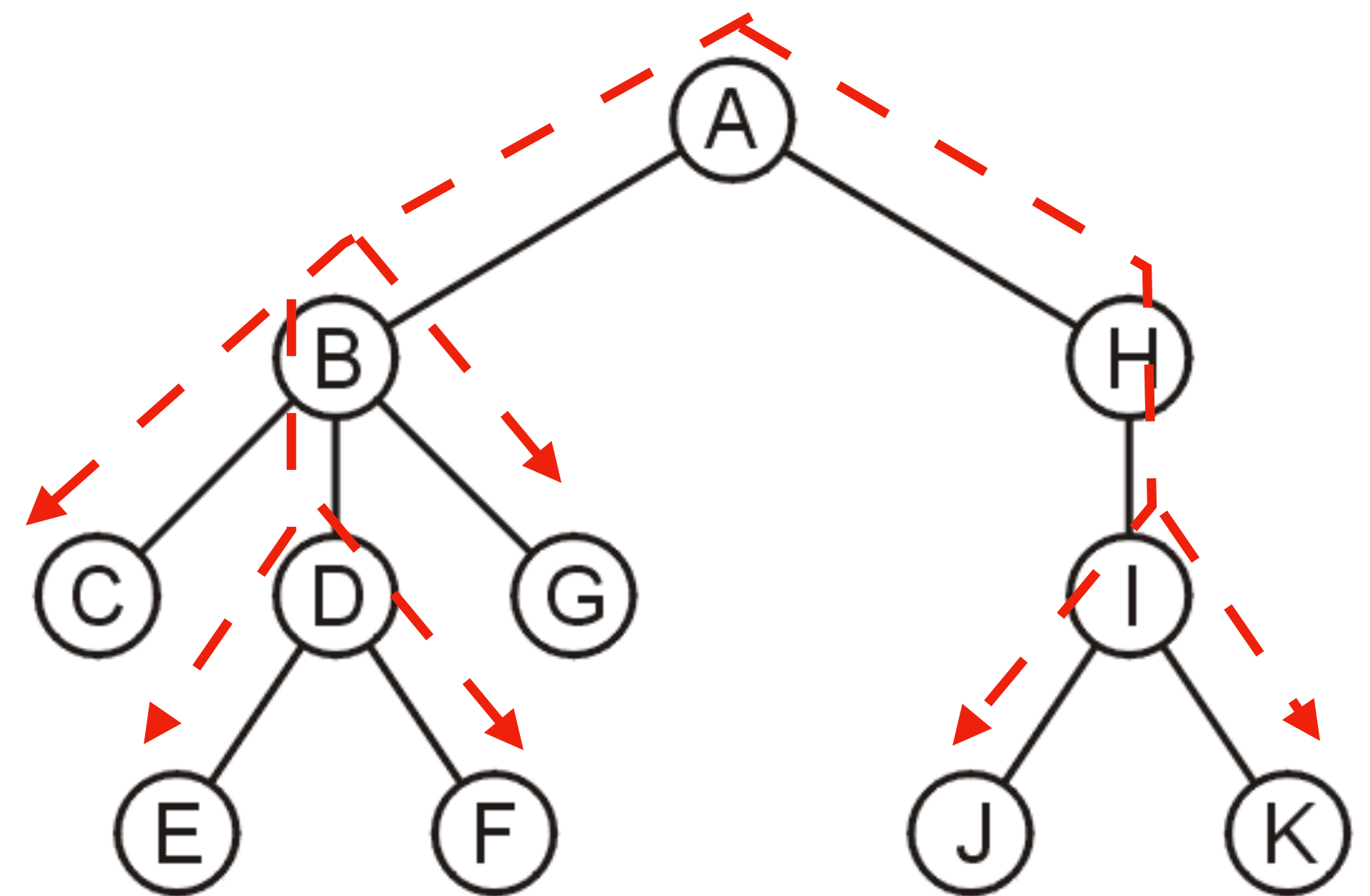
- Assuming map size $N \times M$, what is the time and space complexity?
- Time complexity: $O(N \times M)$
 - Is it $\Theta(N \times M)$ as well? Why?
- Space complexity: $O(N \times M)$

BFS Example

- Assuming map size $N \times M$, what is the time and space complexity?
- Time complexity: $O(N \times M)$
 - Is it $\Theta(N \times M)$ as well? Why?
- Space complexity: $O(N \times M)$
 - Is it $\Theta(N \times M)$ as well? Why?

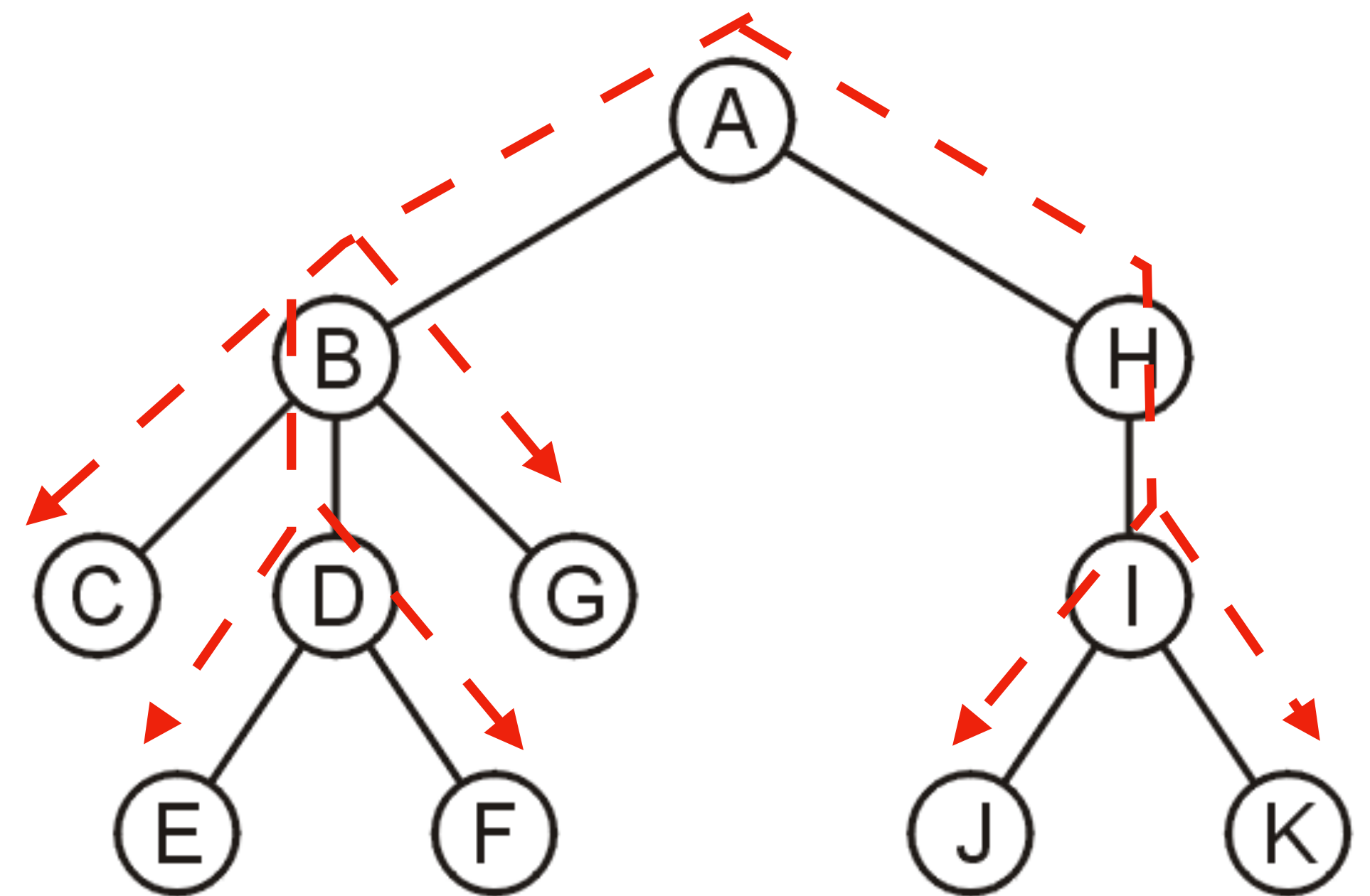
Depth-First Search

- Algorithm for traversing or searching
- Take a single route till the end, then look for other routes
- Implementation using Stack
- Works on trees and graphs



Depth-First Search

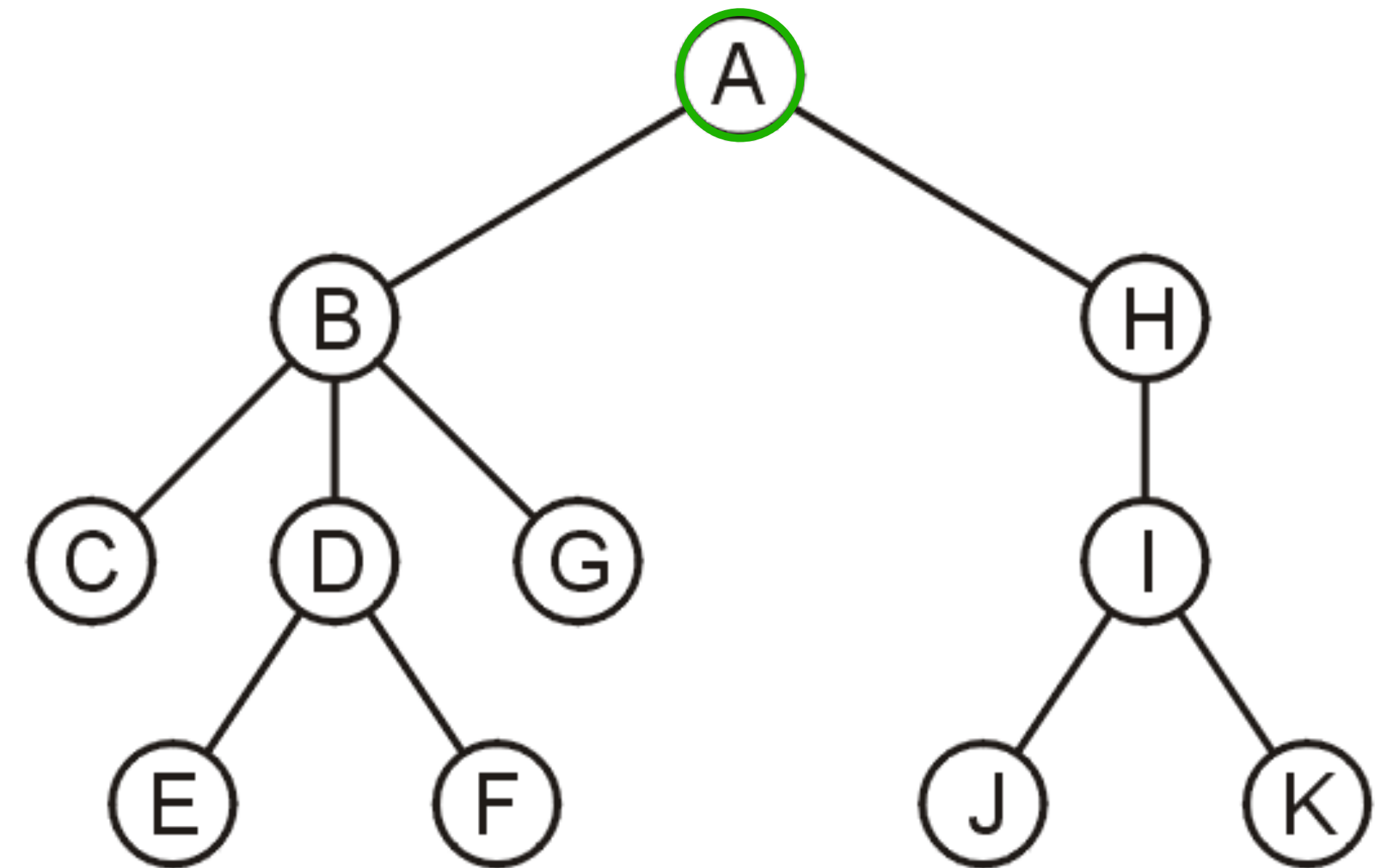
- Depth-First Search (BFS)
- Go all the way till there's no way forward, then go back to find another way
- Classic searching algorithm



