CSCI 150
Introduction to Digital and Computer System Design
Midterm Review I

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Overview

- Focus: Review
- Architecture: Combinational Logic Circuit
- Textbook v4: Ch1-4; v5: Ch1-3
- Core Ideas:
  1. Digital Information Representation (Lecture 1)
  2. Combinational Logic Circuits (Lecture 2)
  3. Combinational Functional Blocks, Arithmetic Blocks (Lecture 3)

Lecture 1: Digital Information Representation

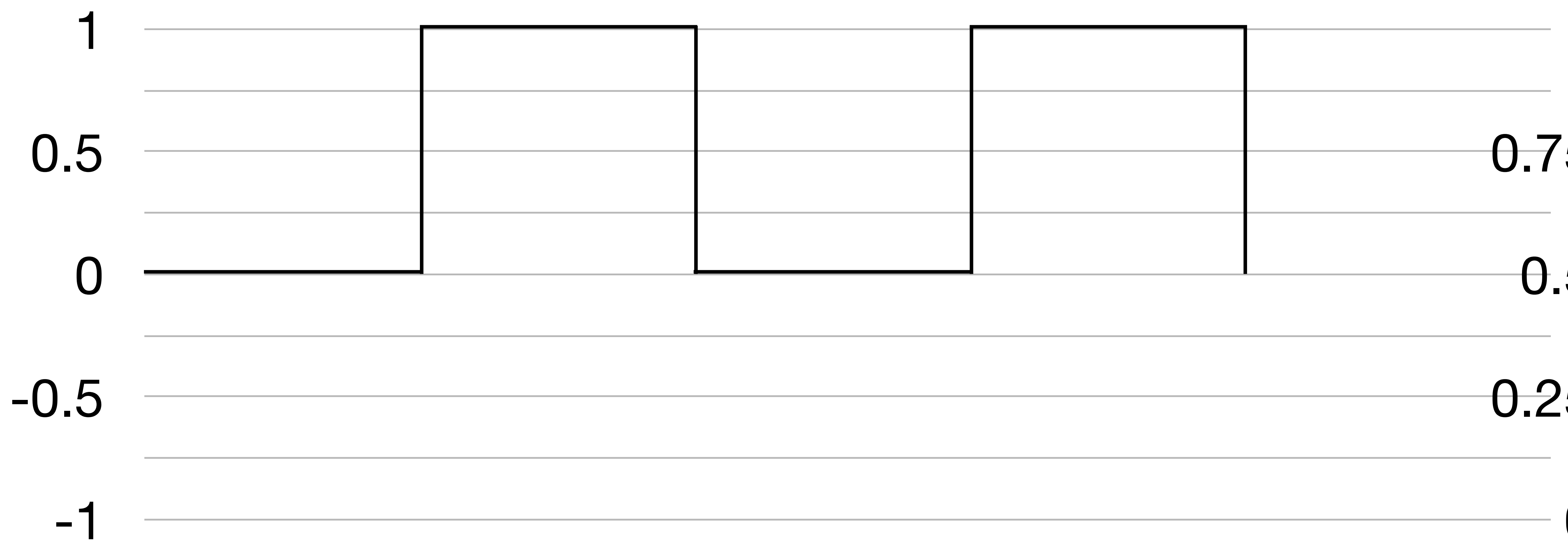
Analog vs Digital circuits; Modern computer architectures; Embedded systems;
Number Systems; Conversions;
Arithmetic Operations; Alphanumeric Codes

Summary

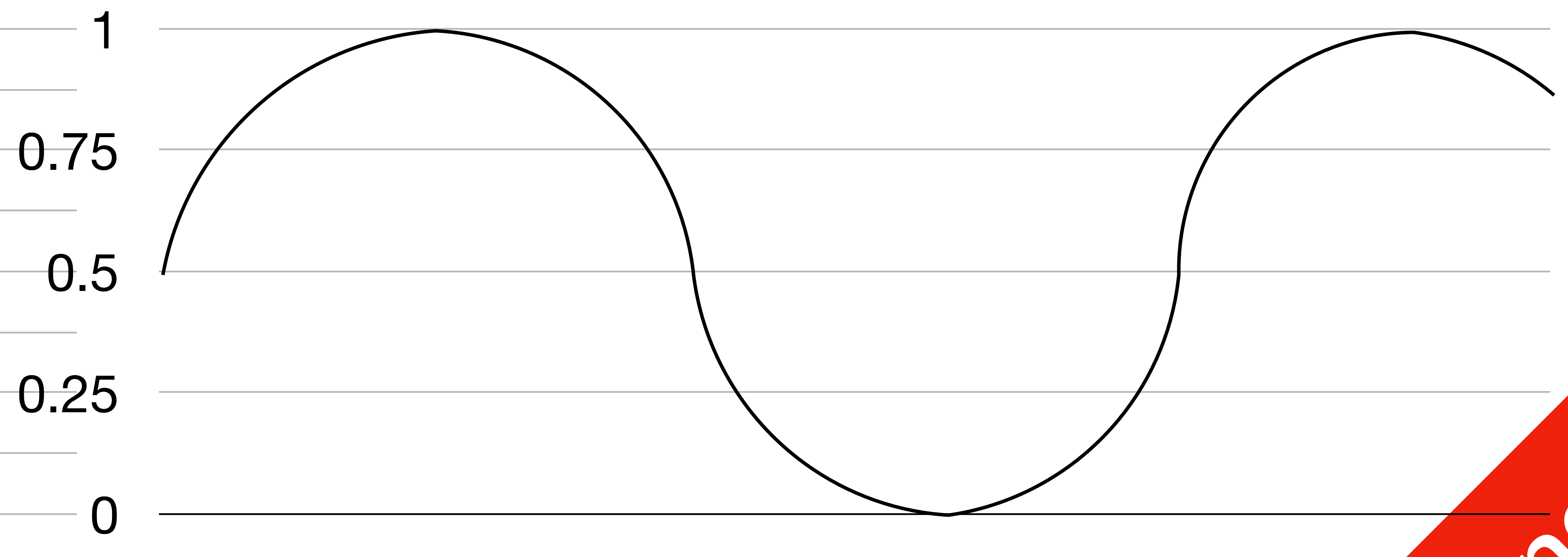
P1.1
Digital Rep.

Analog vs Digital circuits

- Digital Circuits
  - Process digital signals
  - Current/Voltage represent discrete logical and numeric values



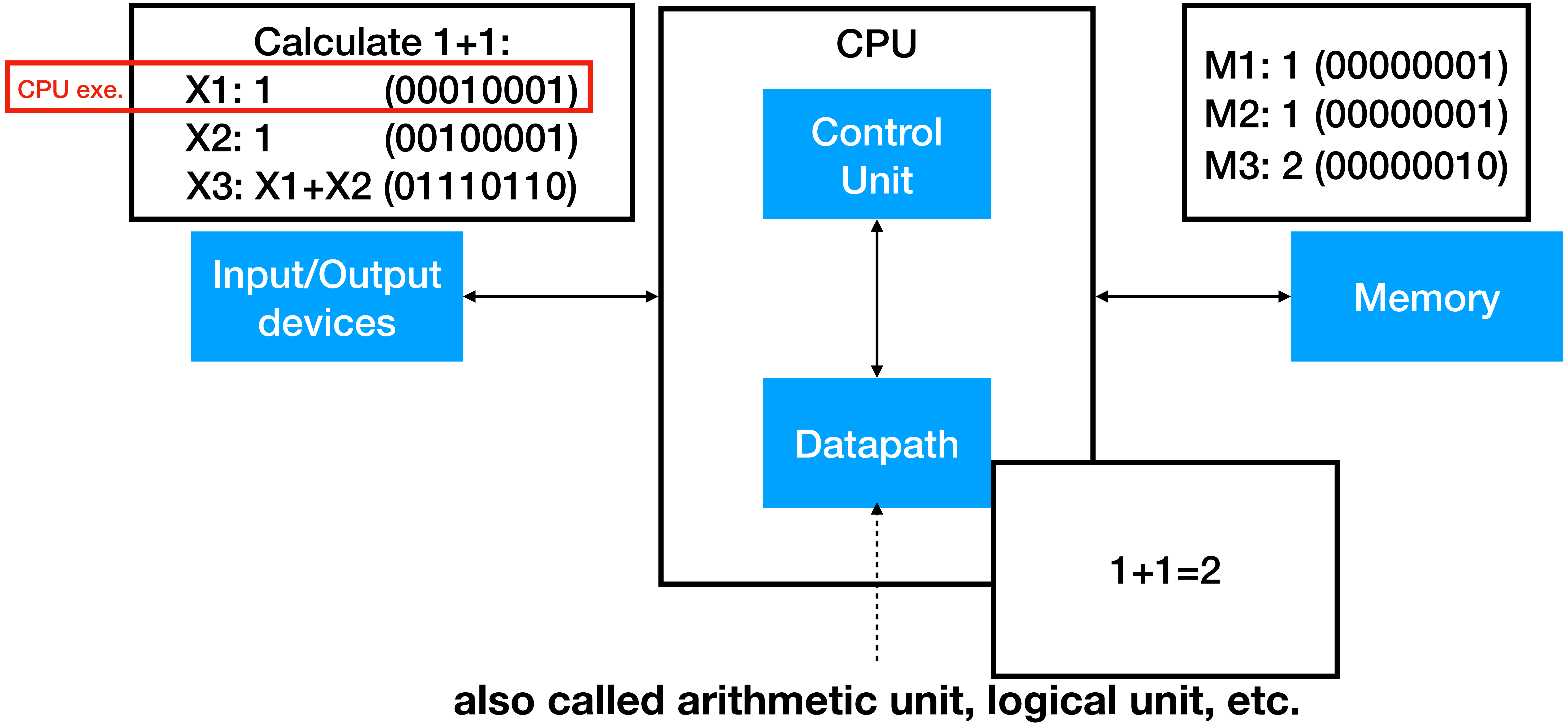
- Analog Circuits
  - Process analog signals
  - Current/Voltage vary continuously to represent information



Concept

Von Neumann Architecture

A very rough example



Demo

Computer

What's it like compared to a human?

- Input/Output devices
  - Interaction (Mouth, hands and feet, eyes, etc.)
- CPU + Memory
  - Processing information, thinking (Brain, short-term memory)
- Storage?
  - Part of I/O devices (Books, long-term memory)

1. Von Neumann Architecture

Embedded Systems

- Similar to computers: processes information
- Difference
  - Function is usually simpler, and very very specific
  - Not programmable

Concept

Decimal System

7	2	4	.	0	5
2	1	0		−1	−2

- Numbers as strings of digits, each ranging from 0-9
- The decimal system is of base(radix) 10

Review

Decimal System

$$\begin{matrix} 7 & 2 & 4. & 0 & 5 \\ 2 & 1 & 0 & -1 & -2 \end{matrix}$$

$$= 7 \times 10^{2} + 2 \times 10^{1} + 4 \times 10^{0} + 0 \times 10^{-1} + 5 \times 10^{-2}$$

- Numbers as strings of digits, each ranging from 0-9
- The decimal system is of base(radix) 10

Review

Numbers of base N

- Default base: 10
- When there are numbers represented in different bases, attach base
  - Decimal: 754.05 -> $(754.05)_{10}$
  - e.g. Base 5: $(432.1)_5 = ?$

$$= 4 \times 5^2 + 3 \times 5^1 + 2 \times 5^0 + 1 \times 5^{-1} = (117.2)_{10}$$