DFS Tree

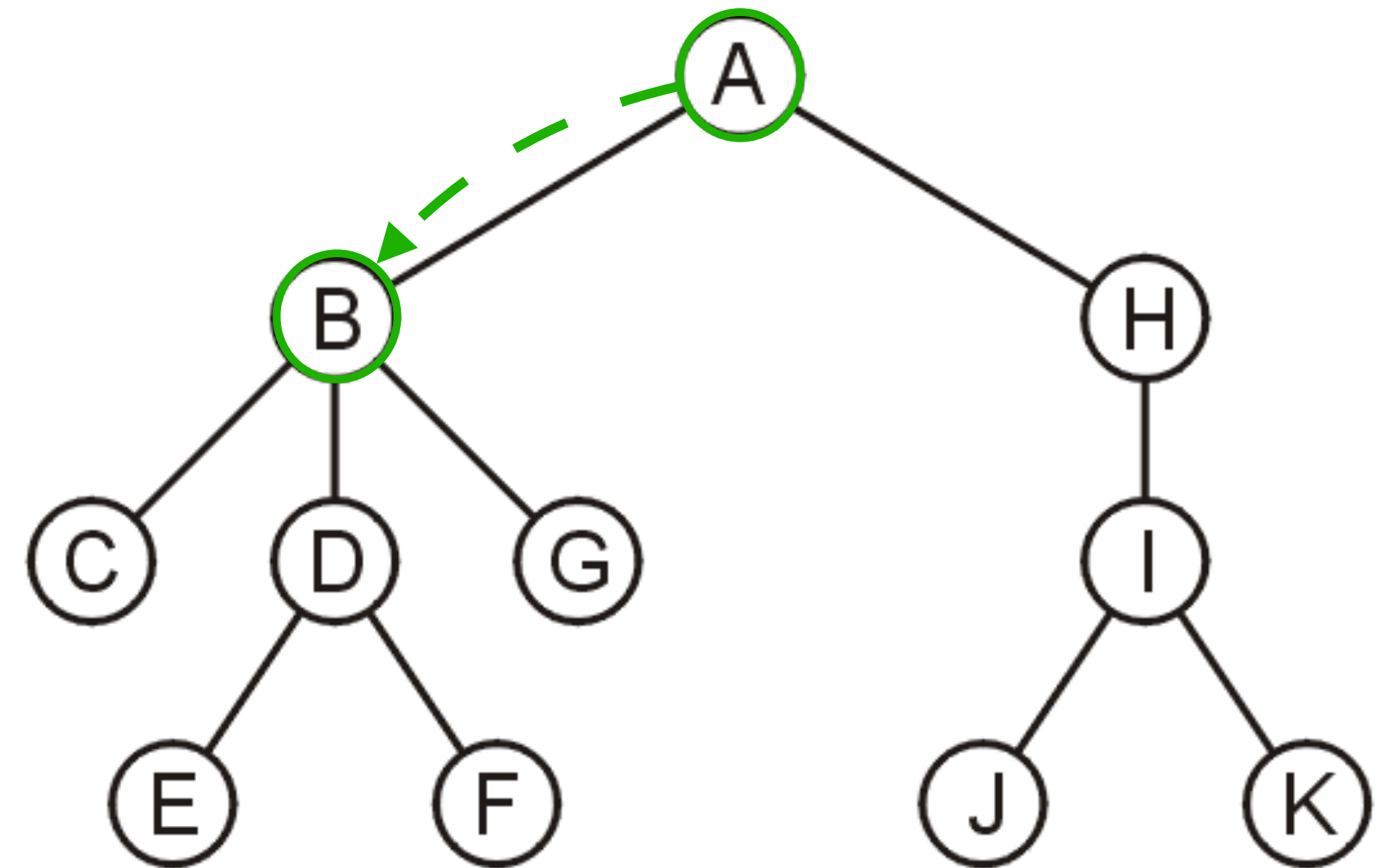
- Push the root directory A

A



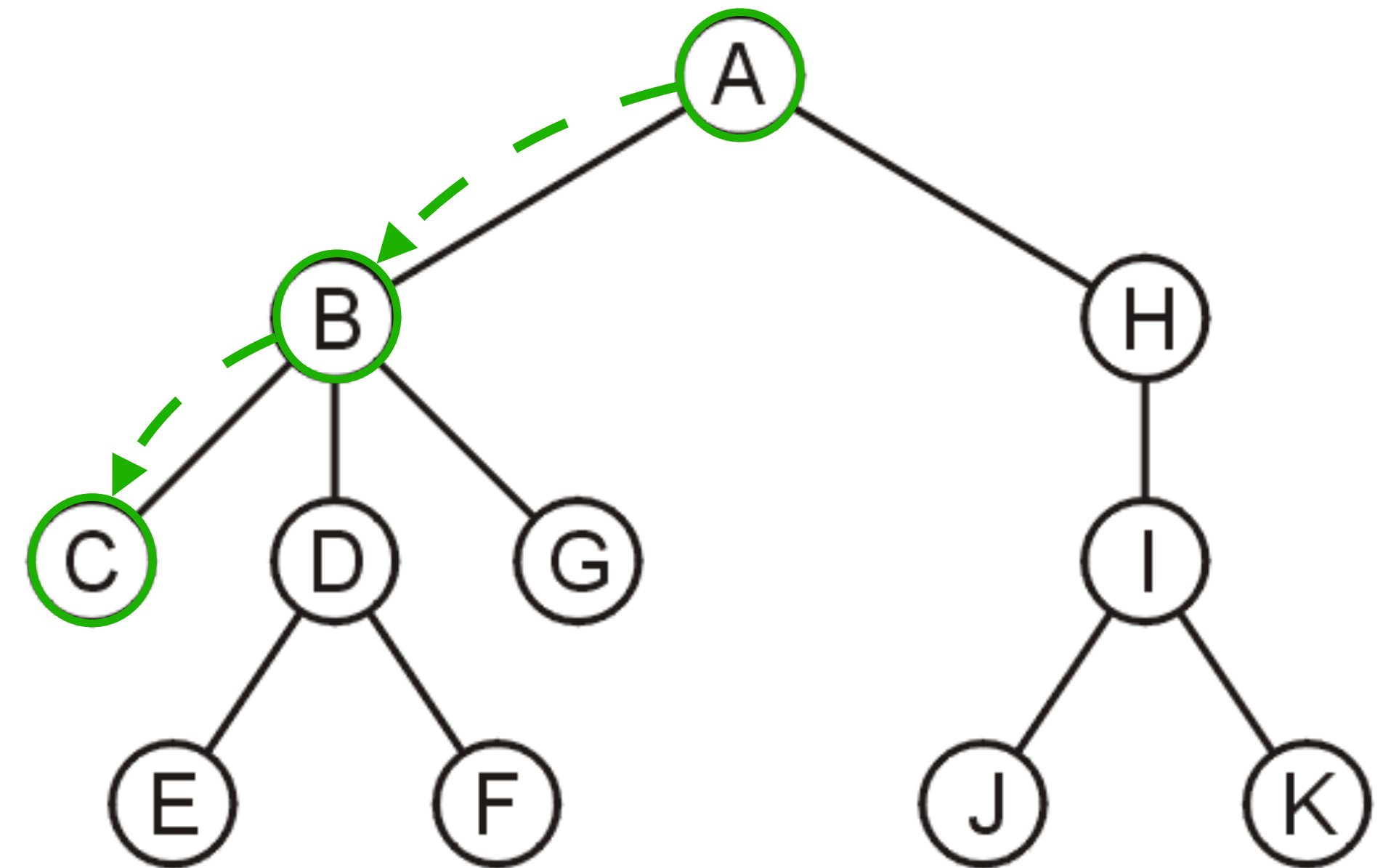
DFS Tree

- Push B



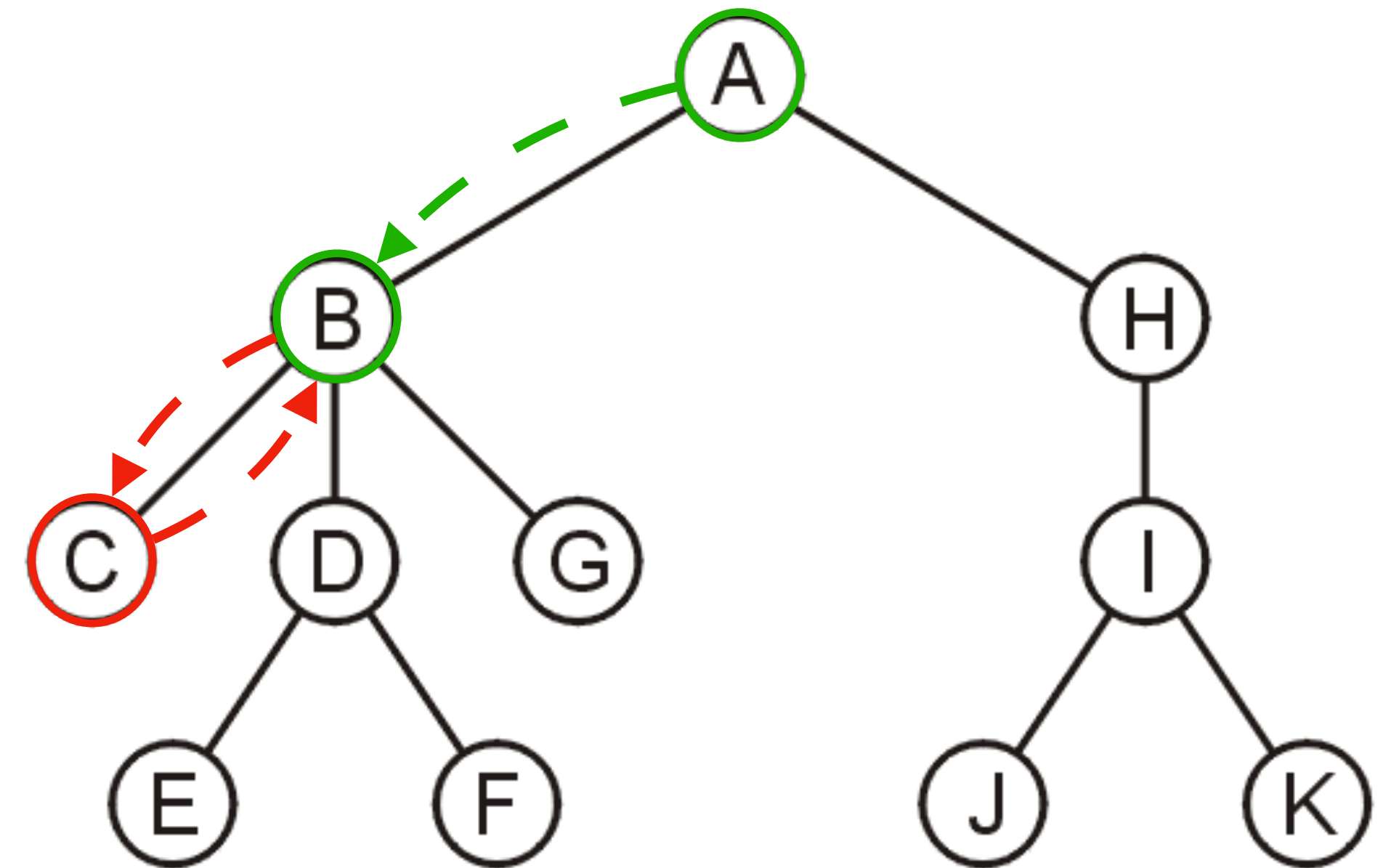
DFS Tree

- Push C



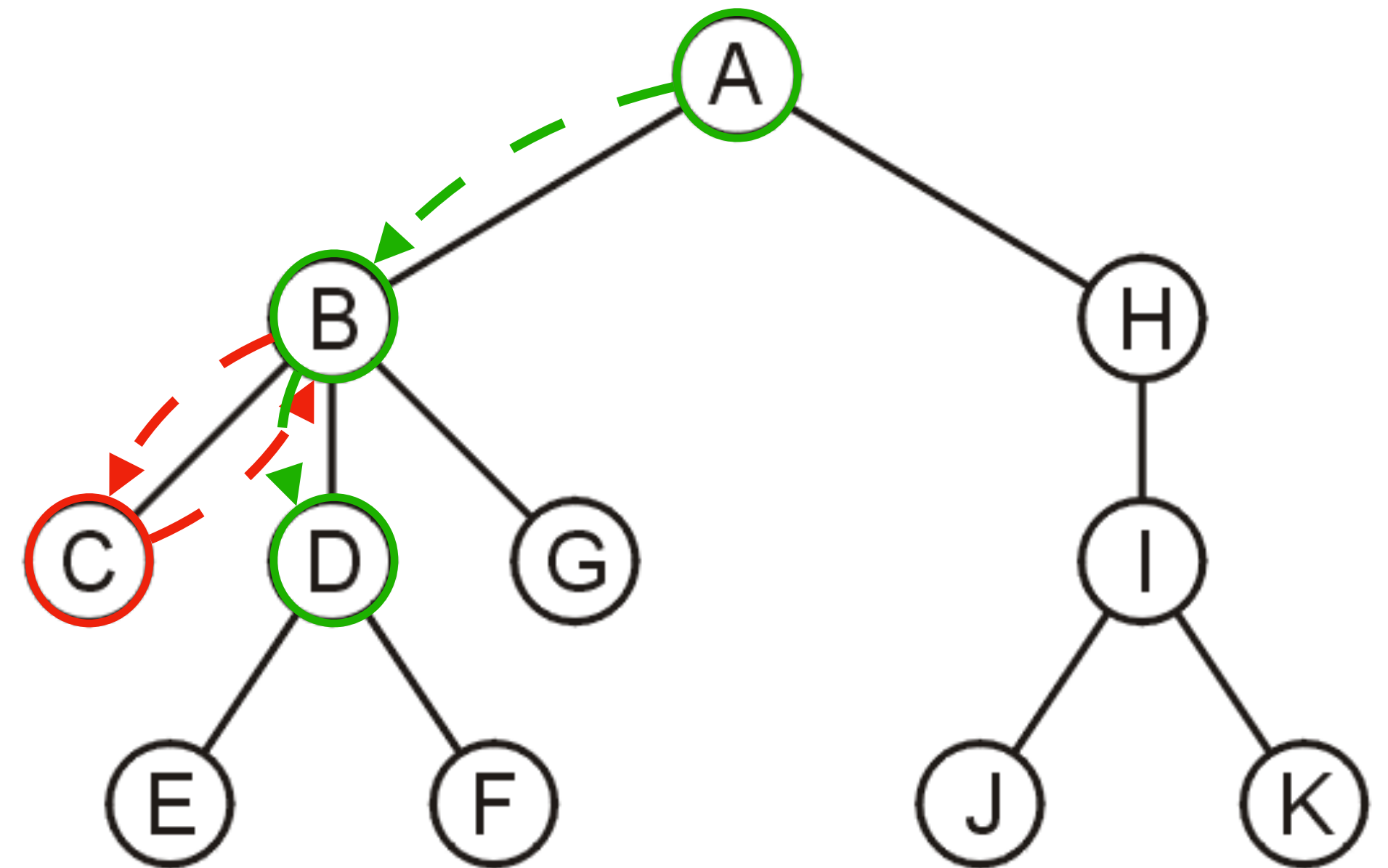
DFS Tree

- Cannot go any further, pop C



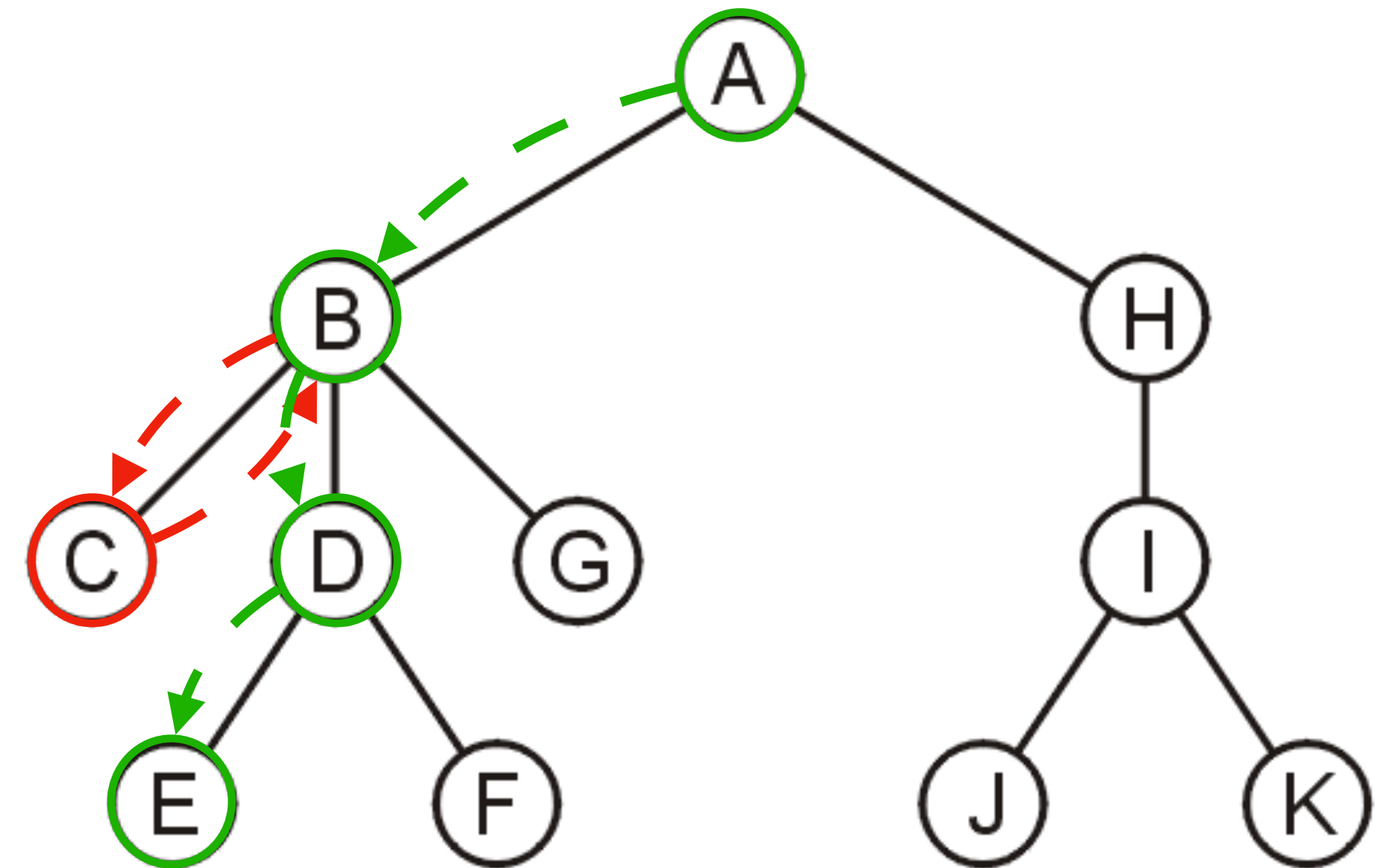
DFS Tree

- Push D, the next not yet visited child of B