Concept

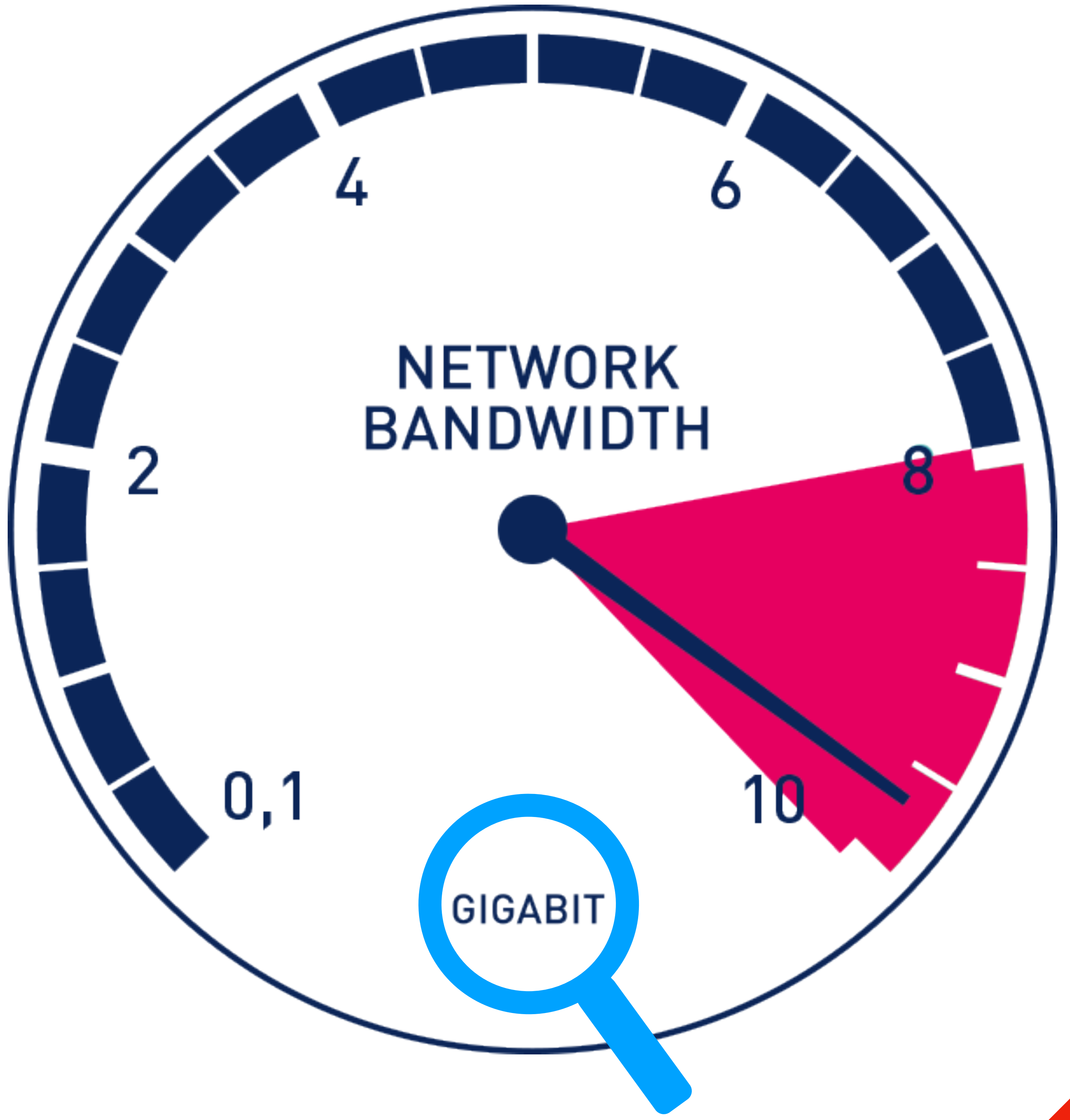
Binary Systems in Computers

- Every 8bit is called a Byte
- 1,024 = 2^{10} is called K (Kilo)
- 1,024 x 1,024 = 2^{20} is called M (Mega)
- 1,024 x 1,024 x 1,024 = 2^{40} is called G (Giga)
- Tera, Peta, Exa, Zetta, Yotta

P1.2
Number Systems

Binary Systems in Computers

- What is the difference between MBps and Mbps?
  - MegaBytes per second vs MegaBits per second
  - 8x difference!



Concept

Octal and Hexadecimal Systems

- Octal: base 8
  - digits: 0-7
- Hexadecimal: base 16
  - digits: 0-9, A-F (10-15)

Concept

P1.2
Number Systems

Conversions

10	9	8	7	6	5	4	3	2	1
1024	512	256	128	64	32	16	8	4	2

- Binary-to-
  Octal: 3bits per octal digit
  Hexadecimal: 4bits per hexa digit
  Decimal: use the chart

- Decimal-to-
  Binary: use the chart
  Oct/Hex: do binary first

Concept

Arithmetics

- The same as decimal (mostly)
-

$$\begin{array}{r} 0010 \\ +0011 \\ \hline 0101 \end{array} \qquad \begin{array}{r} 0101 \\ -0011 \\ \hline 0010 \end{array}$$

Example (binary)

Concept

Arithmetics

OCTAL Multiplication

Octal	Octal	Decimal
$\begin{array}{r} 762 \\ \times \quad 54 \\ \hline 4672 \\ 3710 \\ \hline 43772 \end{array}$	$5 \times 2 = 12$	$10 = (12)_8$
	$5 \times 6 + 1 = 37$	$31 = (37)_8$
	$5 \times 7 + 3 = 46$	$38 = (46)_8$
	$\cdots$	$\cdots$

P1.4 Representations

Signed & Unsigned Integers

- Unsigned 8bit:
  - $(11111111)_2 = 255$
- Signed 8bit (only in digital circuits):
  - 127 -> '01111111'
  - -127 -> '11111111'

First digit:
- **0 for positive**
- **1 for negative**

10001111

(binary, 8bit, signed)

Concept

P1.4
Representations

Signed & Unsigned Integers

- Unsigned 8bit integer: 0 - 255
  - Signed 8bit integer: -128 - 127
- Unsigned 32bit integer: 0 - 4,294,967,295
  - Signed 32bit integer: -2,147,483,648 - 2,147,483,647
- Unless otherwise specified, treat as unsigned

Concept

Binary Coded Decimal

- Decimal numbers, each digit represented in 4bit binary, but separately
- 185 = $(0001\ 1000\ 0101)_{BCD}$ = $(10111001)_2$
- Used in places where using decimals directly is more convenient, such as digital watches etc.

Concept

ASCII

- American Standard Code for Information Interchange
- Assign each character with a 8bit binary code (e.g. '0'-'9', 'A'-'Z', 'a'-'z')
- The first bit is always 0

Demo

Parity Code

- For error detection in data communication
  - e.g. resulting from packet loss or other forms of interference
- One parity bit for n-bits
  - An extra even parity bit: whether the number of 1s is not even
  - An extra odd parity: whether the number of 1s is not odd
  - Can be placed in any fixed position
  - Does it always work?

Demo

Parity Code

Original 7bits	with Even parity	with Odd parity
1000001	<u>0</u>1000001	<u>1</u>1000001
1010100	<u>1</u>1010100	<u>0</u>1010100

Demo

Circuits

- Circuits
  - Digital and Analog
- Integrated systems
  - Von Neumann computers
  - Embedded systems

Review

P1.5
Lect 1 Summary

Number Systems

- Number systems of base N
- Binary systems
- Octal and Hexadecimal systems
- Arithmetics

Review

P1.5
Lect 1 Summary

Number Systems in DC

- Bit, Byte, Representation ranges
- Signed and Unsigned Binary Integers
- BCD, ASCII, UTF8
- Parity bit

Review

P1.5
Lect 1 Summary

Digital to Analog Conversion

- Frequency: number of cycles per second
- Sample rate: number of samples per unit time
- Bitrate: number of bits per second