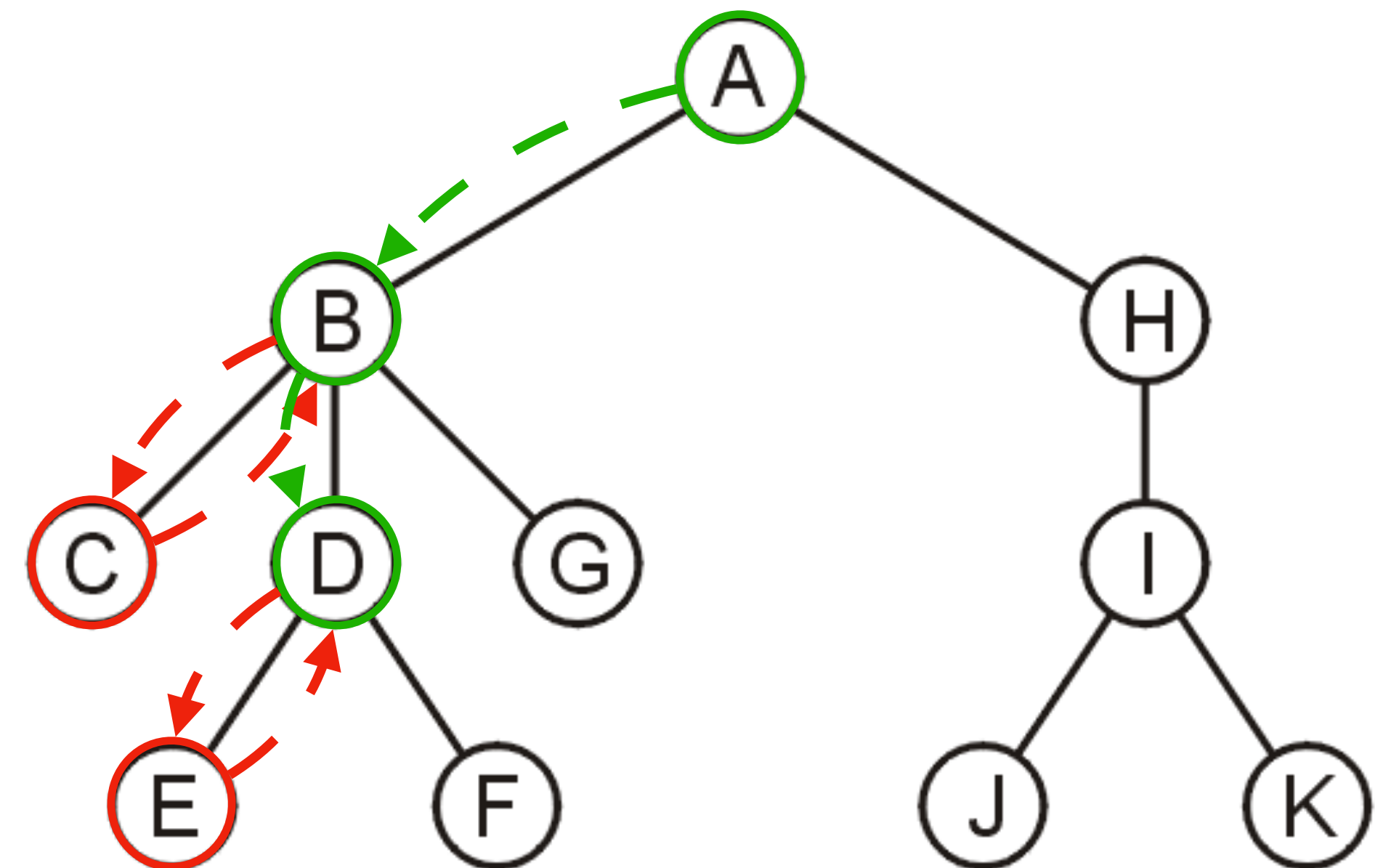
DFS Tree

- Push E



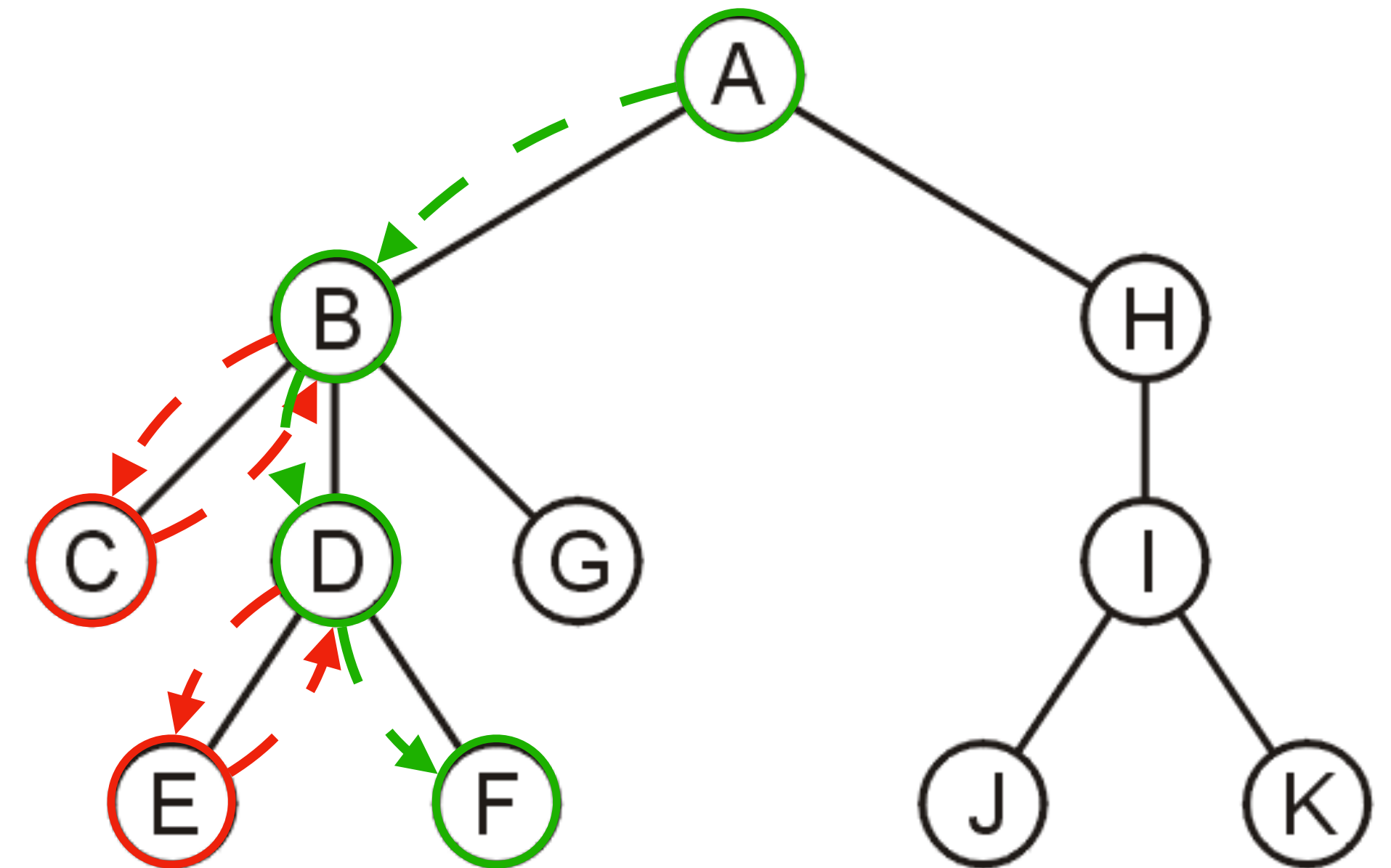
DFS Tree

- Pop E, go back to D



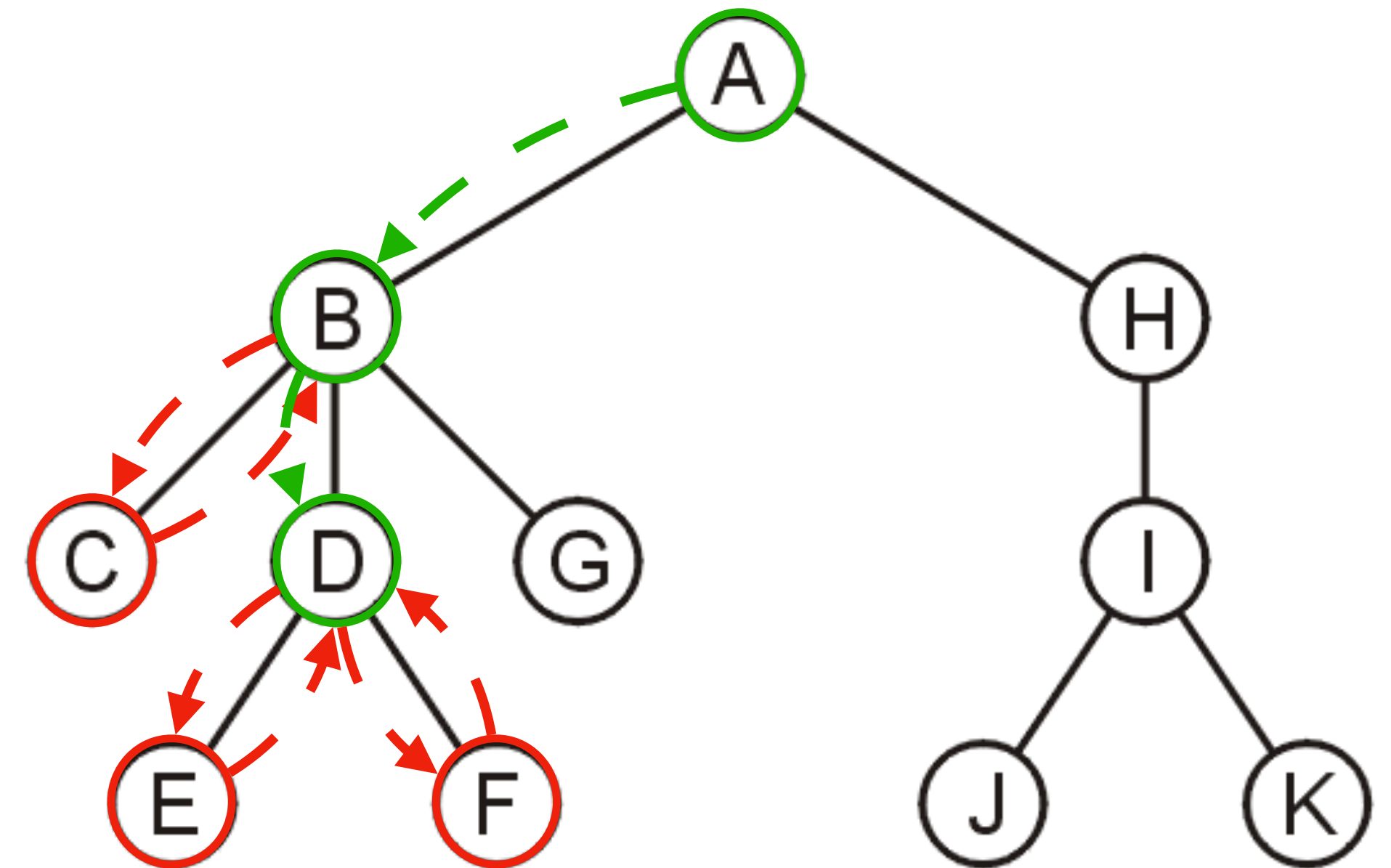
DFS Tree

- Push F



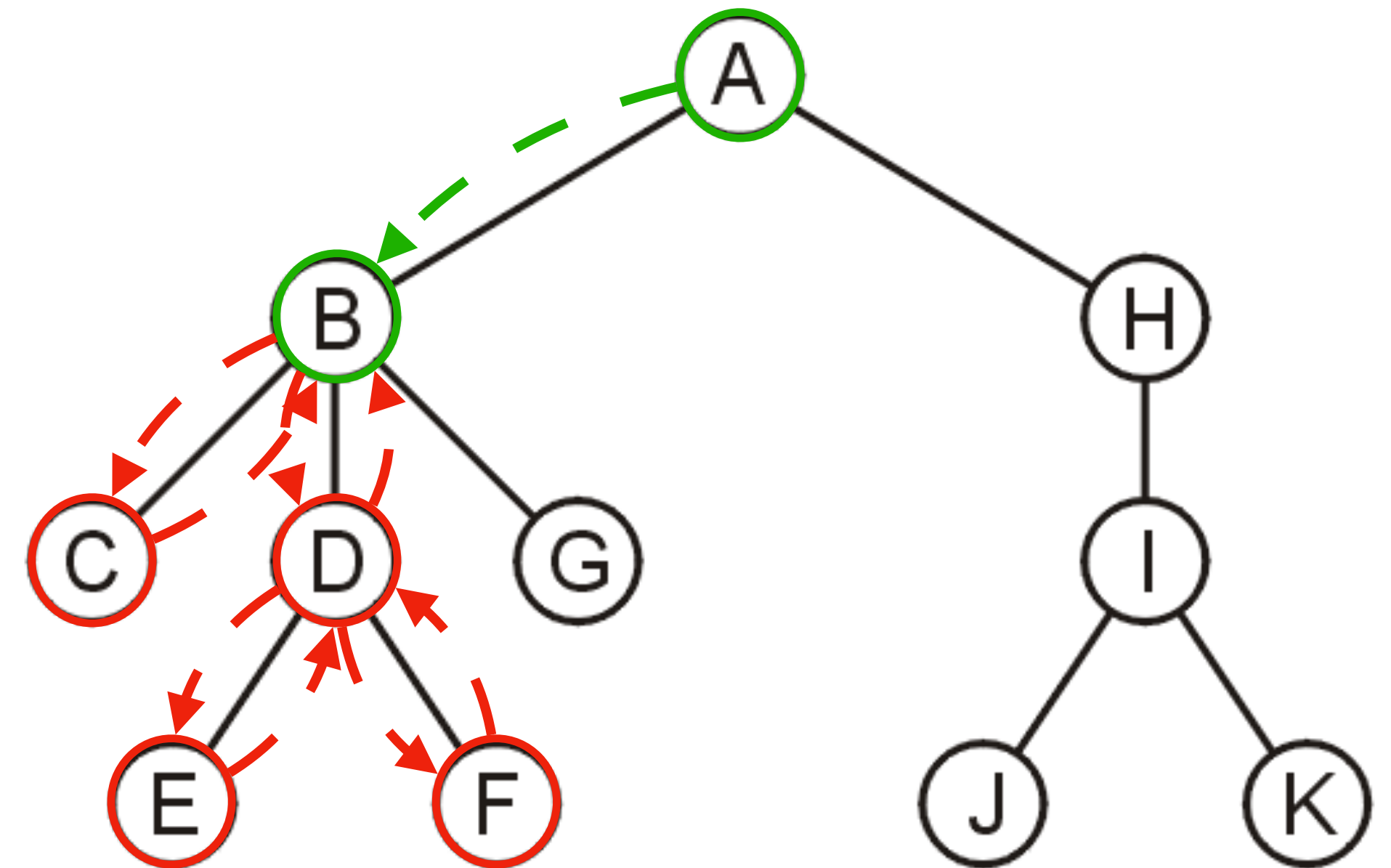
DFS Tree

- Pop F, go back to D



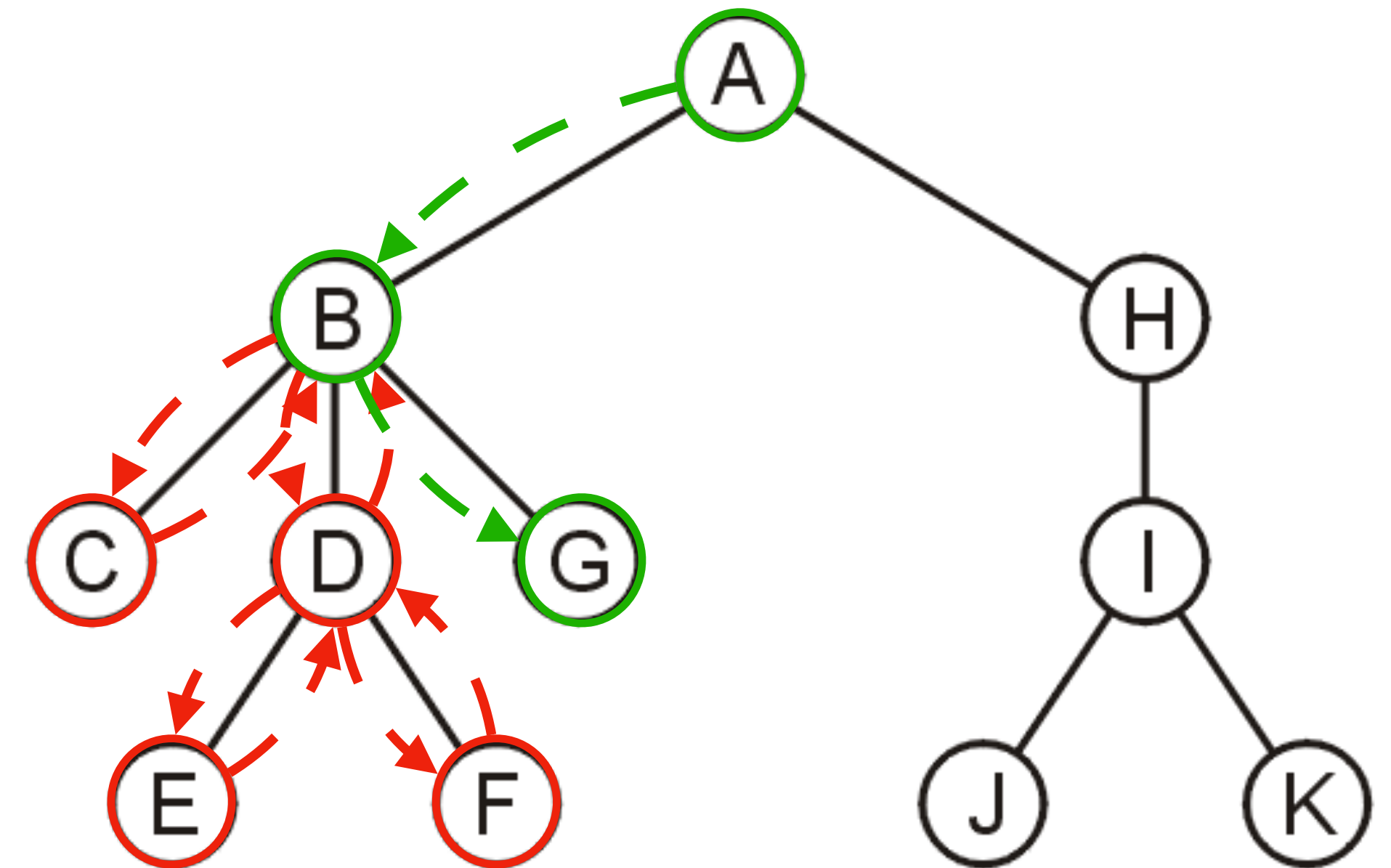
DFS Tree

- Pop D, go back to B



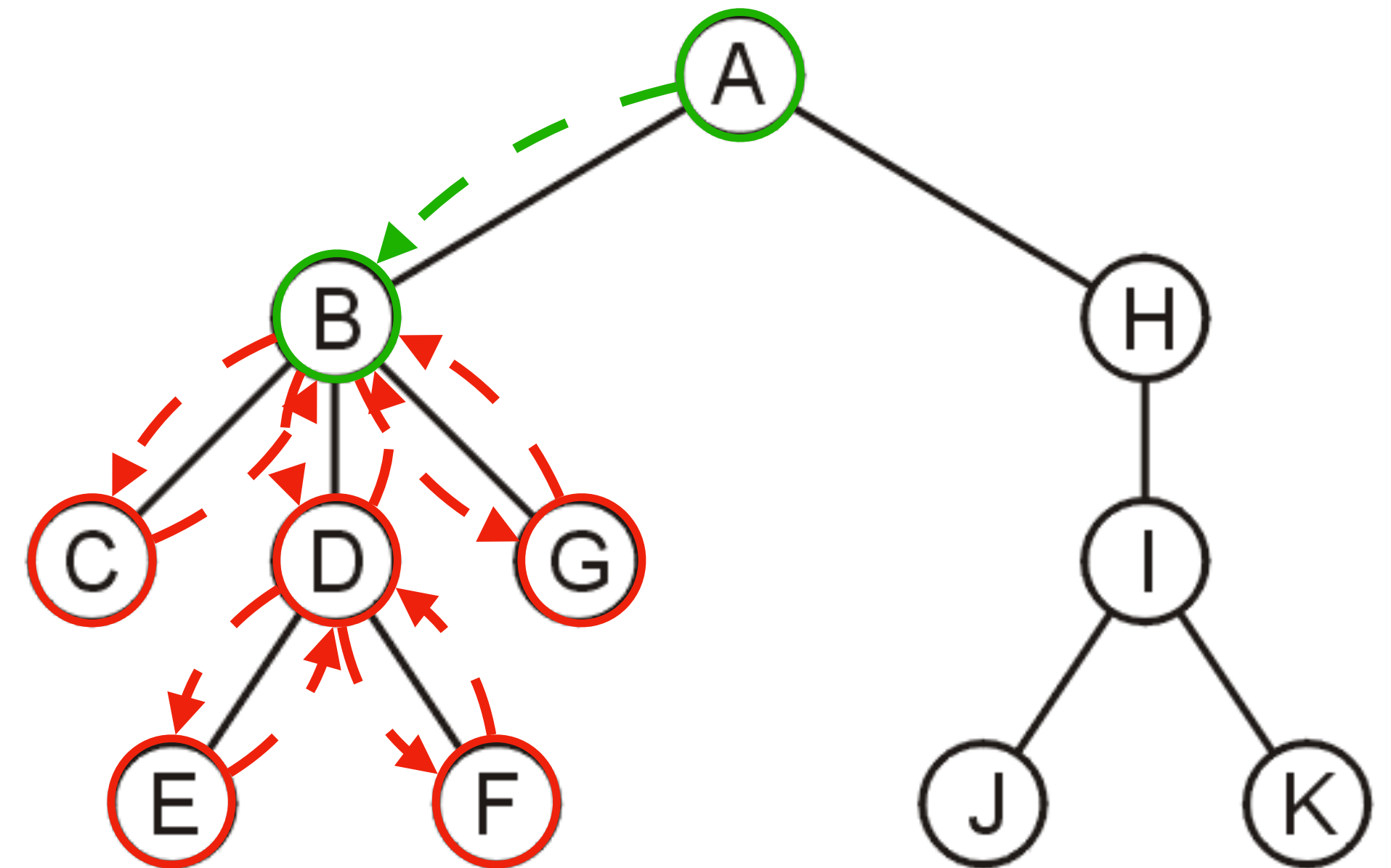
DFS Tree

- Push G



DFS Tree

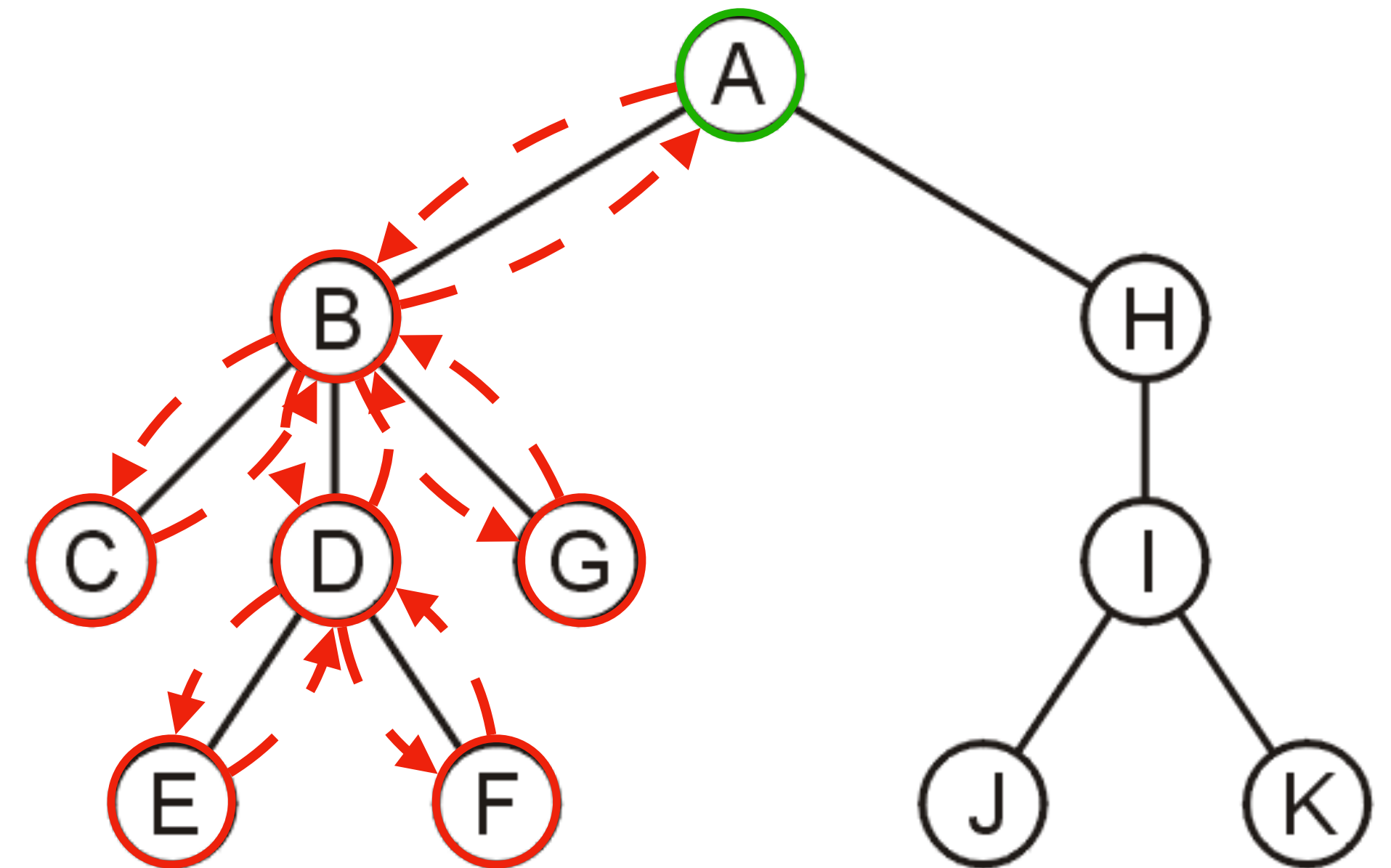
- Pop G, go back to B



DFS Tree

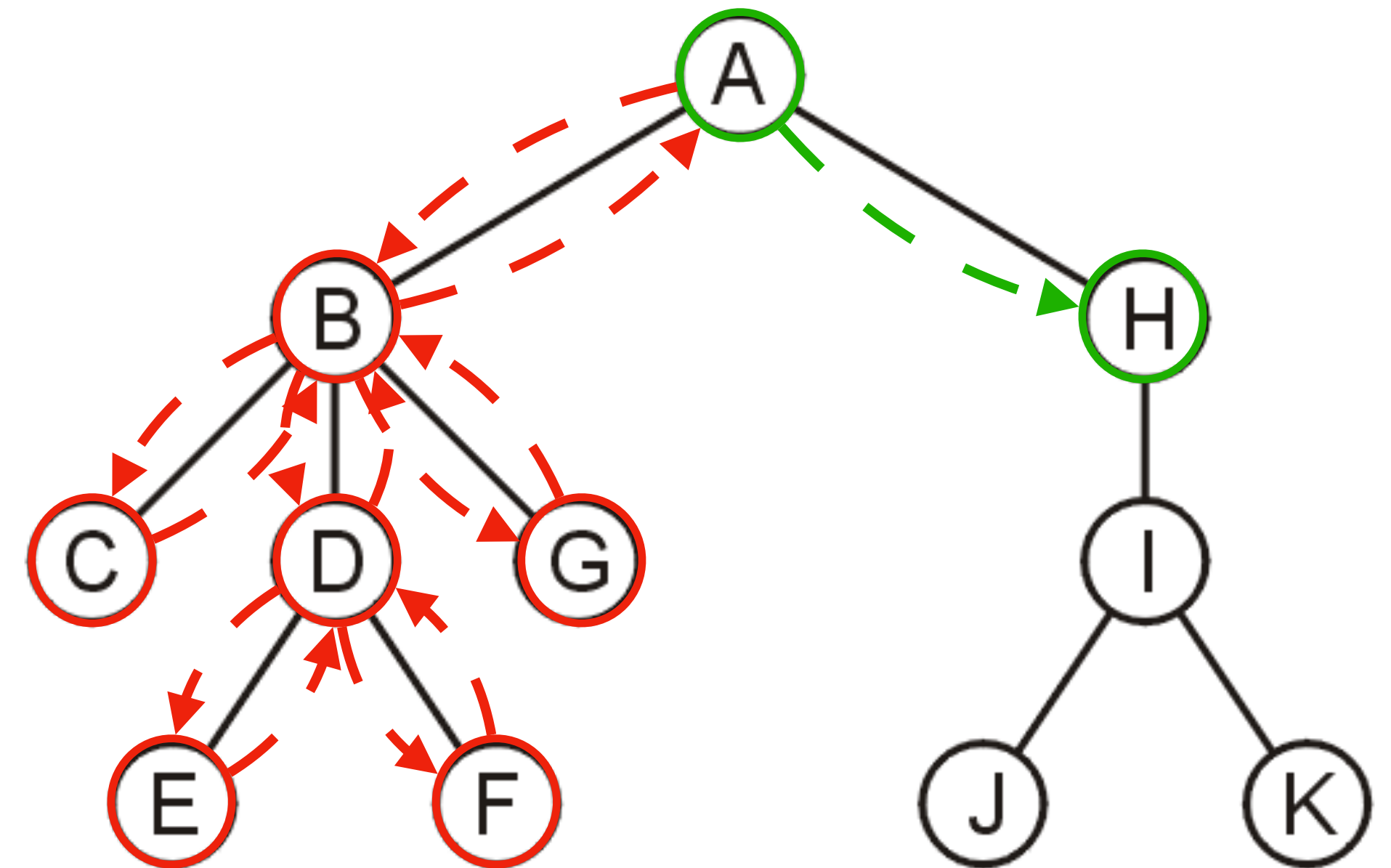
- Pop B, go back to A

A



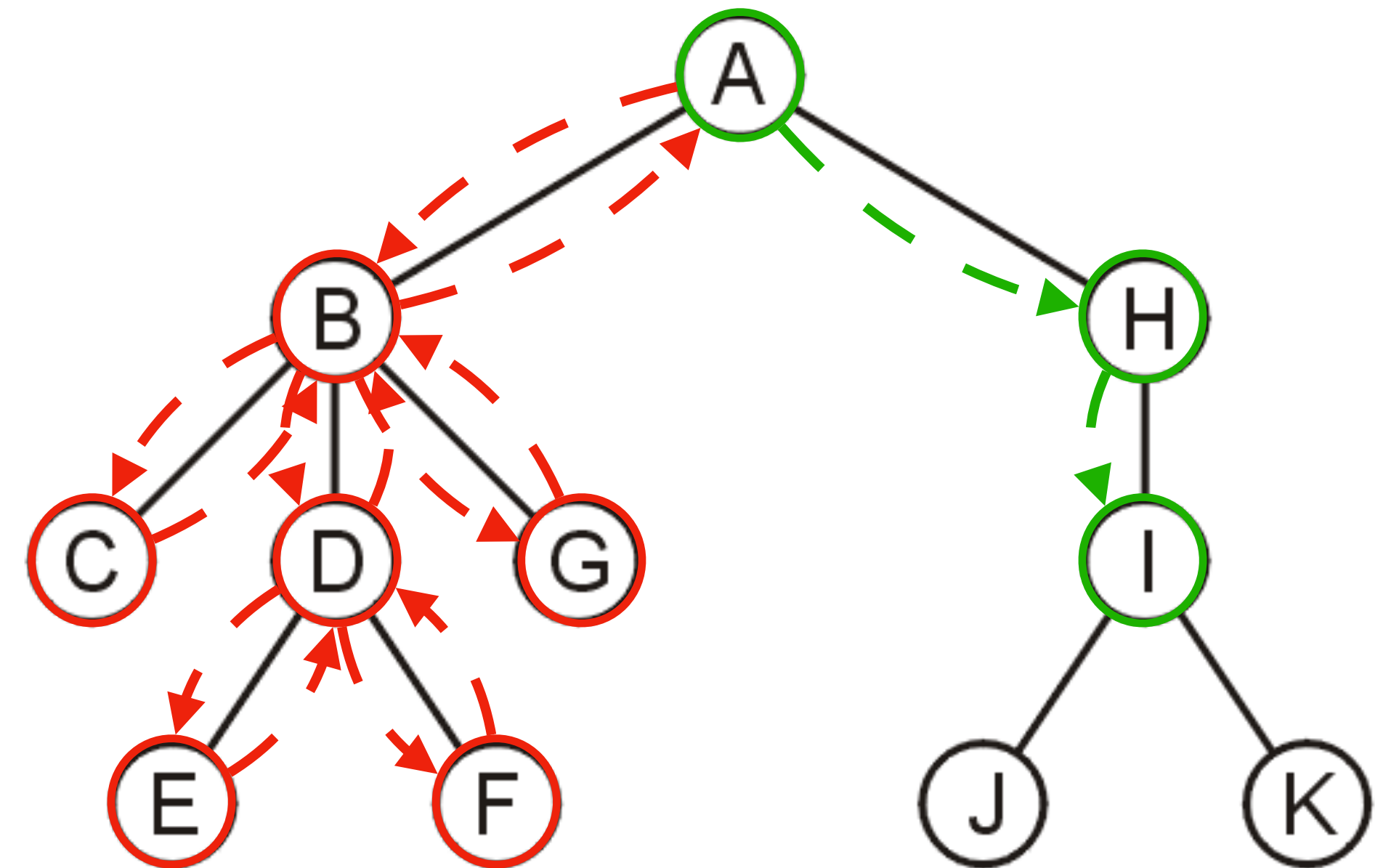
DFS Tree

- Push H



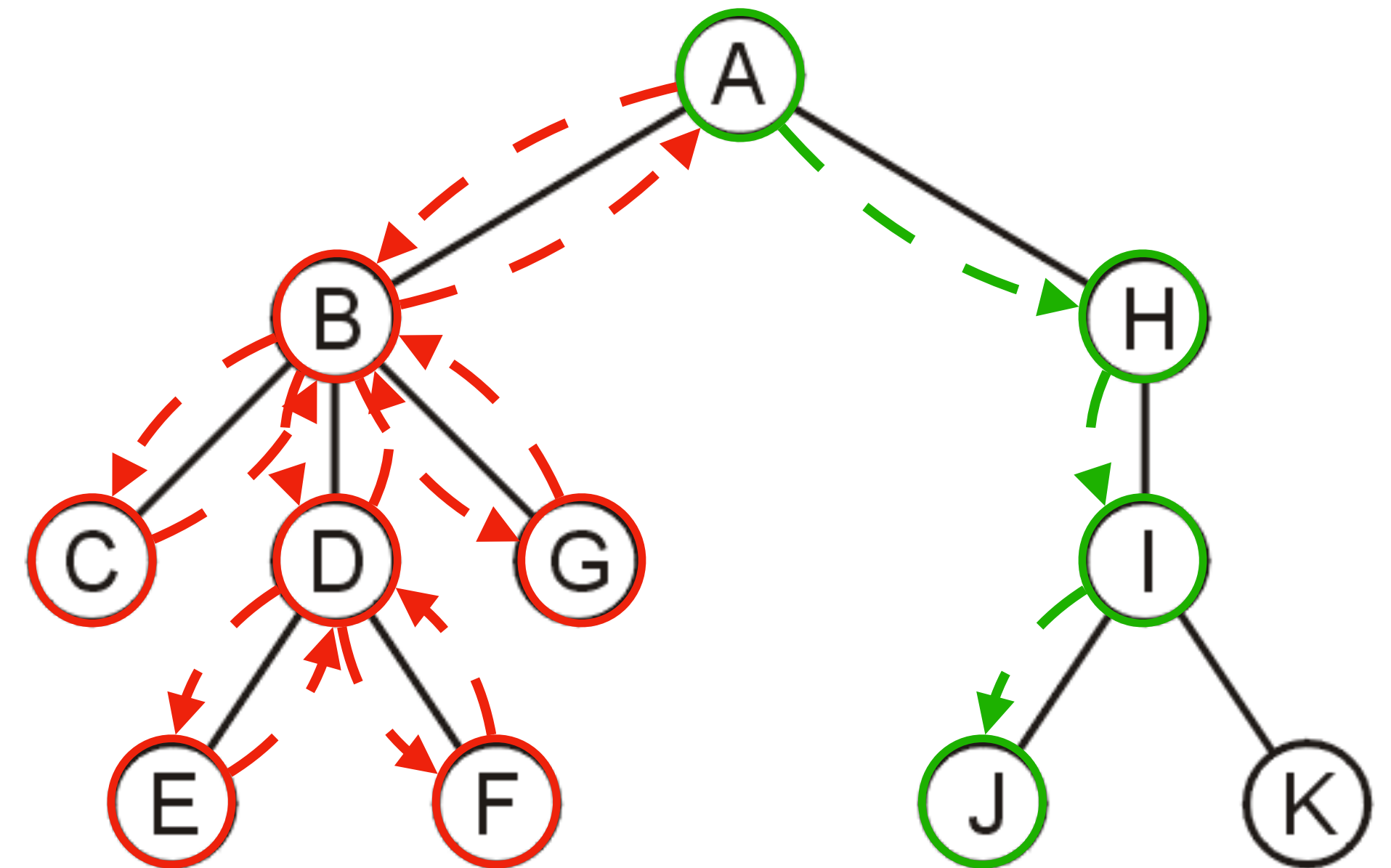
DFS Tree

- Push I



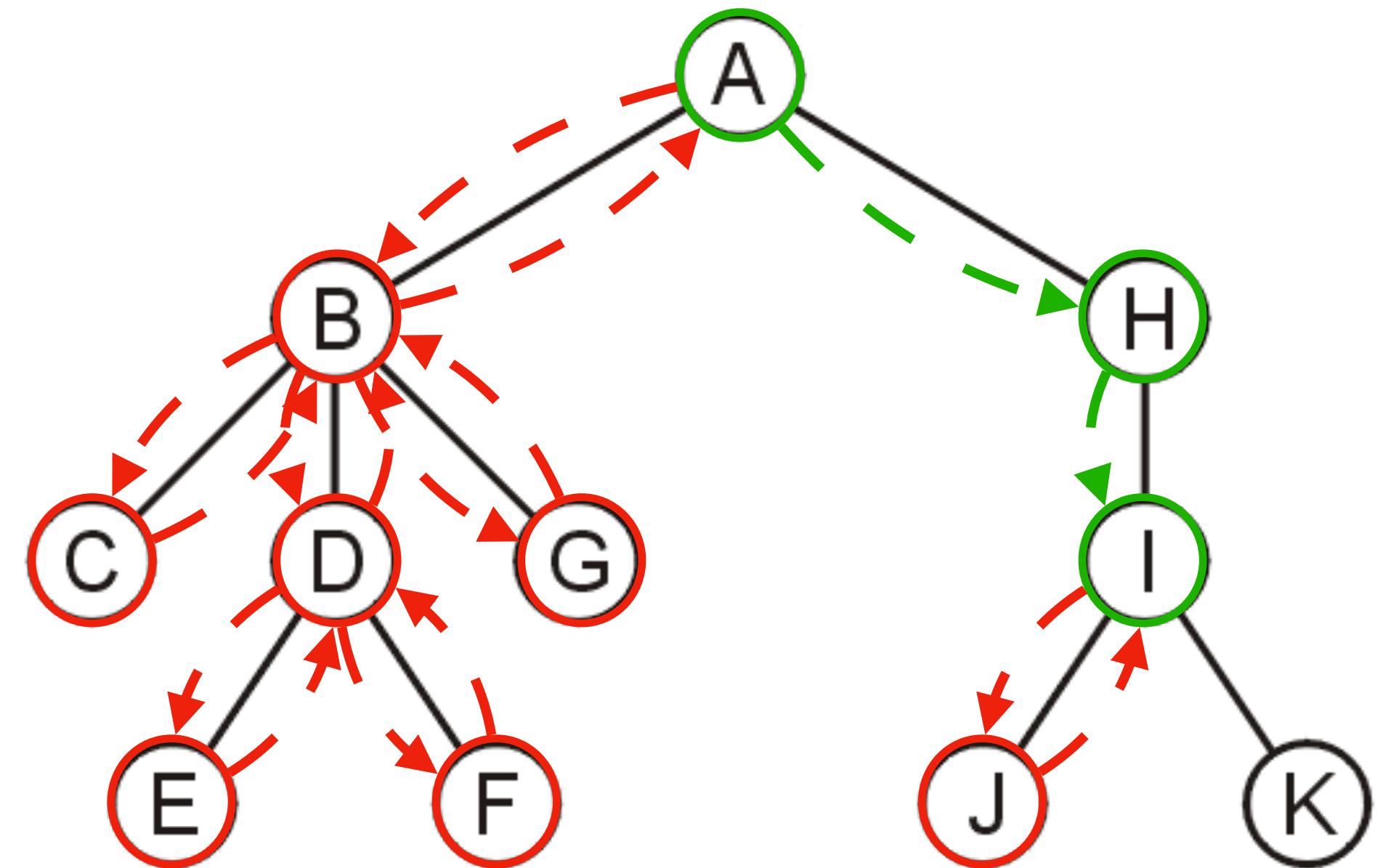
DFS Tree

- Push J



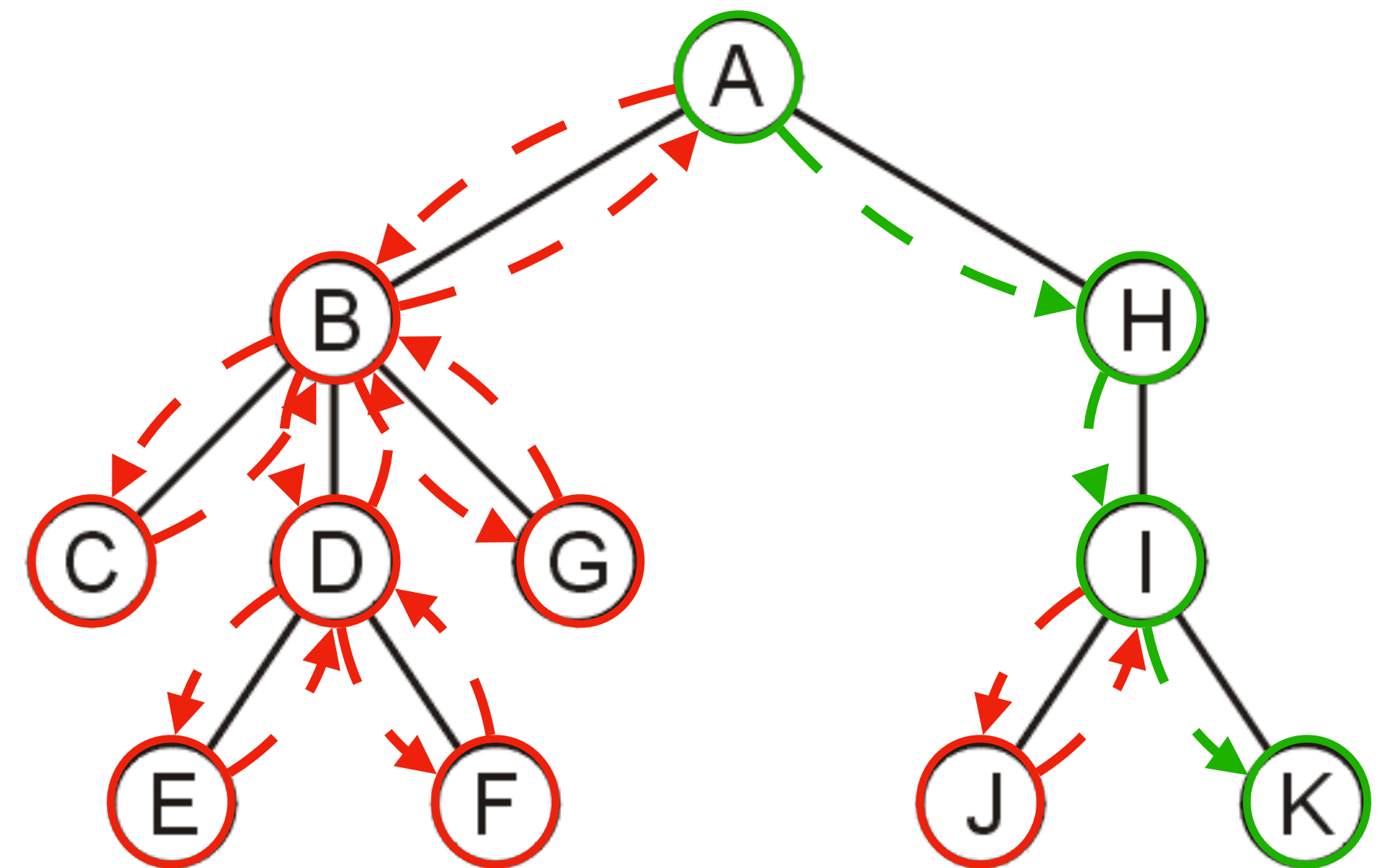
DFS Tree

- Pop J, go back to I