Review

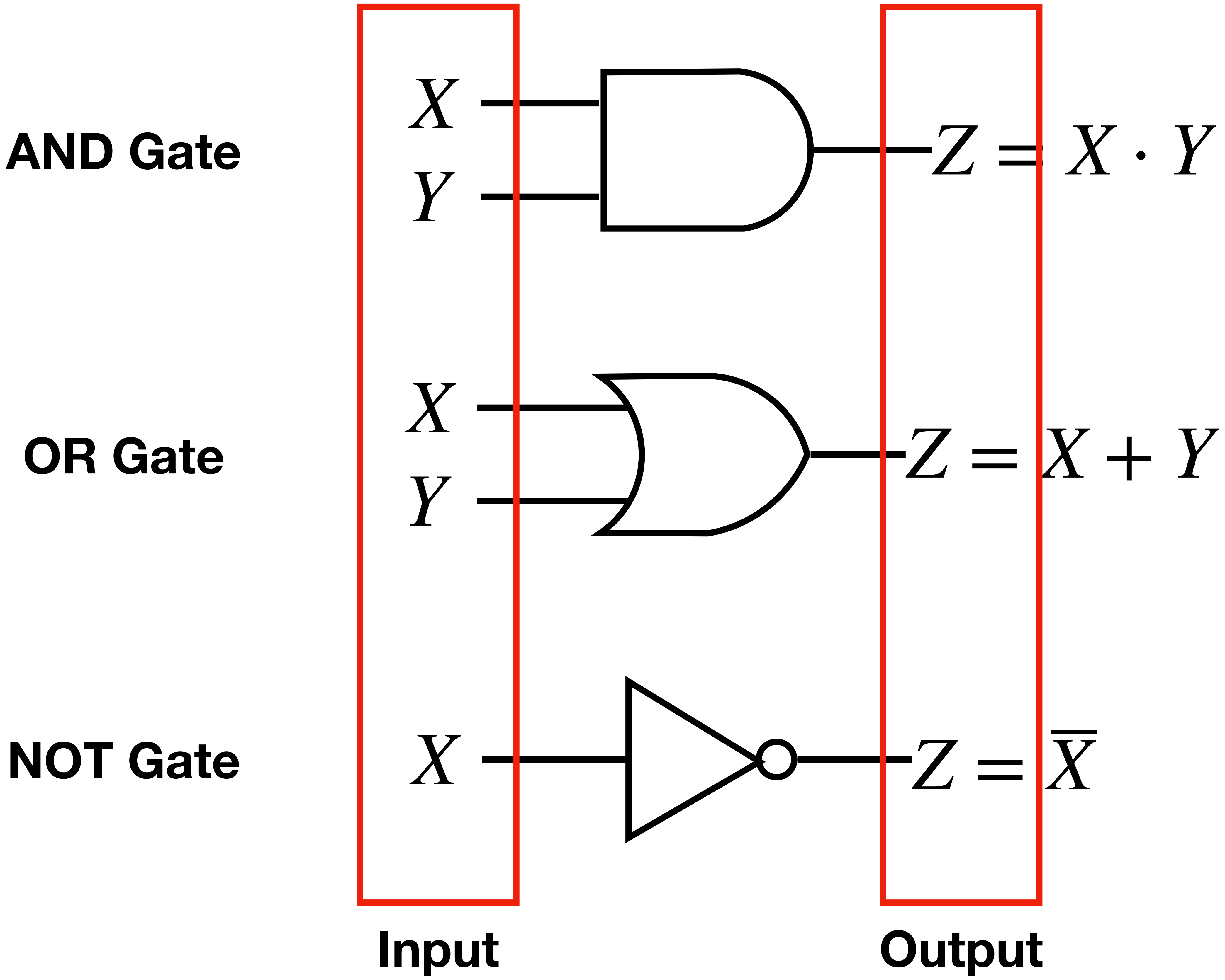
Lecture 2: Combinational Logic Circuits

Logic Gates; Boolean Algebra; Minterm/Maxterm; K-Map; Some Other Gate Types

Summary

P2.1
Logic Gates

First 3 Gates

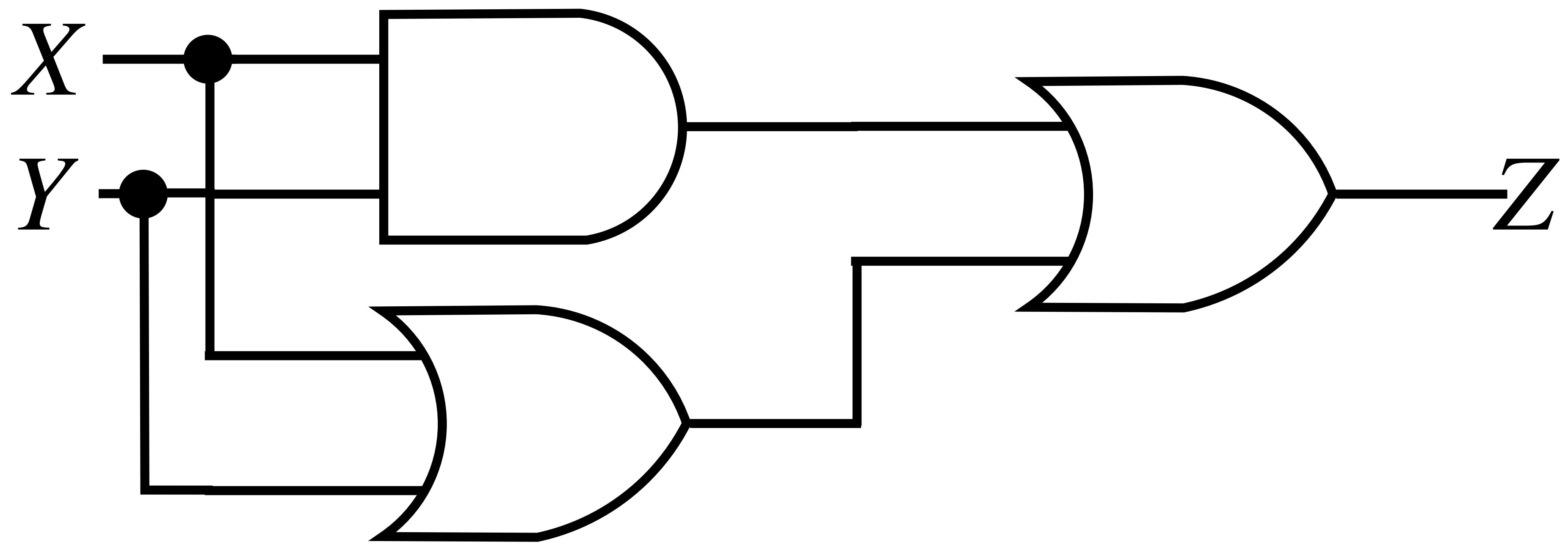


Concept

Truth Table

Truth Table

X	Y	$Z = (X \cdot Y) + (X + Y)$
0	0	0
0	1	1
1	0	1
1	1	1



Example

Basic Identities

- Boolean Algebra solving
  - **Identify** rules **applicable** to the expression
  - **Apply** rules that can help you **simplify** the expression
    - **Simplification**: reducing the number of variables and operators in an expression without changing it’s truth table values
    - **Atomic element:** an element that can’t have the number of its variables and operators reduced any further

Basic Identities

1. $X + 0 = X$
2. $X \cdot 1 = X$
3. $X + 1 = 1$
4. $X \cdot 0 = 0$
5. $X + X = X$
6. $X \cdot X = X$
7. $X + \overline{X} = 1$
8. $X \cdot \overline{X} = 0$
9. $\overline{\overline{X}} = X$

Concept

Basic Identities

- Communicative

10. $X + Y = Y + X$

11. $XY = YX$

- Associative

12. $X + (Y + Z) = (X + Y) + Z$

13. $X(YZ) = (XY)Z$

- Distributive

14. $X(Y + Z) = XY + XZ$

15. $X + (YZ) = (X + Y)(X + Z)$

- DeMorgan’s

16. $\overline{X + Y} = \overline{X} \cdot \overline{Y}$

17. $\overline{X \cdot Y} = \overline{X} + \overline{Y}$

Concept

Basic Identities

A. $X + XY = X$

B. $XY + X\overline{Y} = X$

C. $X + \overline{X}Y = X + Y$

D. $X(X + Y) = X$

E. $(X + Y)(X + \overline{Y}) = X$

F. $X(\overline{X} + Y) = XY$

Complementation

- $\overline{F}$: complement (invert) representation for a function F, obtained from an interchange of 1s to 0s and 0s to 1s for the values of F in the truth table

- Apply DeMorgan's Rule

16. $\overline{X_1 + X_2 + \ldots + X_n} = \overline{X_1} \cdot \overline{X_2} \cdot \ldots \cdot \overline{X_n}$

17. $\overline{X_1 \cdot X_2 \cdot \ldots \cdot X_n} = \overline{X_1} + \overline{X_2} + \ldots + \overline{X_n}$

Concept

P2.2
Boolean Algebra

Algebraic Manipulation

Difficulty: Simple

Simplify the following expressions

- $\overline{X} \cdot \overline{Y} + XYZ + \overline{X}Y$
- $X + Y(Z + \overline{X + Z})$

Exercise

P2.2
Boolean Algebra

Algebraic Manipulation

Difficulty: Mid

Simplify the following expressions

- $\overline{W}X(\overline{Z} + \overline{Y}Z) + X(W + \overline{W}YZ)$
- $(AB + \overline{A}\overline{B})(\overline{C}\overline{D} + CD) + AC$

Exercise

Algebraic Manipulation

Difficulty: Mid

Simplify the following expressions

- $\overline{A} \cdot \overline{C} + \overline{A}BC + \overline{B}C$
- $\overline{A + B + C} \cdot \overline{ABC}$

Exercise

P2.2
Boolean Algebra

Algebraic Manipulation

Difficulty: Mid

Simplify the following expressions

- $AB\overline{C} + AC$
- $\overline{A} \cdot \overline{B}D + \overline{A} \cdot \overline{C}D + BD$

Exercise

Algebraic Manipulation

Difficulty: HARDCORE

Prove the identity of each of the following Boolean equations

- $AB\overline{C} + B\overline{C} \cdot \overline{D} + BC + \overline{C}D = B + \overline{C}D$
- $WY + \overline{W}Y\overline{Z} + WXZ + \overline{W}X\overline{Y} = WY + \overline{W}X\overline{Z} + \overline{X}Y\overline{Z} + X\overline{Y}Z$
- $A\overline{D} + \overline{A}B + \overline{C}D + \overline{B}C = (\overline{A} + \overline{B} + \overline{C} + \overline{D})(A + B + C + D)$

Standard Forms

- Equivalent expressions can be written in a variety of ways
  Standard forms: typical such ways that incorporates some **unique characteristics** -> **simplify the implementation** of these designs

  - **Product terms** (AND terms): e.g. $\overline{X}YZ$
    Literals with inverts connected through only AND operators

  - **Sum terms** (OR terms): e.g. $X + \overline{Y} + Z$
    Literals with inverts connected through only OR operators

Concept

Minterms and Maxterms

- **Minterm**
  Product term; Contains **all variables**; Has only **one Positive row** in the truth table

	X	Y	$m_0 = \overline{X}\overline{Y}$	$m_1 = \overline{X}Y$	$m_2 = X\overline{Y}$	$m_3 = XY$
$(00)_2$=0	0	0	**1**	0	0	0
$(01)_2$=1	0	1	0	**1**	0	0
$(10)_2$=2	1	0	0	0	**1**	0
$(11)_2$=3	1	1	0	0	0	**1**

Concept

Minterms and Maxterms

- **Maxterm**
  Sum term; Contains **all variables**; Has only **one Negative row** in the truth table

$$M_i = \overline{m_i}$$

X	Y	$M_0 = X + Y$	$M_1 = X + \overline{Y}$	$M_2 = \overline{X} + Y$	$M_3 = \overline{X} + \overline{Y}$
0	0	**0**	1	1	1
0	1	1	**0**	1	1
1	0	1	1	**0**	1
1	1	1	1	1	**0**

Concept

P2.3
Standard Forms

Minterms and Maxterms

- e.g. $M_3 = X + \overline{Y} + \overline{Z} = \overline{\overline{X}YZ} = \overline{m_3}$

- Sum of Minterms

  - e.g. $F = \overline{X}\,\overline{Y}\,\overline{Z} + \overline{X}Y\overline{Z} + X\overline{Y}Z + XYZ = m_0 + m_2 + m_5 + m_7$
    $= \Sigma m(0,2,5,7)$

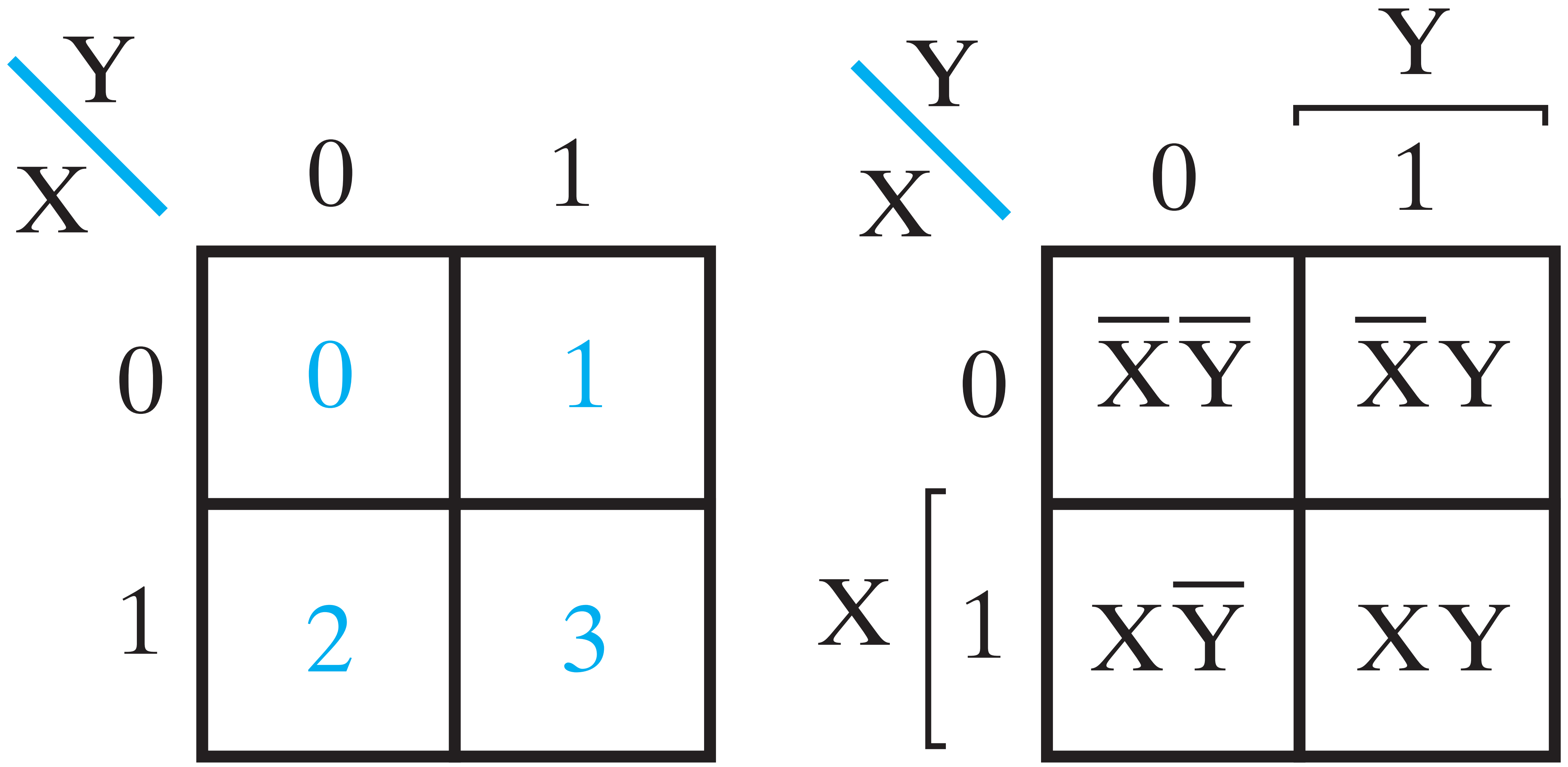
- Product of Maxterm

  - e.g. $F = (X + Y + Z)(X + \overline{Y} + Z)(\overline{X} + Y + \overline{Z})(\overline{X} + \overline{Y} + \overline{Z})$
    $= M_0 M_2 M_5 M_7$
    $= \Pi M(0,2,5,7)$

Concept

P2.4
K-Map

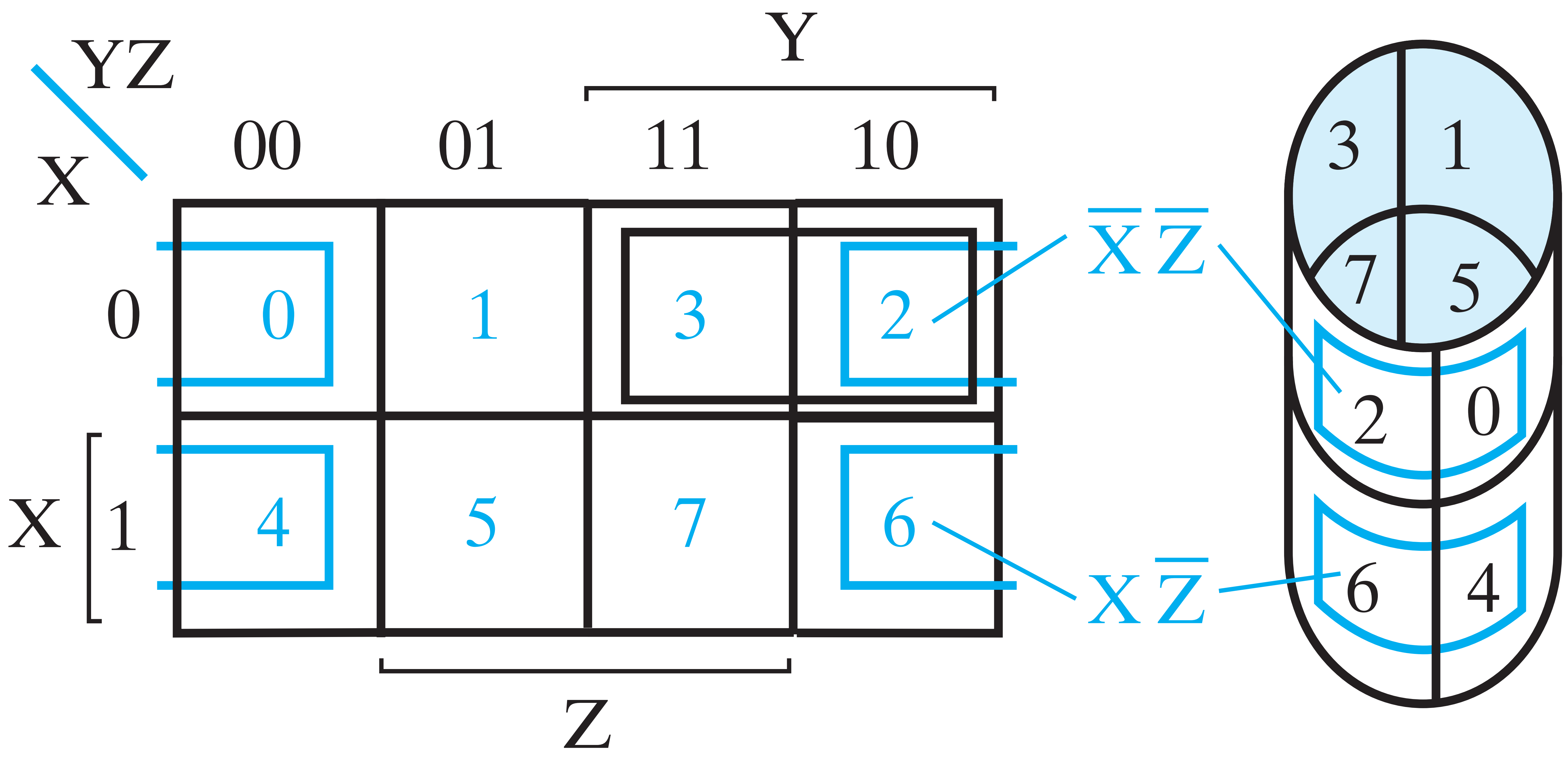
Two Variable Maps



- Number of squares in each map is equal to the number of minterms for the same number of variables, light blue digit above is the index (of minterm)
- Two squares are adjacent if they only differ in one variable
- Binary value inside at each position indicates the truth table value for that term

Concept

Three Variable Maps



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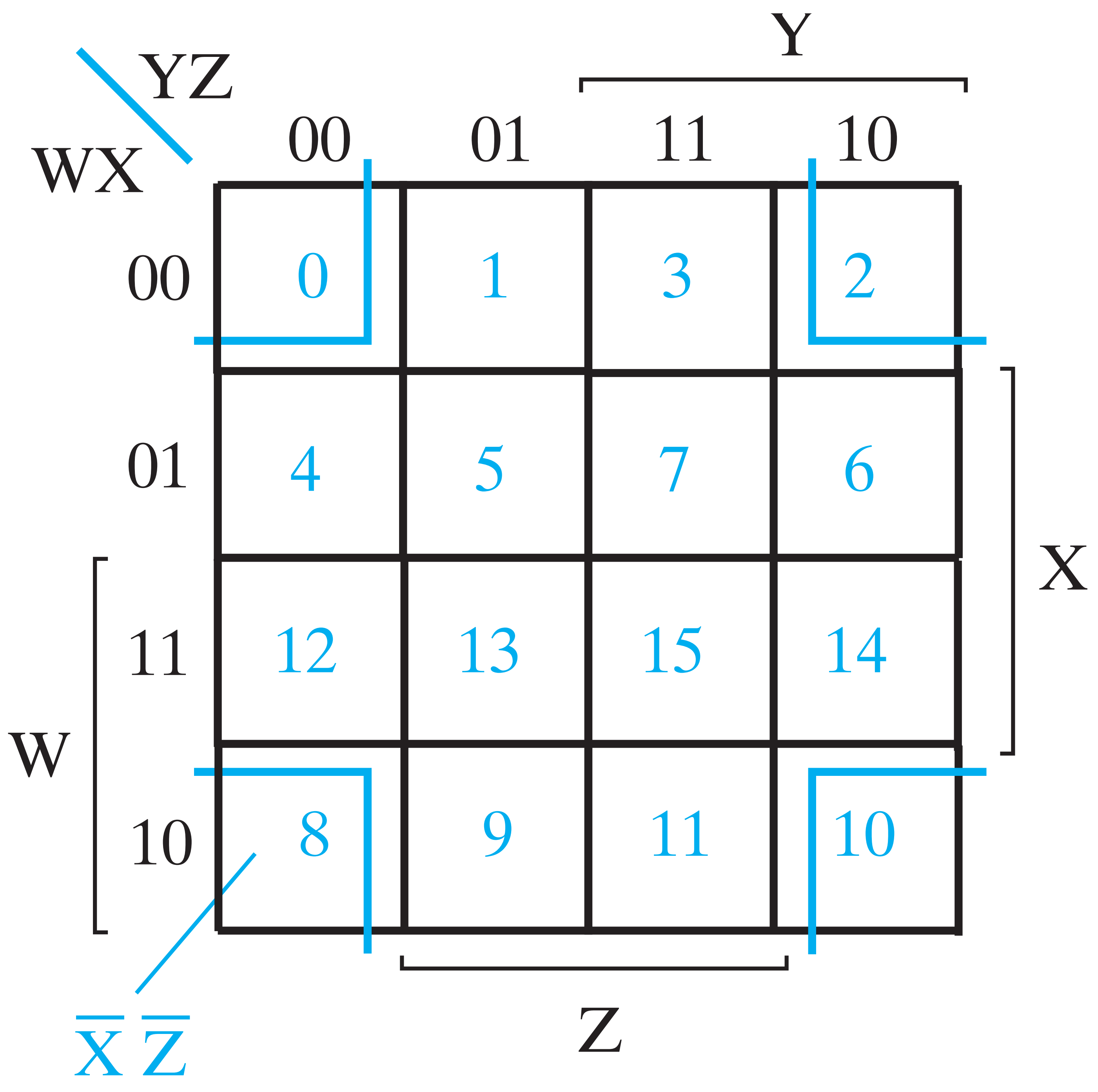
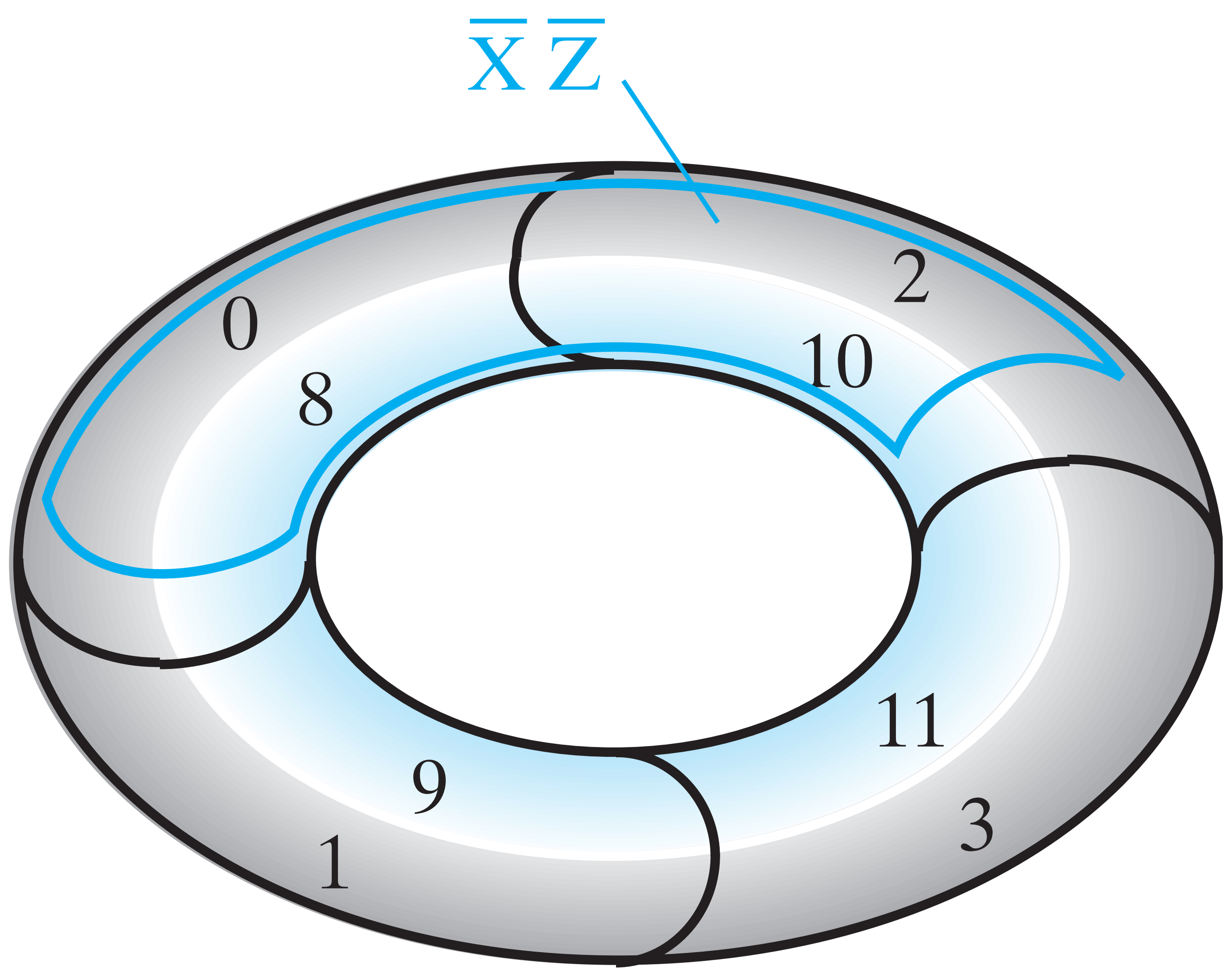
- Two squares are adjacent if they only differ in one variable
- Binary value inside at each position indicates the truth table value for that term