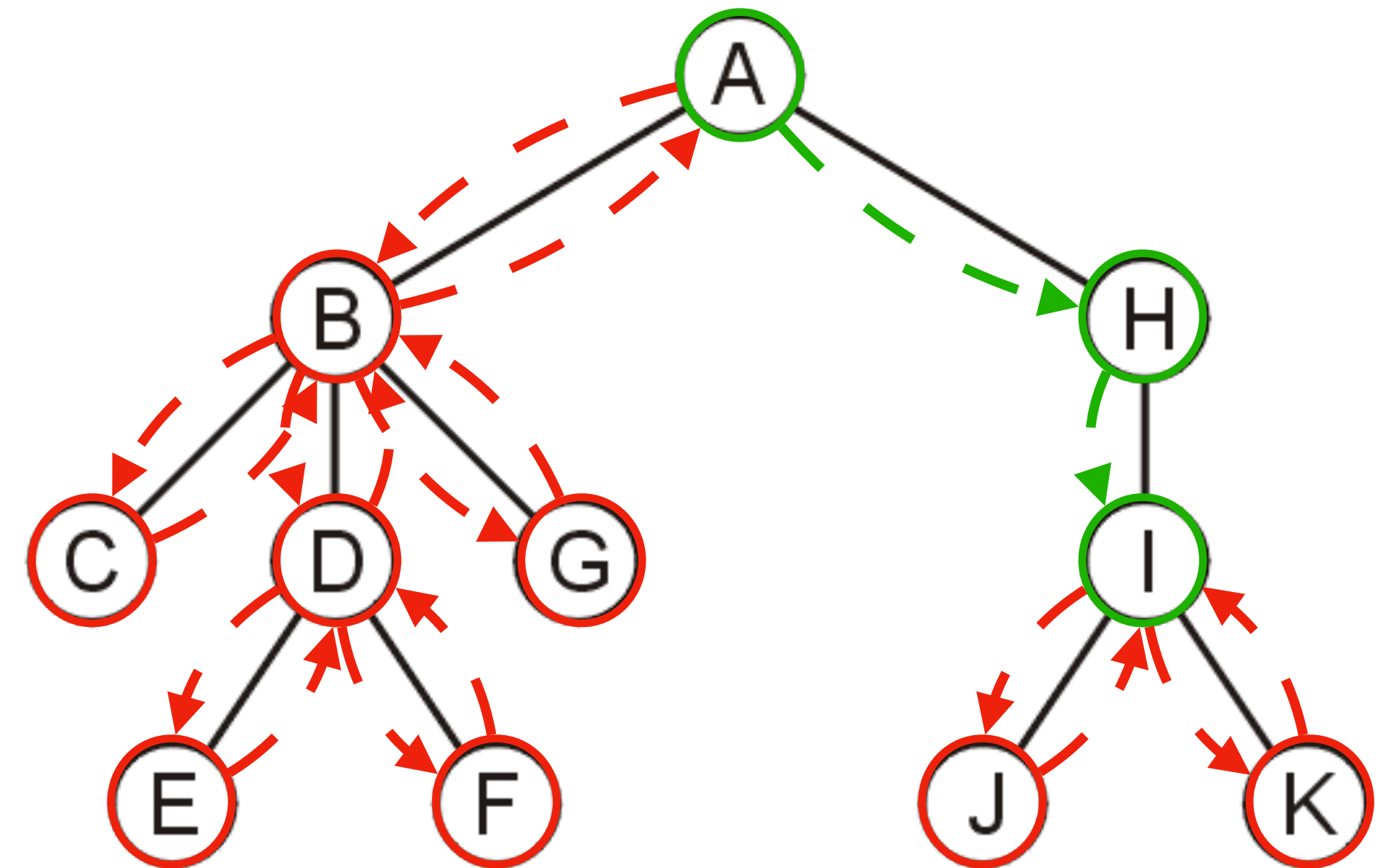
DFS Tree

- Push K



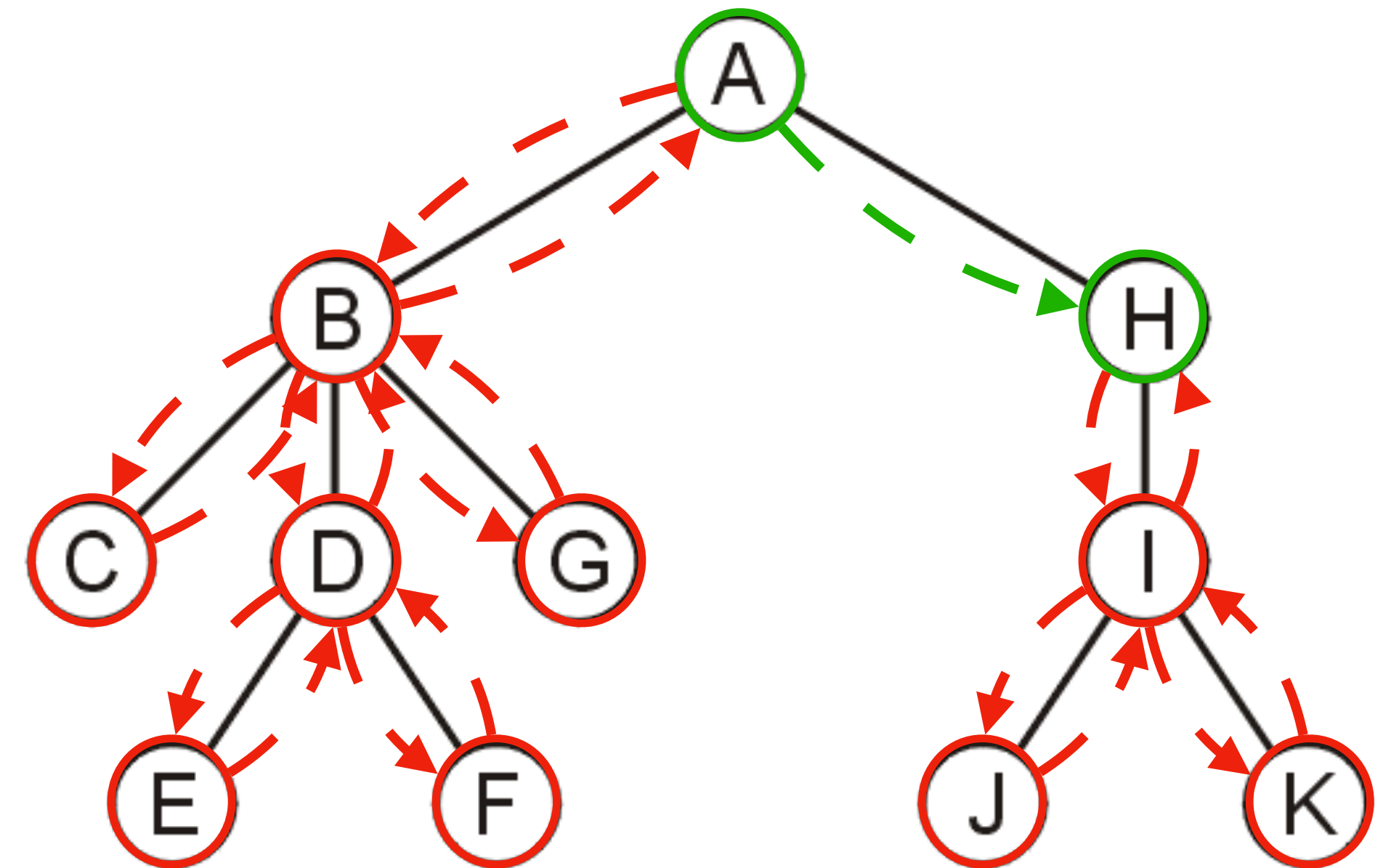
DFS Tree

- Pop K, go back to I



DFS Tree

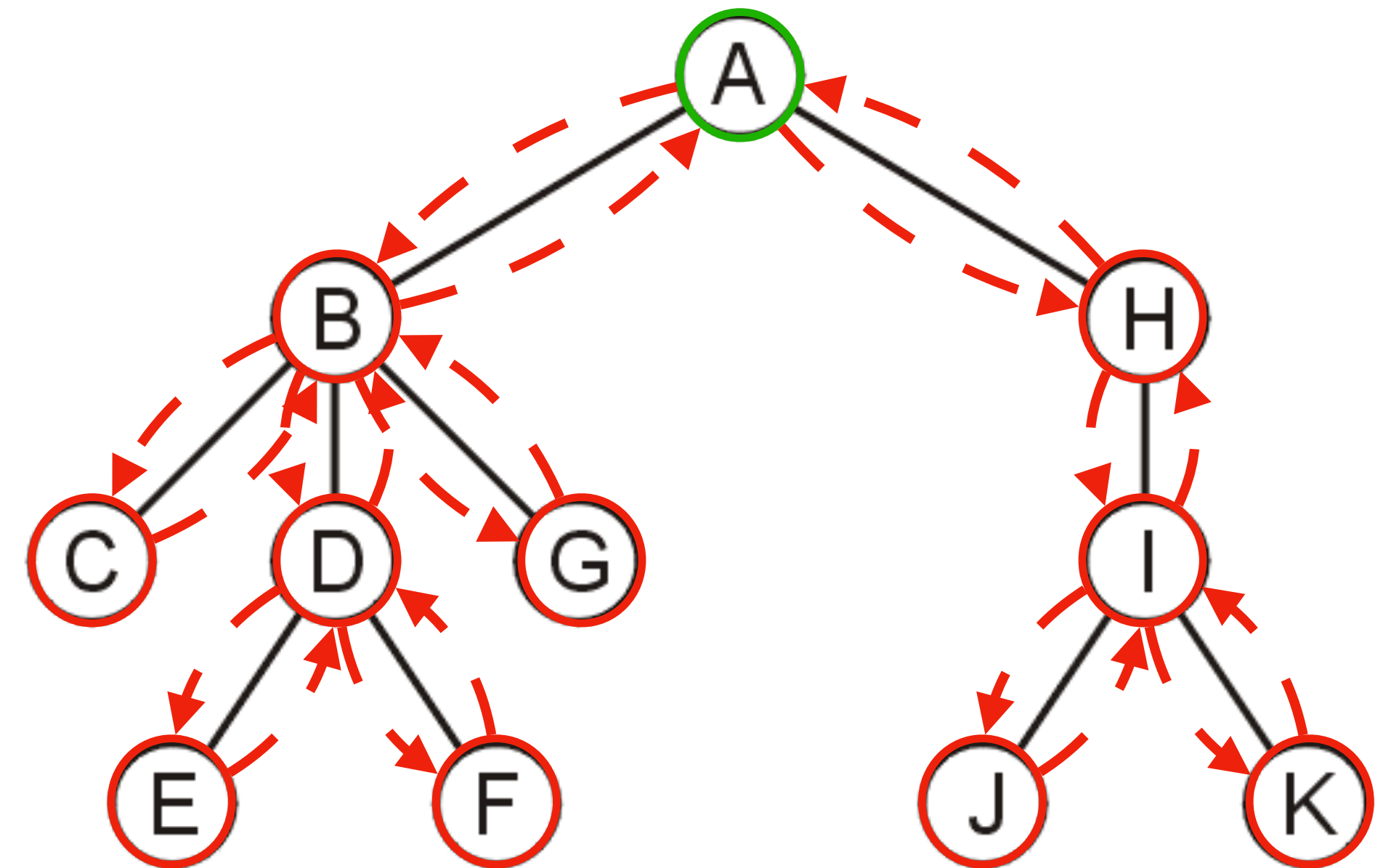
- Pop I, go back to H



DFS Tree

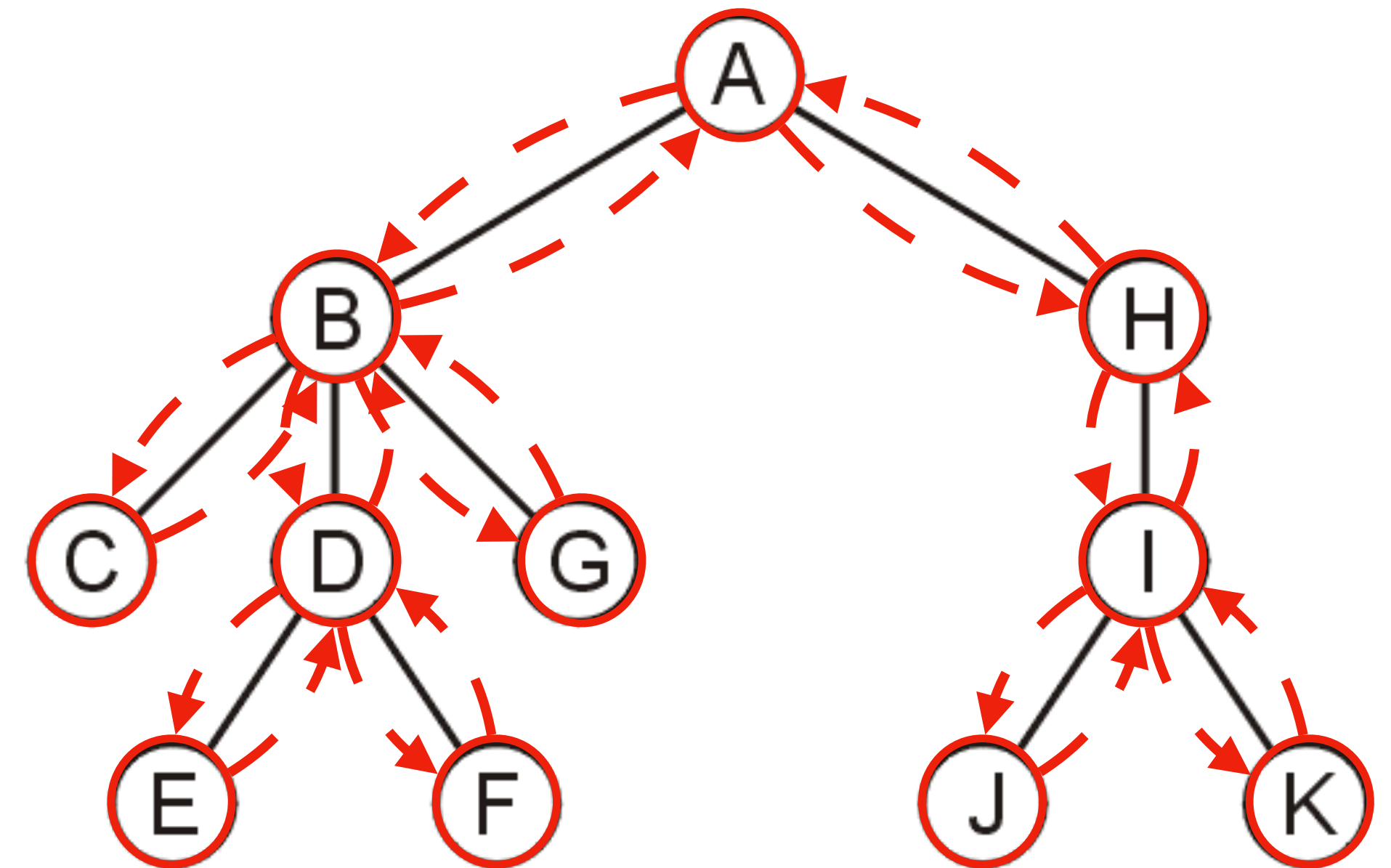
- Pop H, go back to A

A



DFS Tree

- Pop A, traversal completed
- Traversal Sequence
A-B-C-D-E-F-G-H-I-J-K



DFS Tree Analysis

DFS Tree Analysis