Concept

Four Variable Maps



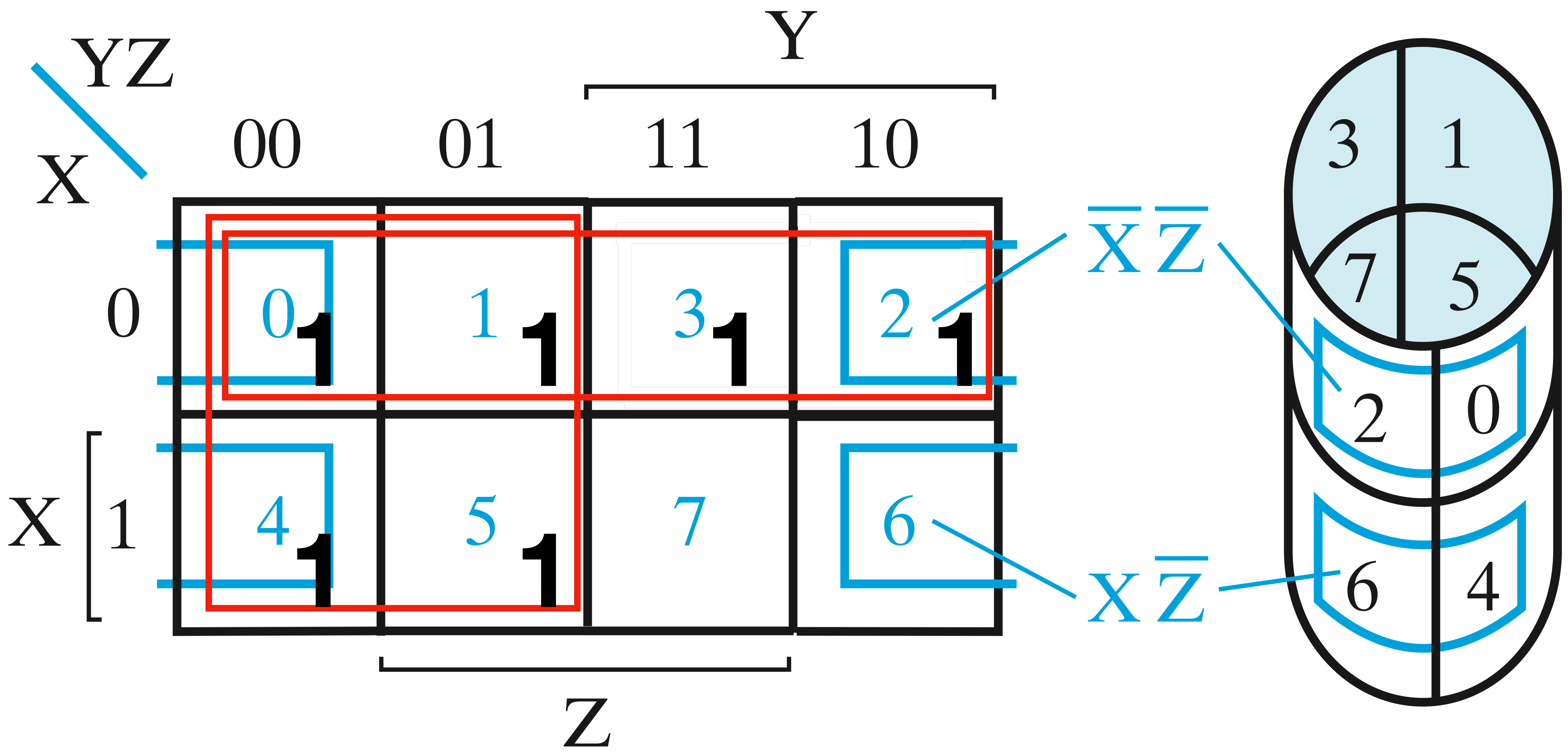


- Number of squares in each map is equal to the number of minterms for the same number of variables, light blue digit above is the index (of minterm)
- Two squares are adjacent if they only differ in one variable
- Binary value inside at each position indicates the truth table value for that term

Concept

P2.4
K-Map

K Map Optimisation



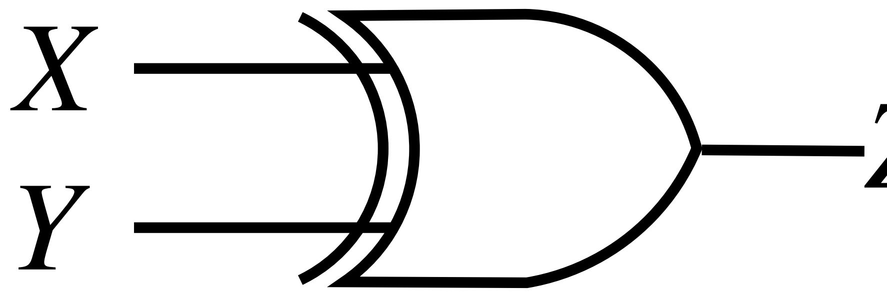
$$F(X,Y,Z) = \Sigma m(0,1,2,3,4,5)$$
$$= \overline{X} + \overline{Y}$$

- Step 1: Enter the values
- Step 2: Identify the set of **largest** rectangles in which **all values are 1**, covering **all** 1s
- Step 3: **Read off** the selected rectangles. Rectangle has to have 2^k length edges

Concept

XOR Gate

XOR Gate
Exclusive-OR



$$Z = X \oplus Y$$
$$= X\overline{Y} + \overline{X}Y$$

- $X \oplus 0 = X$
- $X \oplus 1 = \overline{X}$
- $X \oplus X = X$
- $X \oplus \overline{X} = 1$
- $X \oplus \overline{Y} = \overline{X \oplus Y}$
- $\overline{X} \oplus Y = \overline{X \oplus Y}$

XOR Truth Table

X	Y	$Z = X \oplus Y$
0	0	0
0	1	1
1	0	1
1	1	0

Concept

XOR Gate

- $X \oplus 0 = X$
- $X \oplus 1 = \overline{X}$
- $X \oplus X = X$
- $X \oplus \overline{X} = 1$
- $X \oplus \overline{Y} = \overline{X \oplus Y}$
- $\overline{X} \oplus Y = \overline{X \oplus Y}$

Concept

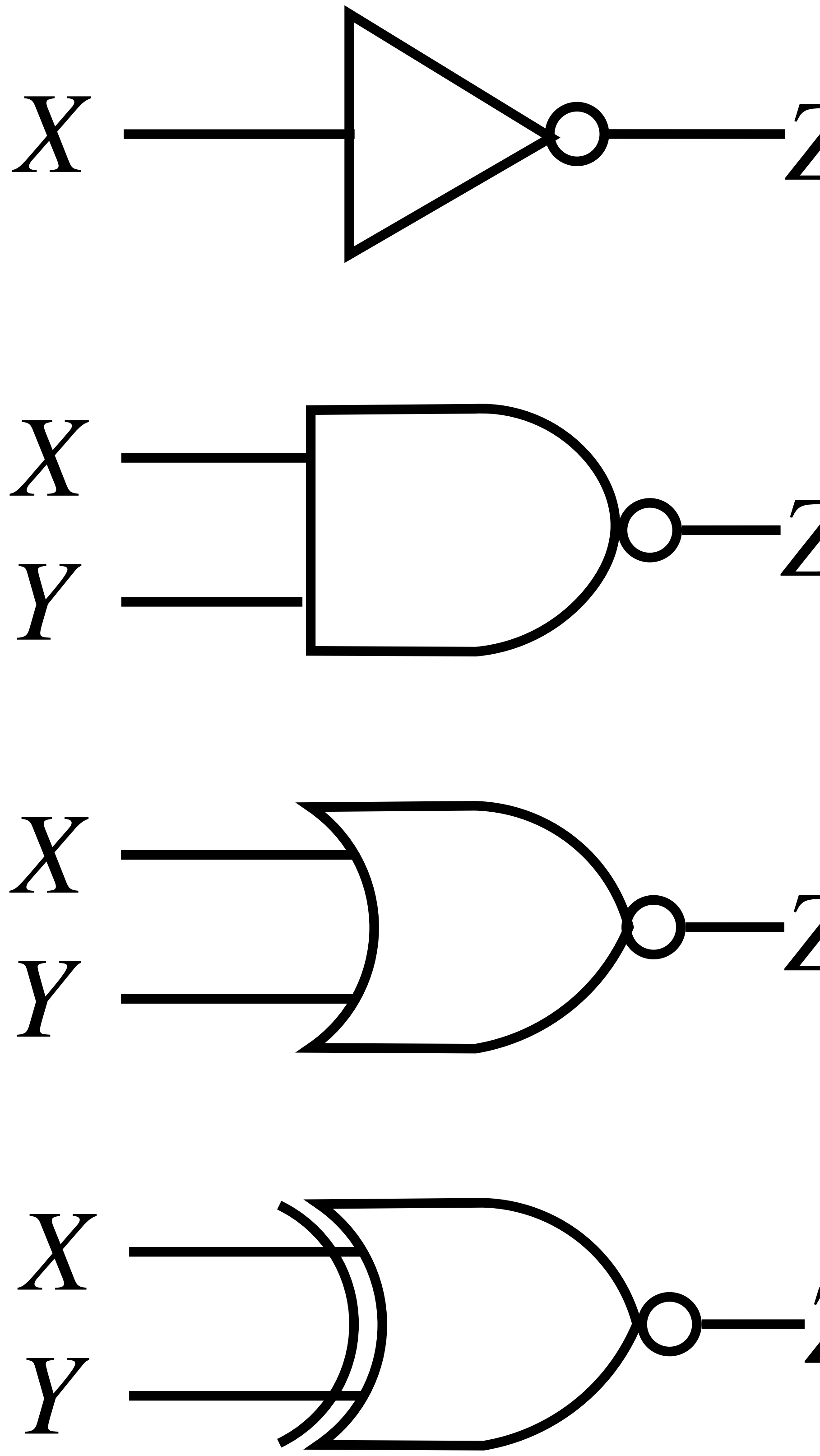
N-Gates

NOT Gate — $Z = \overline{X}$

NAND Gate — $Z = \overline{X \cdot Y}$

NOR Gate — $Z = \overline{X + Y}$

XNOR Gate — $Z = \overline{X \oplus Y}$



Concept

Boolean Algebra

I. AND, OR, NOT Operators and Gates
- Simple digital circuit implementation
- Algebraic manipulation using Binary Identities

II. Standard Forms
- Minterm & Maxterm
- Sum of Products & Product of Sums

III. Optimisation Using K-Map (For 2,3,4 Variables)

IV. XOR, NAND, NOR, XNOR

Review

Lecture 3: Combinational Logic Design

5 Steps Systematic Design Procedures; Functional Blocks; Decoder, Enabler, Multiplexer; Arithmetic Blocks

Summary

Systematic Design Procedures

1. **Specification**: Write a specification for the circuit

2. **Formulation**: Derive relationship between inputs and outputs of the system e.g. using truth table or Boolean expressions

3. **Optimisation**: Apply optimisation, minimise the number of logic gates and literals required

4. **Technology Mapping**: Transform design to new diagram using available implementation technology

5. **Verification**: Verify the correctness of the final design in meeting the specifications

Concept

Hierarchical Design

- "divide-and-conquer"
- Circuit is broken up into individual functional pieces (blocks)
  - Each block has explicitly defined **Interface** (I/O) and **Behaviour**
  - A single block can be **reused** multiple times to simplify design process
  - If a single block is too complex, it can be **further divided into smaller blocks**, to allow for easier designs

Concept

Value-Fixing, Transferring, and Inverting

① **Value-Fixing**: giving a constant value to a wire

- $F = 0; F = 1;$

② **Transferring**: giving a variable (wire) value from another variable (wire)

- $F = X;$

③ **Inverting**: inverting the value of a variable

- $F = \overline{X}$

Concept

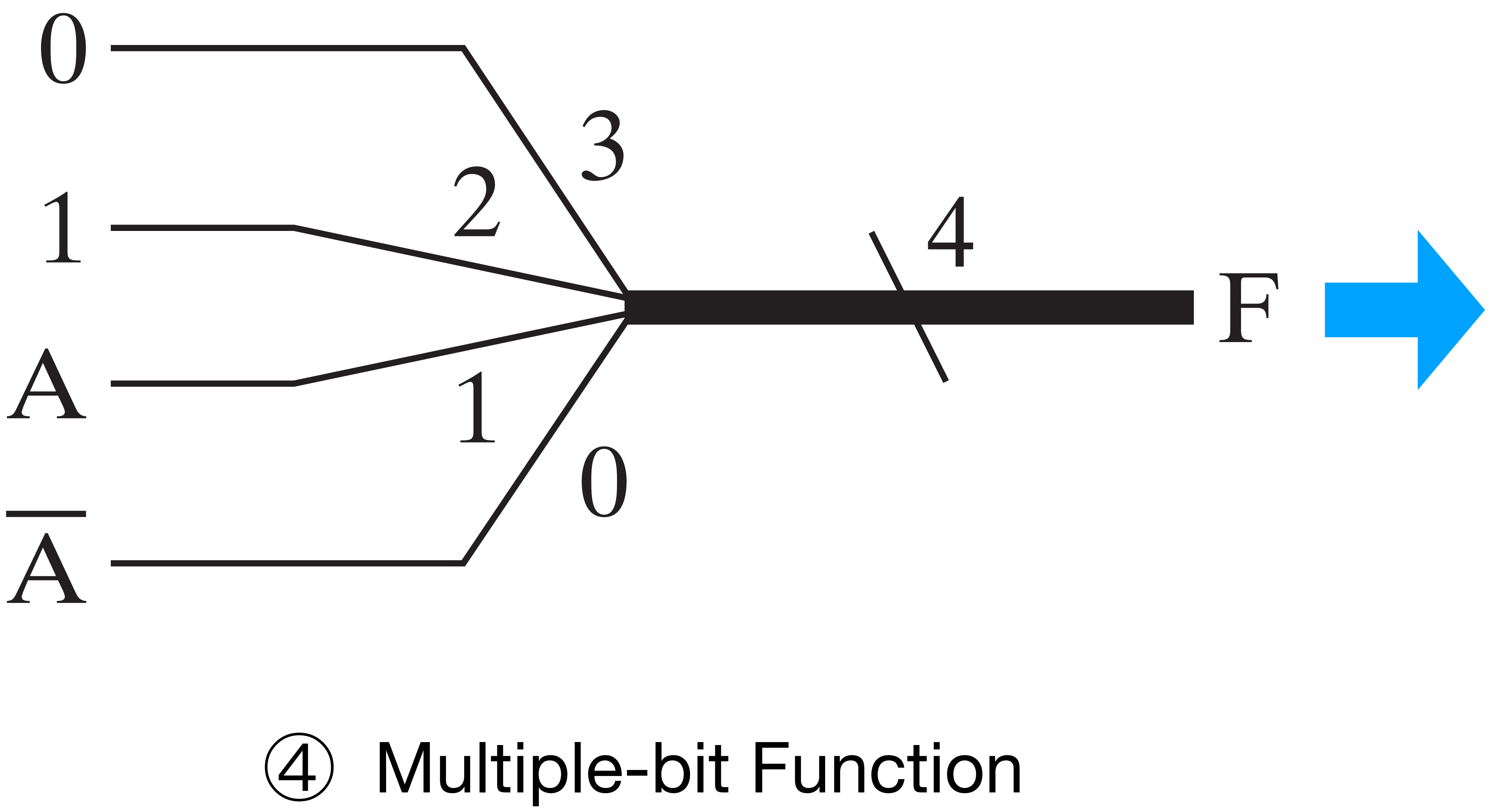
Vector Denotation

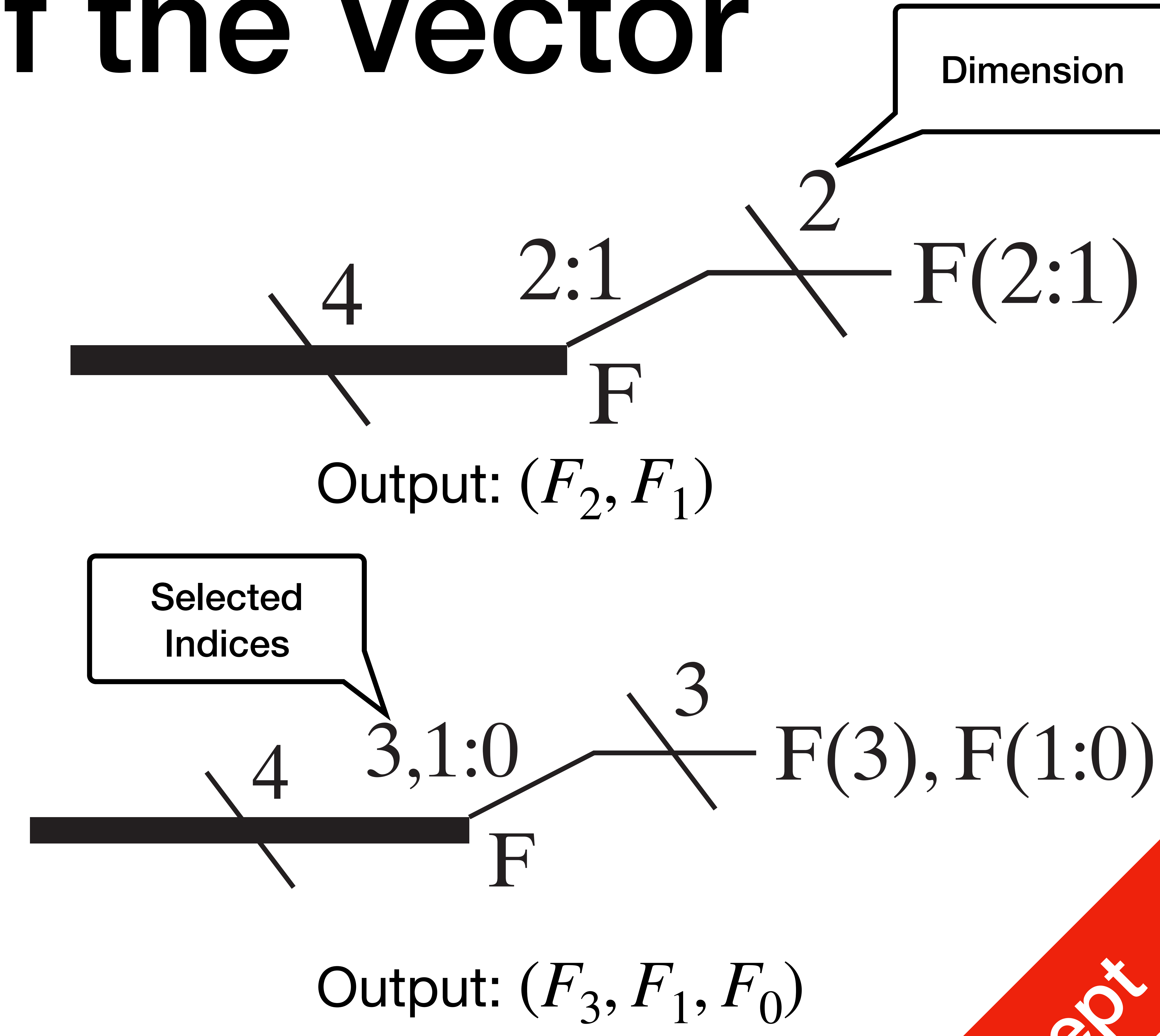
④ **Multiple-bit Function**

- Functions we've seen so far has only one-bit output: 0/1
- Certain functions may have n-bit output
  - $F(n-1:0) = (F_{n-1}, F_{n-2}, \ldots, F_0)$, each F_i is a one-bit function
  - Curtain Motor Control Circuit: $F = (F_{\mathrm{Motor}_1}, F_{\mathrm{Motor}_2}, F_{\mathrm{Light}})$

Concept

Taking part of the Vector





Concept

Enabler

⑤ Enabler

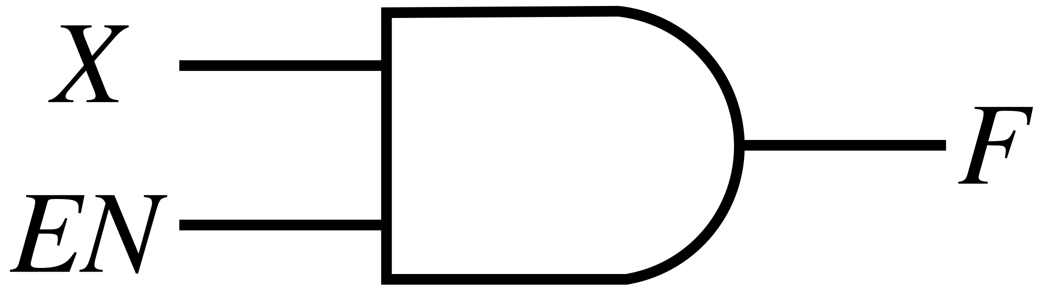
- Transferring function, but with an additional EN signal acting as switch