- Assuming N nodes, M connections, what is the time and space complexity?

DFS Tree Analysis

- Assuming N nodes, M connections, what is the time and space complexity?
- Time complexity: $O(N + M)$

DFS Tree Analysis

- Assuming N nodes, M connections, what is the time and space complexity?
- Time complexity: $O(N + M)$
 - Is it $\Theta(N + M)$ as well? Why?

DFS Tree Analysis

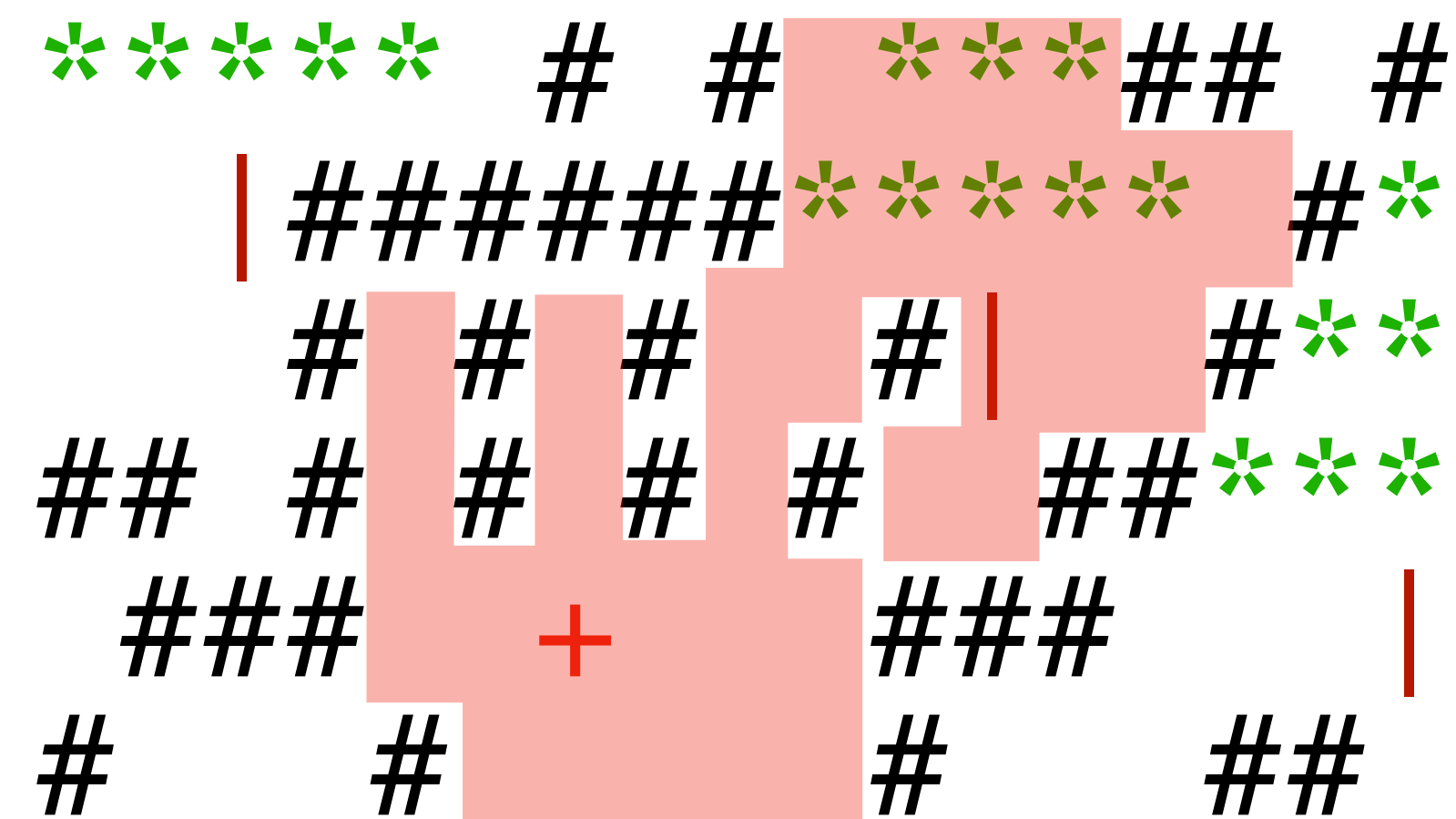
- Assuming N nodes, M connections, what is the time and space complexity?
- Time complexity: $O(N + M)$
 - Is it $\Theta(N + M)$ as well? Why?
- Space complexity: $\Theta(N + M)$

DFS Tree Analysis

- Assuming N nodes, M connections, what is the time and space complexity?
- Time complexity: $O(N + M)$
 - Is it $\Theta(N + M)$ as well? Why?
- Space complexity: $\Theta(N + M)$
 - Is it $\Theta(N + M)$ as well? Why?