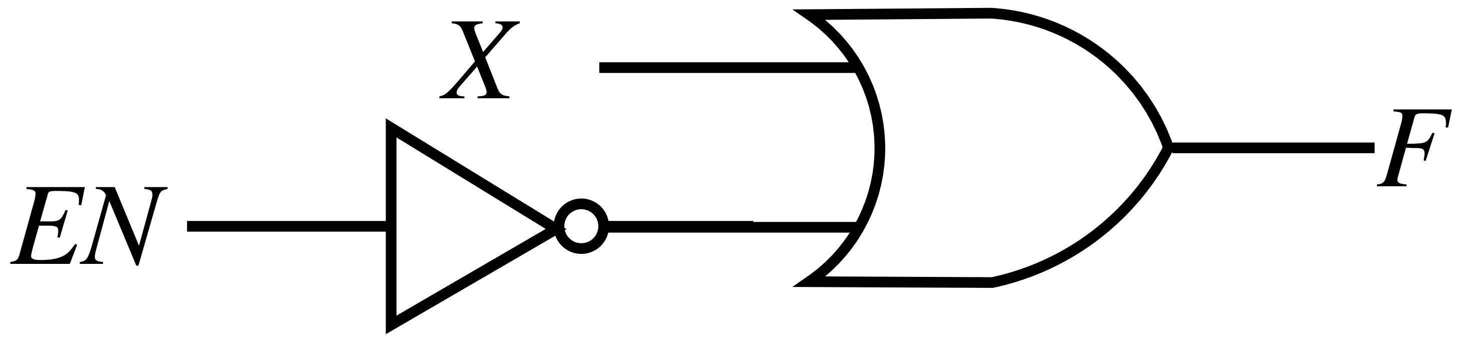
EN	X	F
0	X	0
1	0	0
1	1	1

Enabler

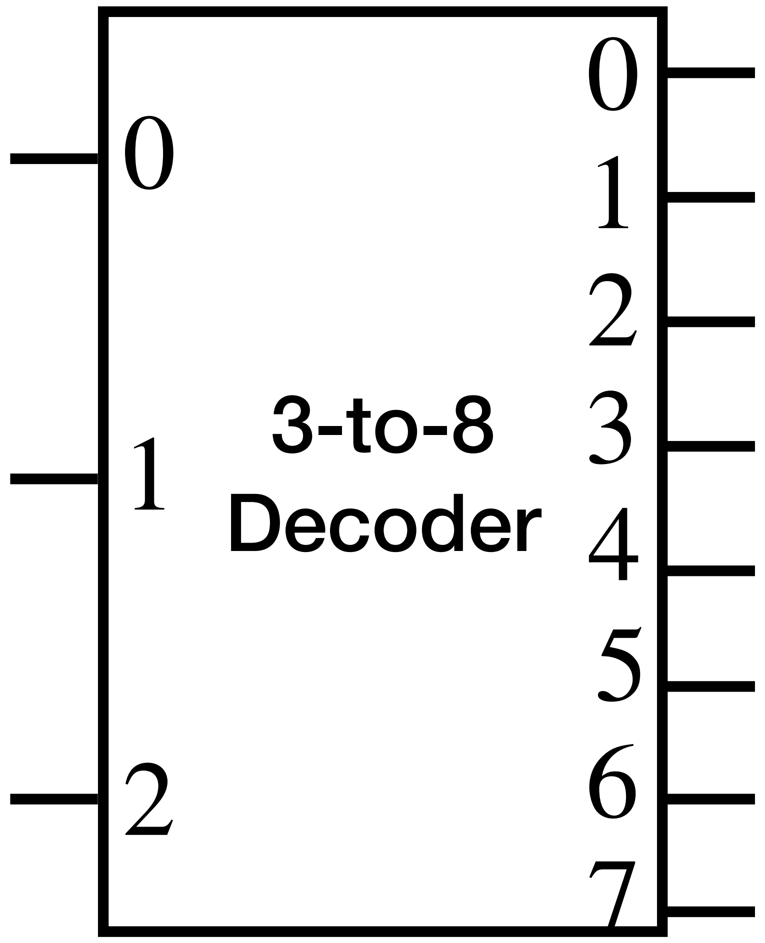
⑤ Enabler

- Transferring function, but with an additional EN signal acting as switch



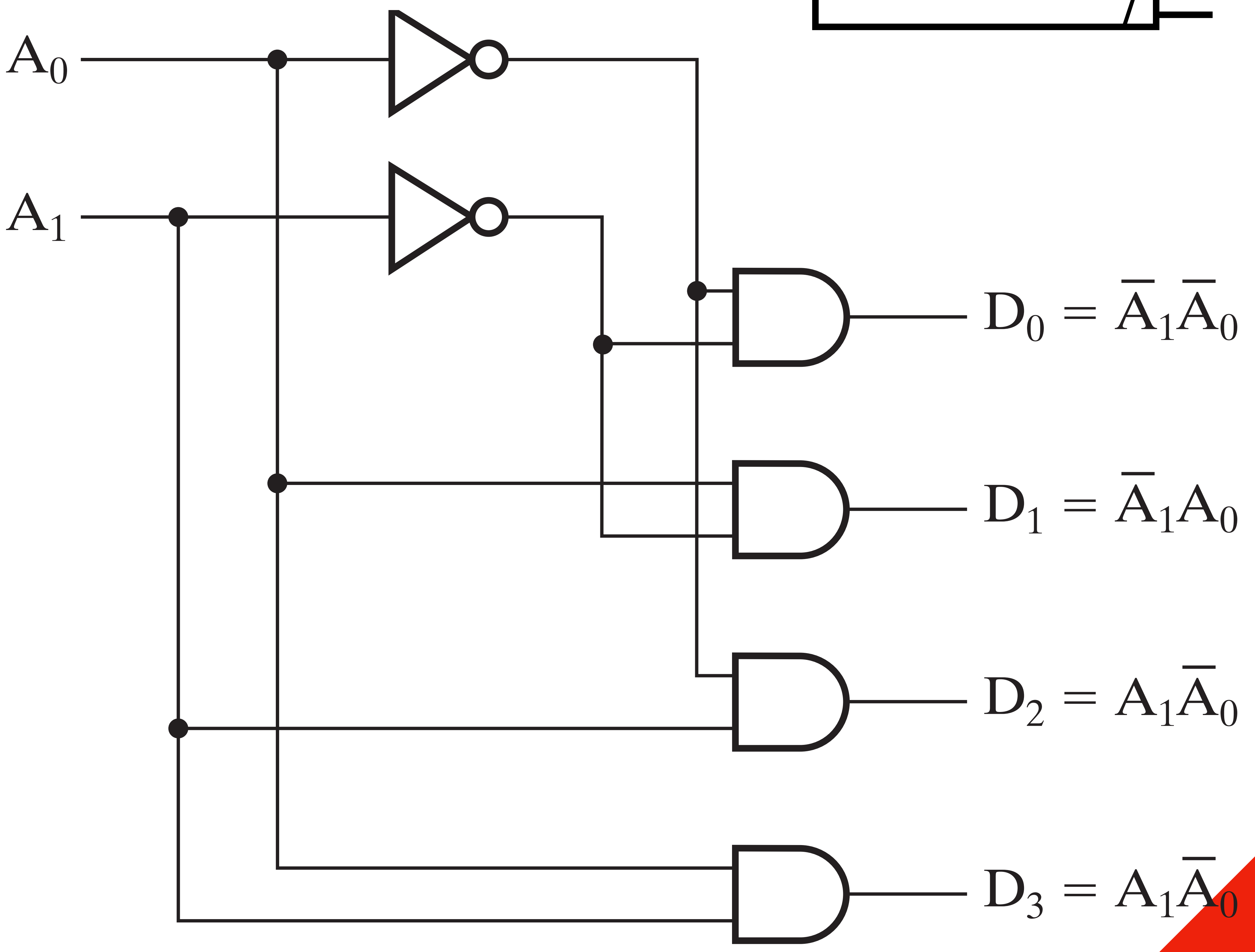


Decoder



- n-bit input, 2^nbits output
  - $D_i = m_i$
- Design: use hierarchical designs!

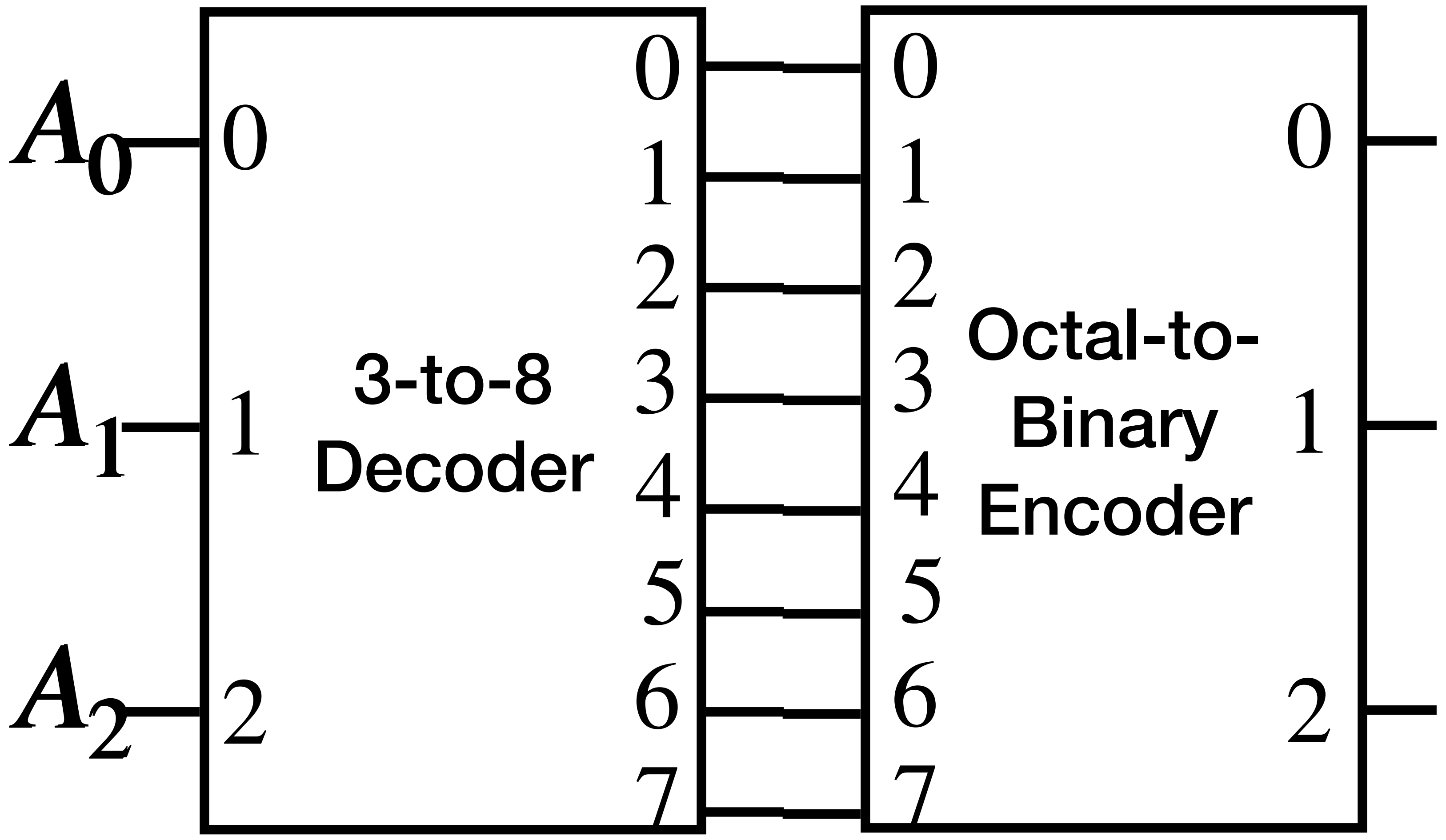
A_1	A_0	D_0	D_1	D_2	D_3
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1



Concept

Encoder

- Inverse operation of a decoder
- 2^n inputs, only one is giving positive input[1]
- n outputs



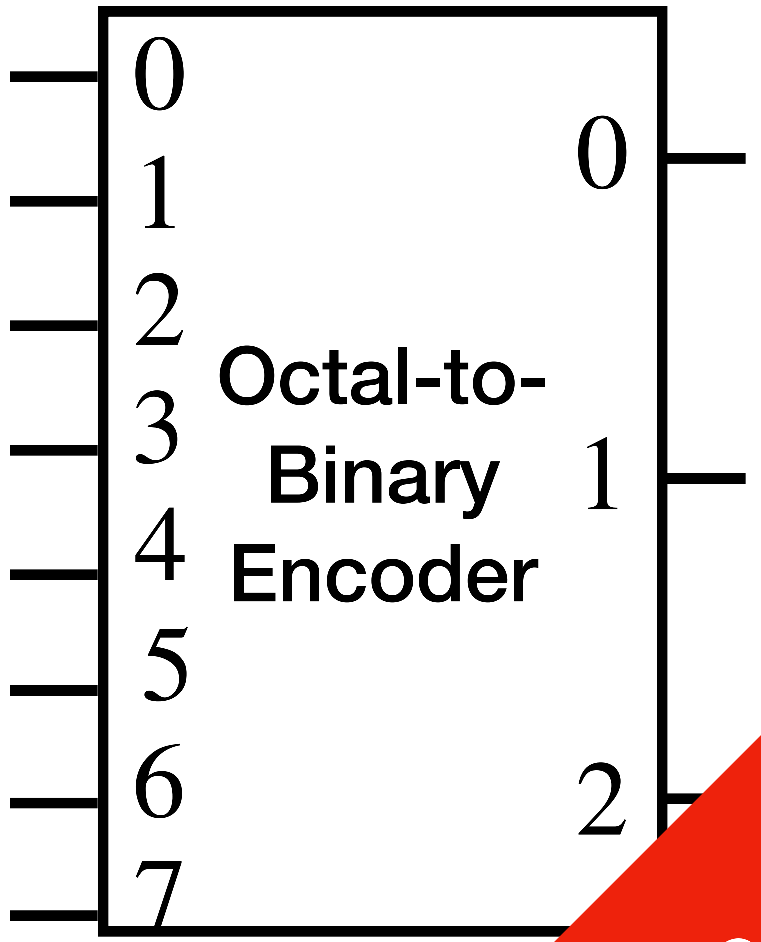
1. In reality, could be less

P3.3
Adv. Func. Blocks

Encoder

D_7	D_6	D_5	D_4	D_3	D_2	D_1	D_0	A_2	A_1	A_0
							1	0	0	0
						1		0	0	1
					1			0	1	0
				1				0	1	1
			1					1	0	0
		1						1	0	1
	1							1	1	0
1								1	1	1