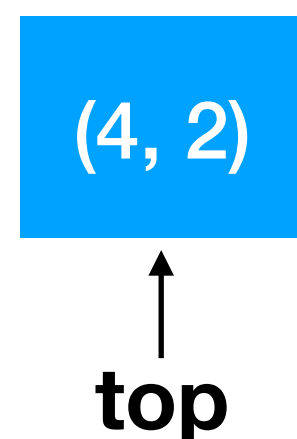
DFS Example

- Village of the Sorcerer
- Certain areas of the map are not reachable (fences, houses)
- Certain areas might be entirely blocked off by fences
- **+**: player
- **Light Red**: player accessible regions



DFS Example

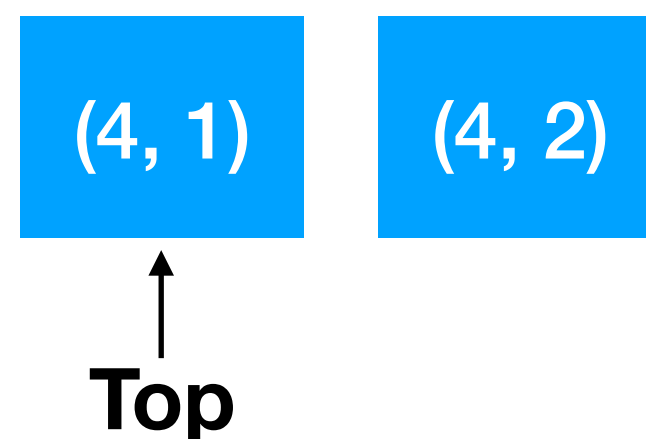
- We start by initialising a Stack to store coordinates
- `push(playerCoordinate)`, which is 4,2
 - `push ((4, 2)) ;`
- In C++, you can simulate this using 2 `int`-based stacks



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#		+	#	#
5		#		#	

DFS Example

- For every frontal element, traverse randomly until cannot further
- $\text{top}() = (4, 2);$
- $\text{push}(4, 1)$



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#		+	#	#
5		#		#	

Demo

DFS Example

- For every frontal element, traverse randomly until cannot further
 - `top() = (4, 1);`
 - `push(3, 1);`

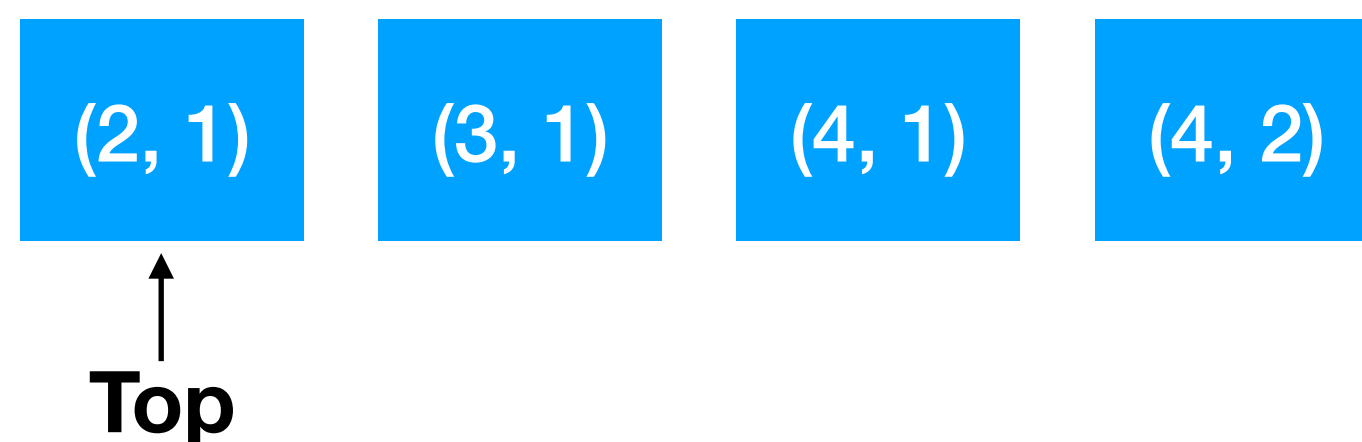


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#			#	#
5		#		#	

Demo

DFS Example

- For every frontal element, traverse randomly until cannot further
- `top() = (3, 1);`
- `push(2, 1)`



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#		+	#	#
5		#		#	

Demo

DFS Example

- For every frontal element, traverse randomly until cannot further
- $\text{top}() = (2, 1);$
- $\text{pop}()$

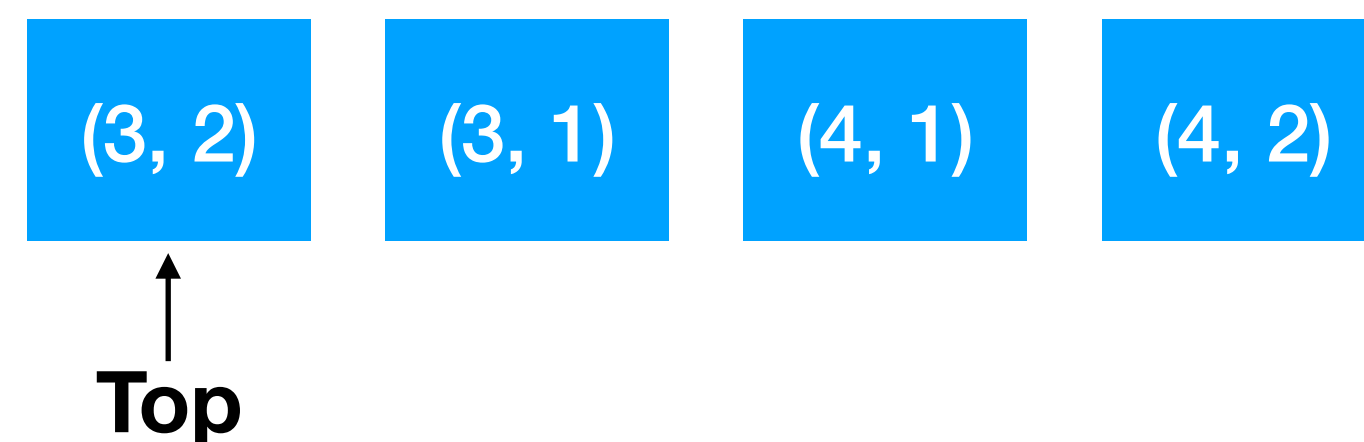


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#		+	#	#
5		#		#	

Demo

DFS Example

- For every frontal element, traverse randomly until cannot further
- $\text{top}() = (3, 1);$
- $\text{push}(3, 2)$



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#			#	#
5		#		#	

Demo

DFS Example

- For every frontal element, traverse randomly until cannot further
- $\text{top}() = (3, 2);$
- $\text{push}(3, 3)$



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#				#
5		#		#	

Demo

DFS Example

- For every frontal element, traverse randomly until cannot further
- $\text{top}() = (3, 3);$
- $\text{push}(2, 3)$



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#		+	#	#
5		#		#	

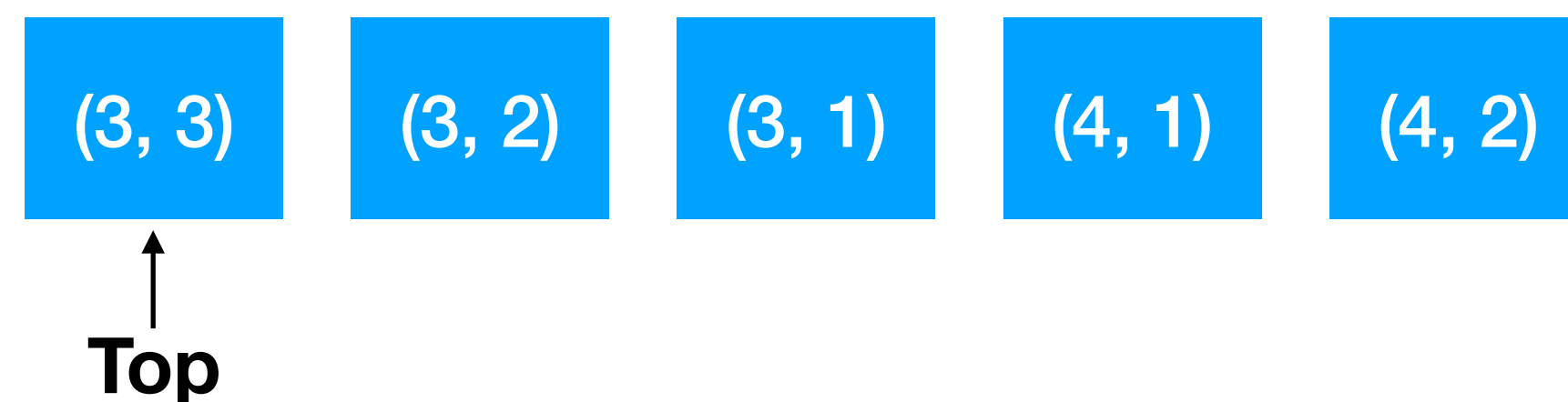
Demo

DFS Example

- For every frontal element, traverse randomly until cannot further

- $\text{top}() = (2, 3);$

- $\text{pop}()$

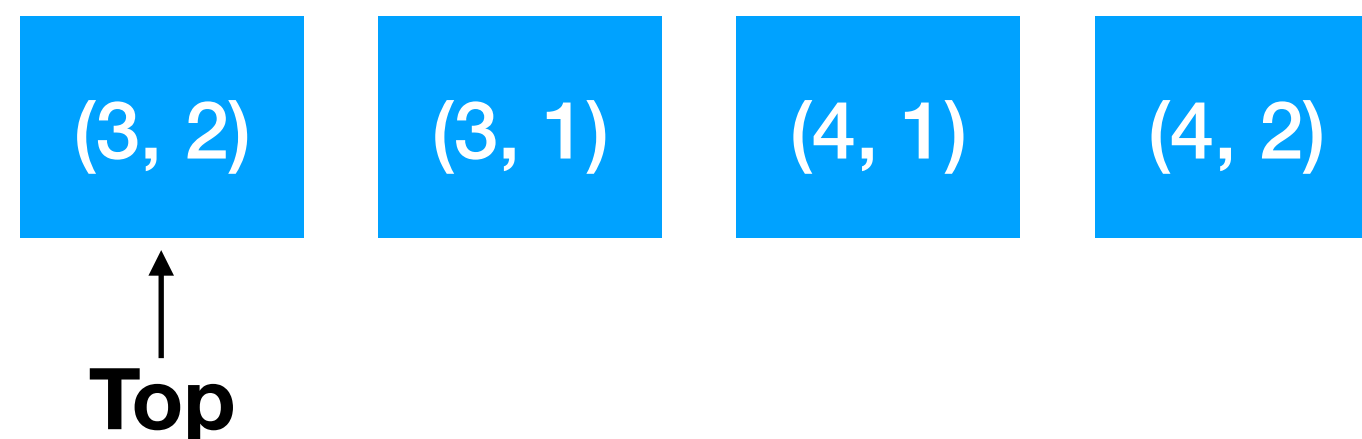


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#		+	#	#
5		#		#	

Diagram illustrating a DFS search on a grid. The grid is 6x6. The search path is highlighted in green, starting from (4, 2) and moving to (3, 2), (3, 1), (2, 1), and (2, 2). The current position is (2, 2), which is highlighted in red. Arrows indicate the movement from (4, 2) to (3, 2), (3, 2) to (3, 1), (3, 1) to (2, 1), and (2, 1) to (2, 2). A blue diagonal banner in the bottom right corner says "Demo".

DFS Example

- For every frontal element, traverse randomly until cannot further
- $\text{top}() = (3, 3);$
- $\text{pop}()$



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#		+	#	#
5		#		#	

Demo

DFS Example

- For every frontal element, traverse randomly until cannot further
- $\text{top}() = (3, 2);$
- $\text{pop}()$

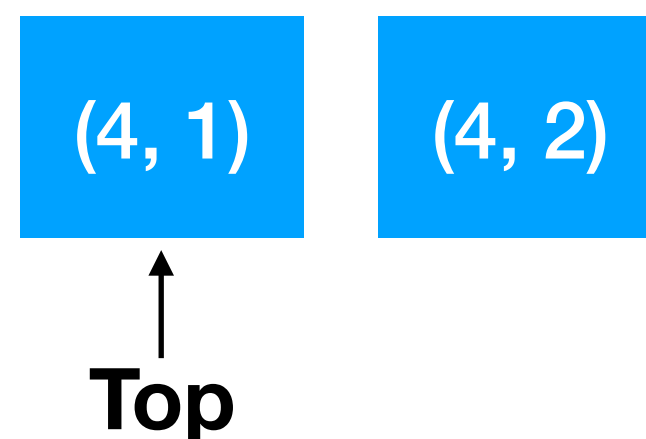


	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#		+	#	#
5		#		#	

Demo

DFS Example

- For every frontal element, traverse randomly until cannot further
- $\text{top}() = (3, 1);$
- $\text{pop}()$



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#		+	#	#
5		#		#	

Demo

DFS Example

- For every frontal element, traverse randomly until cannot further
- $\text{top}() = (4, 1);$
- $\text{pop}()$

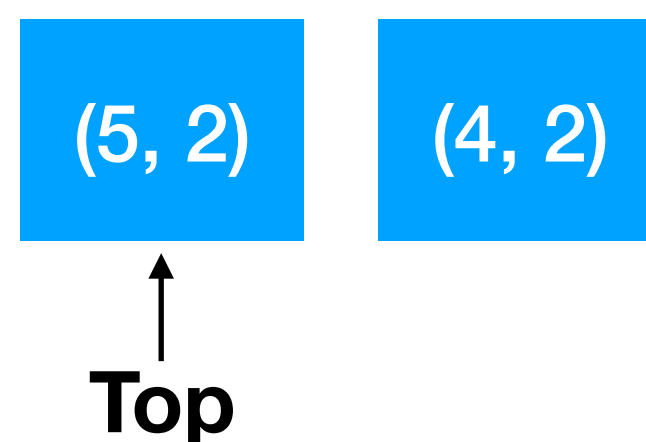
(4, 2)
↑
Top

	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#		+	#	#
5		#		#	

Demo

DFS Example

- For every frontal element, traverse randomly until cannot further
- $\text{top}() = (4, 2);$
- $\text{push}(5, 2)$



	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#			#	#
5		#		#	

Demo

DFS Example

- For every frontal element, traverse randomly until cannot further
- $\text{top}() = (5, 2);$
- $\text{pop}()$

(4, 2)
↑
Top

	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#			#	#
5		#		#	

Demo

DFS Example

- For every frontal element, traverse randomly until cannot further
- $\text{top}() = (4, 2);$
- $\text{pop}()$

↑
Top

	0	1	2	3	4
0	*	#			
1	#	#	#	#	#
2	#		#		#
3	#				#
4	#			#	#
5		#		#	

Demo

DFS Example

DFS Example

- Assuming map size $N \times M$, what is the time and space complexity?

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DFS Example

- Assuming map size $N \times M$, what is the time and space complexity?
- Time complexity: $O(N \times M)$
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- Space complexity: $O(N \times M)$

DFS Example

- Assuming map size $N \times M$, what is the time and space complexity?
- Time complexity: $O(N \times M)$
 - Is it $\Theta(N \times M)$ as well? Why?
- Space complexity: $O(N \times M)$
 - Is it $\Theta(N \times M)$ as well? Why?

Summary

- Lecture 7: Data Structure 1-4
 - Linked List
 - Stack
 - Queue
 - Θ and Big-O Notation
 - BFS, DFS