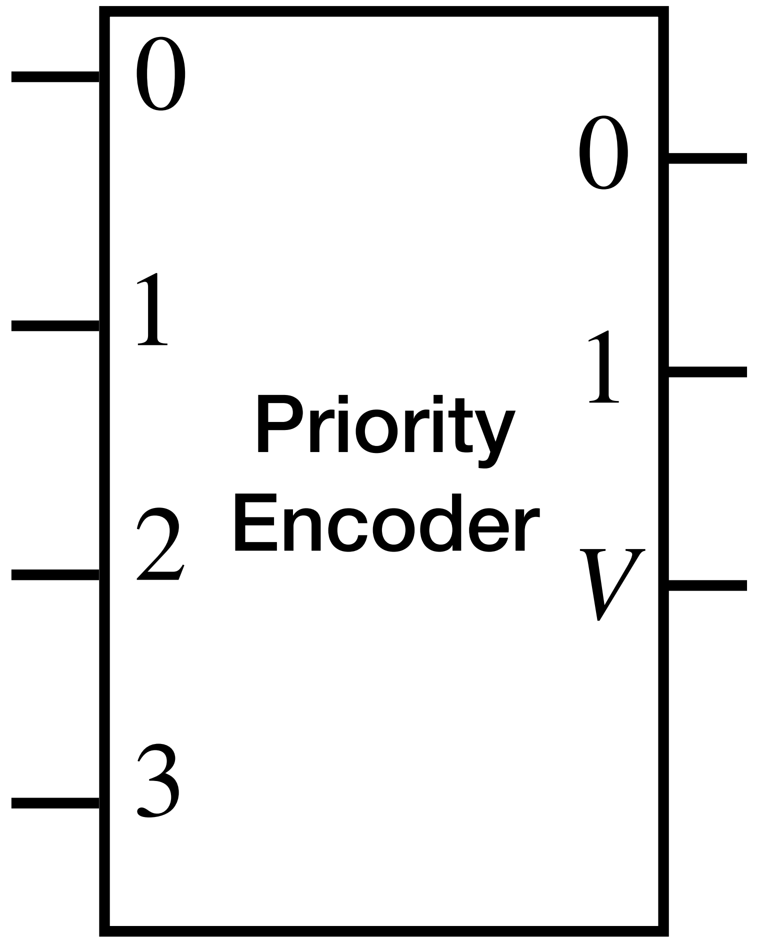
$$A_0 = D_1 + D_3 + D_5 + D_7$$
$$A_1 = D_2 + D_3 + D_6 + D_7$$
$$A_2 = D_4 + D_5 + D_6 + D_7$$



Concept

Priority Encoder

- Additional Validity Output V
  - Indicating whether the input is valid (contains 1)
- Priority
  - Ignores $D_{<i}$ if $D_i = 1$



Concept

P3.3
Adv. Func. Blocks

Priority Encoder

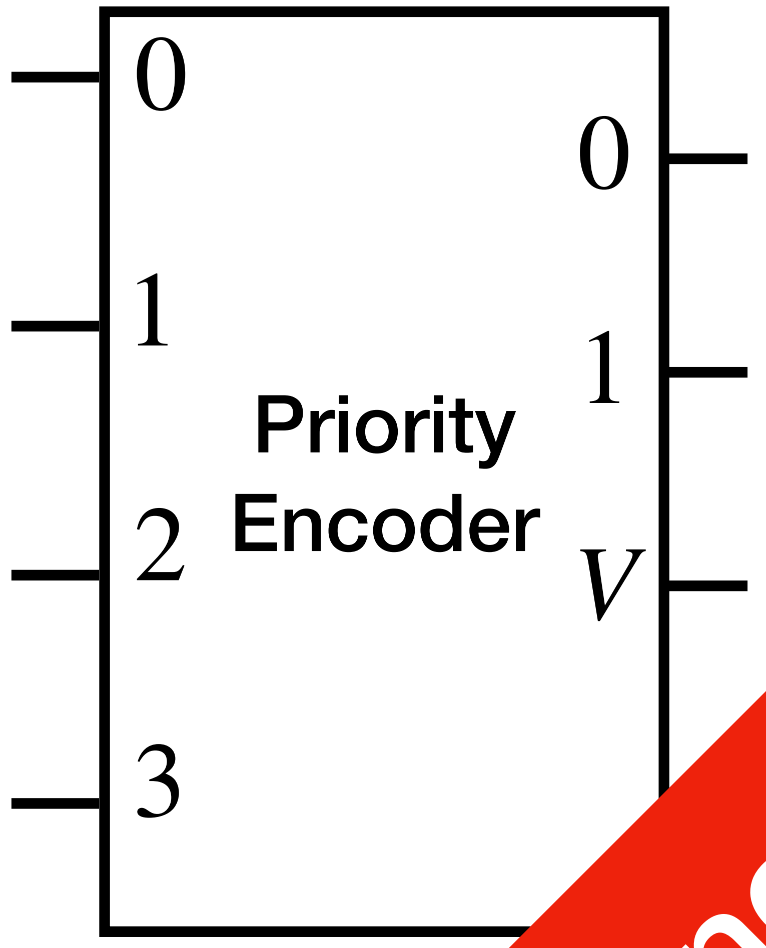
D_3	D_2	D_1	D_0	A_1	A_0	V
0	0	0	0	0	0	0
0	0	0	1	0	0	1
0	0	1	X	0	1	1
0	1	X	X	1	0	1
1	X	X	X	1	1	1

$$V = D_3 + D_2 + D_1 + D_0$$

$$A_1 = D_3 + \overline{D_3}D_2 = D_2 + D_3$$

$$A_0 = \overline{D_3}\,\overline{D_2}D_1 + D_3$$

$$= \overline{D_2}D_1 + D_3$$



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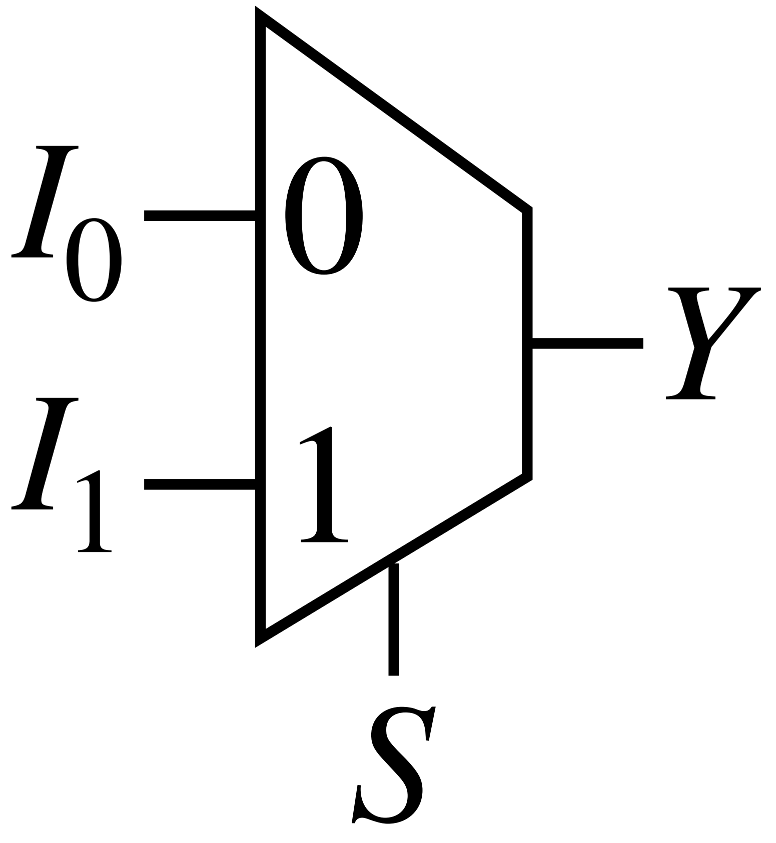
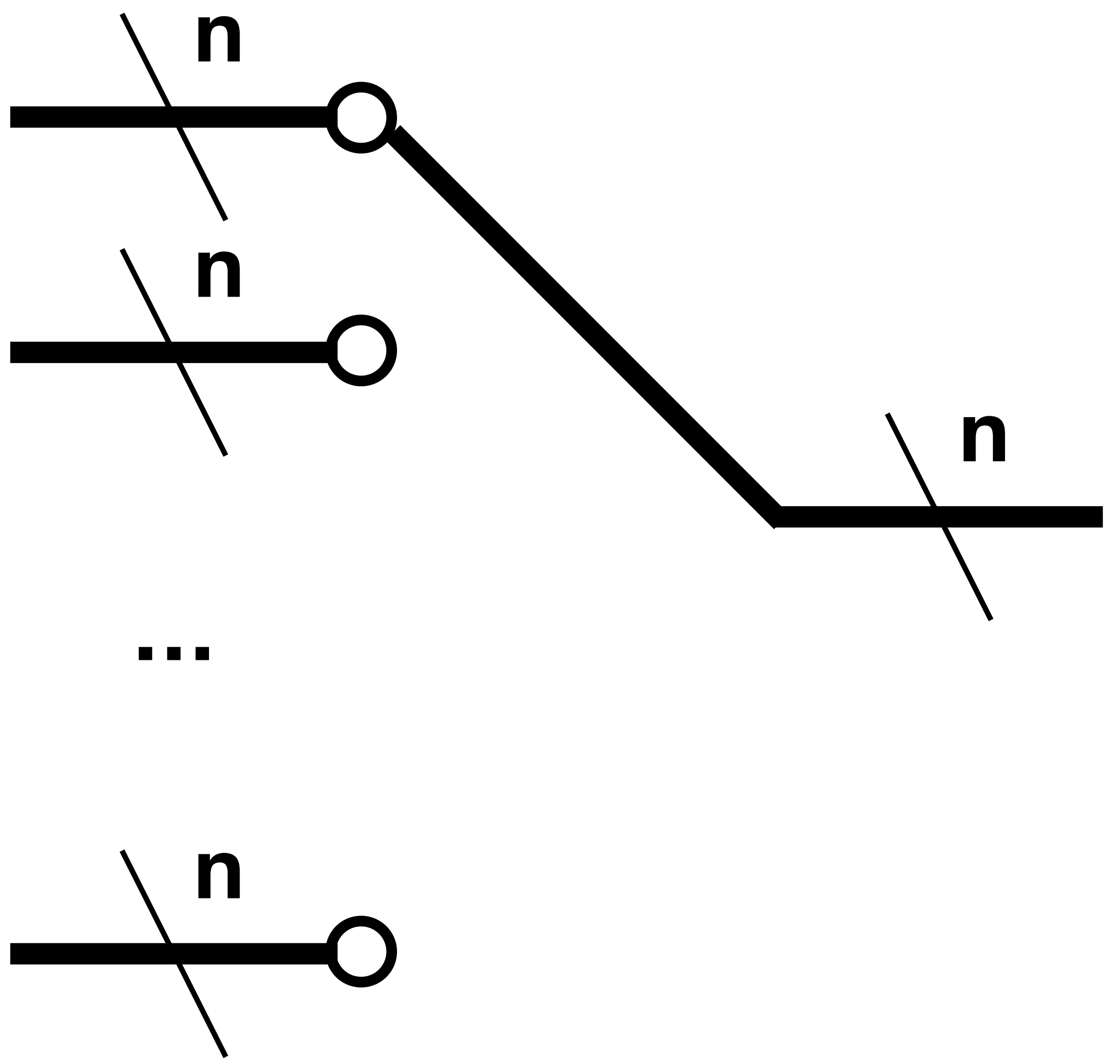
Multiplexer



- Multiple n-variable input vectors
- Single n-variable output vector
- Switches: which input vectors to output



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